## OPTIMAL TRANSMIT STRATEGY FOR MIMO CHANNELS WITH JOINT SUM AND PER-ANTENNA POWER CONSTRAINTS Phuong Le Cao and Tobias J. Oechtering OCH KONST School of Electrical Engineering and ACCESS Linnaeus Centre, KTH Royal Institute of Technology, Stockholm, Sweden Why Joint Sum and Per-antenna Power Constraints? **APPROACHING METHOD** • Sum power constraint (SPC) is imposed by regulations or to limit the • Sequence of optimization problems: energy consumption. • Per-antenna power constraints (PAPC) are imposed by hardware $f(\mathbf{Q})$ limitation of each RF chain Hz] $f(\mathbf{Q})$ • Problems with **SPC** and **PAPC** have been studied by many researchers []pb Intersection points before for both point-to-point and multi-user channels • Both motivations are simultaneously relevant for practical sys $f(\mathbf{Q}) = \max_{\mathbf{Q} \in \mathcal{S}(\mathcal{A})} f(\mathbf{Q}),$ tems, thus we consider a system with a **joint sum and per-antenna** power constraints (JSPC)[1]• Re-assign the antenna coefficient order, form $\mathbf{Q}^{\star} = \begin{bmatrix} \mathbf{Q}_P & \mathbf{q}^H \\ \mathbf{q} & \mathbf{Q}_S \end{bmatrix}$ $\neg \nabla - P_{tot} = \hat{P}_1 = \hat{P}_2 = \hat{P}_3 = 25$ 5.4 MIMO Channels with JSPC $P_{tot} = \hat{P}_2 = \hat{P}_3 = 25$ • Finding the remaining optimal power allocation in $\mathbf{Q}_{S}$ by solving the reduced **MIMO** Transmission $- \diamond - P_{tot} = \hat{P}_2 = 25, \ \hat{P}_3 = 10$ optimization problem $- \star - P_{tot} = 25, \hat{P}_2 = 7, \hat{P}_3 = 10$ Joint sum and per-antenna power constraints $\in \mathcal{S}(\mathcal{P}(k))$ $P_{tot,A} = \min(\hat{P}_1, \hat{P}_2), \ P_{tot,B} = \sum_{i=1}^2 \hat{P}_i$ $\hat{P}_1$ [Watt] $\mathbf{H} \in \mathbb{C}^{m \times n}$ a) $P_{tot} < P_{tot,A}$ (SPC) using the generalized water-filling solution[2]: b) $P_{tot} > P_{tot,B}$ (PAPC) • The more restricted per-antenna power constraint, the less optimal transc) $P_{tot,A} \leq P_{tot} \leq P_{tot,B}$ ∕⁺Lemma 2 mission rate since adding more per-antenna power constraints, we have less The optimal solution of the transmit strategy $\mathbf{Q}^{\star}(k)$ in k-th iteration freedom to allocate the optimal transmit power $\mathbf{Q}^{\star}(k) = (\mathbf{D}^{-\frac{1}{2}}[\mathbf{U}]_{:,1:L}[\mathbf{U}]_{:,1:L}^{H}\mathbf{D}^{-\frac{1}{2}} - [\mathbf{U}]_{:,1:L}\mathbf{\Lambda}^{-1}[\mathbf{U}]_{:,1:L}^{H})^{+},$ • Intersection points: points where the power constraints on some antennas change their state from active to inactive. Definition where $\Lambda$ and $[\mathbf{U}]_{:,1:L}$ are obtained from eigenvalue decomposition $\mathbf{H}^{H}\mathbf{H}$ . The diag-• $\mathcal{S}(\mathcal{A}) := \{ \mathbf{Q} \succeq 0 : \operatorname{tr}(\mathbf{Q}) \leq P_{tot}, P_i = \}$ $\mathbf{e}_i^T \mathbf{Q} \mathbf{e}_i \leq \hat{P}_i, \forall i \in \mathcal{A} := \{1 \dots n\}\}$ onal elements of $L \times L$ , $L = \min(n, m)$ , diagonal matrix $\Lambda$ are positive real values • $f(\mathbf{Q}) := \log |\mathbf{I}_m + \mathbf{H}\mathbf{Q}\mathbf{H}^H|$ in decreasing order. CONCLUSIONS **?** Optimization Problem ALGORITHM $\max_{\mathbf{Q}} f(\mathbf{Q}), \quad \text{s.t.} \mathbf{Q} \in \mathcal{S}(\mathcal{A}).$ • Optimal transmit strategy in closed-form using generalized water-filling $k = 1, \mathcal{P}(1) = \emptyset$ • An unconstraint optimal power allocation of an antenna exceeds a perantenna power constraint, then it is optimal to allocate the maximal power in the constraint optimal transmit strategy including the per-antenna power Compute $\mathbf{Q}^{\star}(k)$ PROPERTIES constraints 2 Propisition 1 • Highly relevant and interesting for massive MIMO since the results might be The maximum transmission rate $R^{\star}$ can be achieved when the optimal transmit strategy $\mathcal{P}(k+1) = \mathcal{P}(k) \cup \{i \in \mathcal{P}^{\mathsf{c}}(k) : P_i(k) > \hat{P}_i\}$ extended to have power constraints for groups of antennas which are driven $\mathbf{Q}^{\star}$ uses full power $P_{tot}$ by one amplifier which has an own power constraint. $\bigstar$ Lemma 1 $k \leftarrow k+1$ Let $\mathcal{A}' \subseteq \mathcal{A}, \ \mathcal{S}(\mathcal{A}') := \{ \mathbf{Q} \succeq 0 : \operatorname{tr}(\mathbf{Q}) \leq P_{tot}, P_j^{\mathcal{S}(\mathcal{A}')} = \mathbf{e}_j^T \mathbf{Q} \mathbf{e}_j \leq \hat{P}_j, j \in \mathcal{A}' \}$ , and $\mathcal{P} := \{ i \in \mathcal{A}^{\mathsf{c}} : P_i^{\mathcal{S}(\mathcal{A}')} > \hat{P}_i \}$ with $\mathcal{A}^{\mathsf{c}} = \mathcal{A} \setminus \mathcal{A}'$ . Then the optimal power can be $\mathcal{P}(k+1) = \mathcal{P}(k)$ References $P_i(k) \leftarrow \hat{P}_i$ $P_{tot}(k) \leftarrow P_{tot} - \sum_{i \in \mathcal{P}(k)} P_i$ [1] P. L. Cao, T.J Oechtering, R.Schaefer and M. Skoglund, "Optimal Transmit Strategy for MISO Channels allocated as with Joint Sum and Per-antenna Power Constraints," IEEE Transactions on Signal Processing, vol.64, no.16, pp.4296-4306, Aug 2016. $\mathbf{Q}^{\star} \leftarrow \mathbf{Q}^{\star}(k)$ [2] C. Xing, Z. Fei, Y. Zhou, and Z. Pan, "Matrix-field Water-filling Architecture for MIMO Transceiver Designs

$$\begin{cases} P_i^{\star} = P_i^{\mathcal{S}(\mathcal{A}')}, \, \forall i \in \mathcal{A}^{\mathsf{c}} \text{ if } \mathcal{P} = \emptyset \\ P_i^{\star} = \hat{P}_i, \quad \forall i \in \mathcal{P} \quad \text{otherwise} \end{cases}$$



$$\max_{\mathbf{Q}(k)} f(\mathbf{Q}(k)) \text{ s.t. } \mathbf{Q}(k) \in$$





- with Mixed Power Constraints," in *PIMRC*, Sep. 2016.

