



**CENTER FOR
CYBER DEFENCE AND
INFORMATION SECURITY**



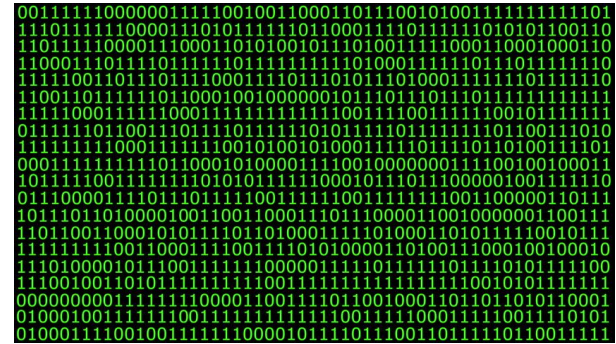
Provable Security
Roberto Guanciaie, KTH
Mads Dam, Henrik Karlsson
CDIS Spring Conference 2022
KTH Royal Institute of Technology
Tuesday, May 24, 2022

Formal methods

- build our systems in new ways
 - more rigorous techniques.
- **prove** using mathematics and logic the **impossibility** of vulnerabilities in systems

Heavy duty, but remarkable successes for critical SW

Key is to analyse the critical part of the system and guarantee that the other parts do not compromise these properties

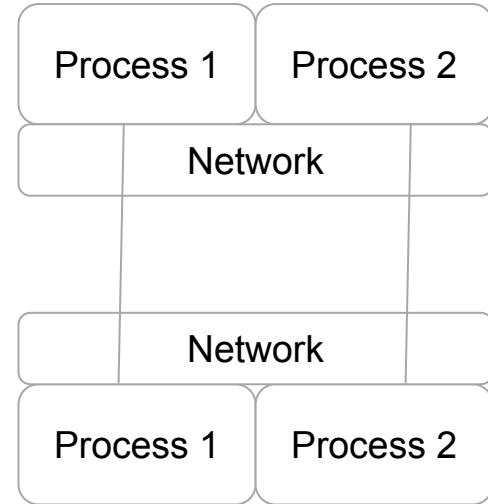


Formal methods

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Let's look at an autonomous grass cutter

WiFi

FTP

Motor Control

Voice command

Battery

Schedules

Proximity sensor

Mapping

Bluetooth

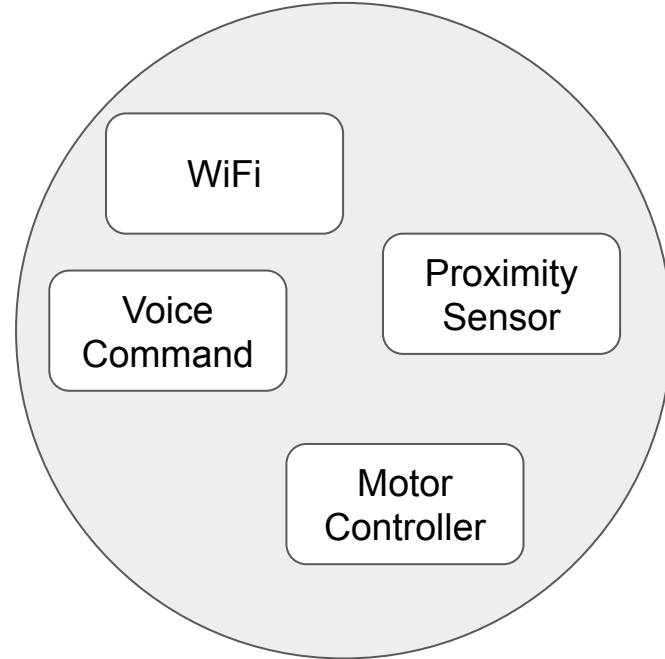


Camera

- varying quality
 - varying trust
 - different supply chains
-

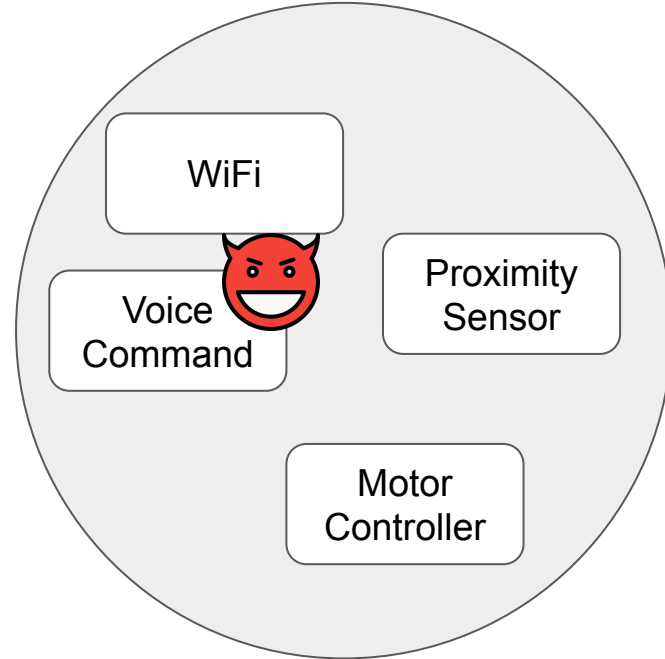
Key problem for security

Complexity



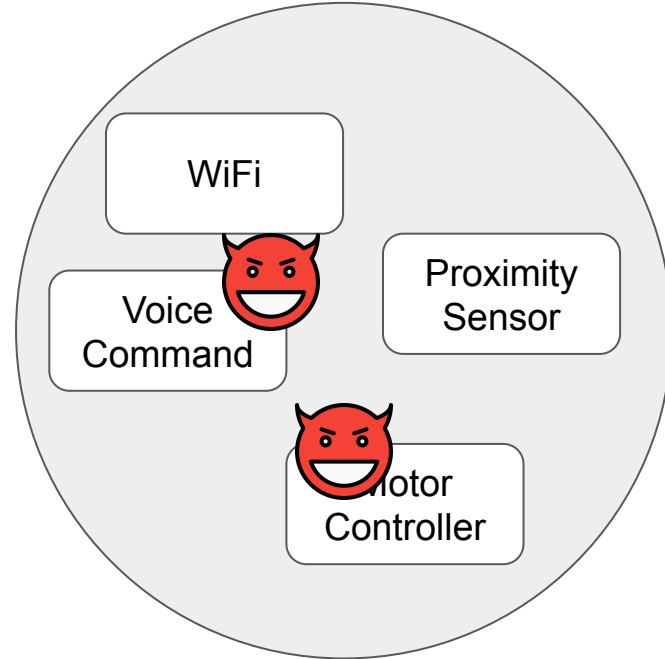
Key problem for security

Complexity



Key problem for security

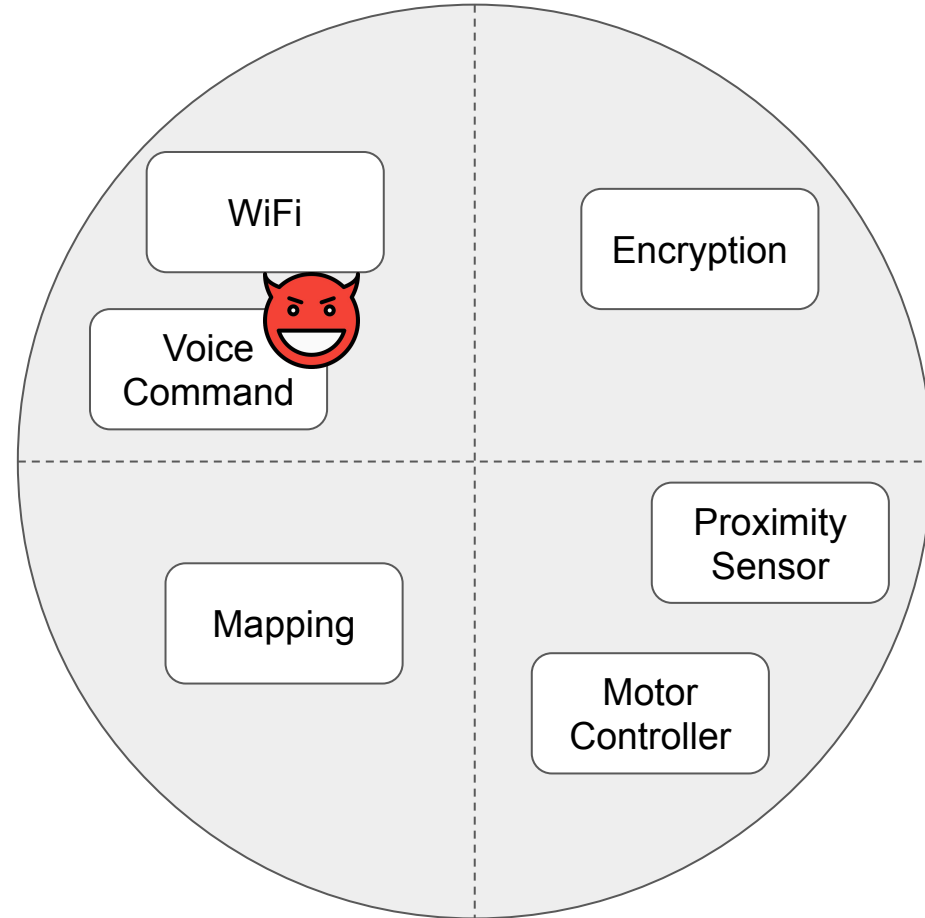
Complexity



Standard OSs (Linux)

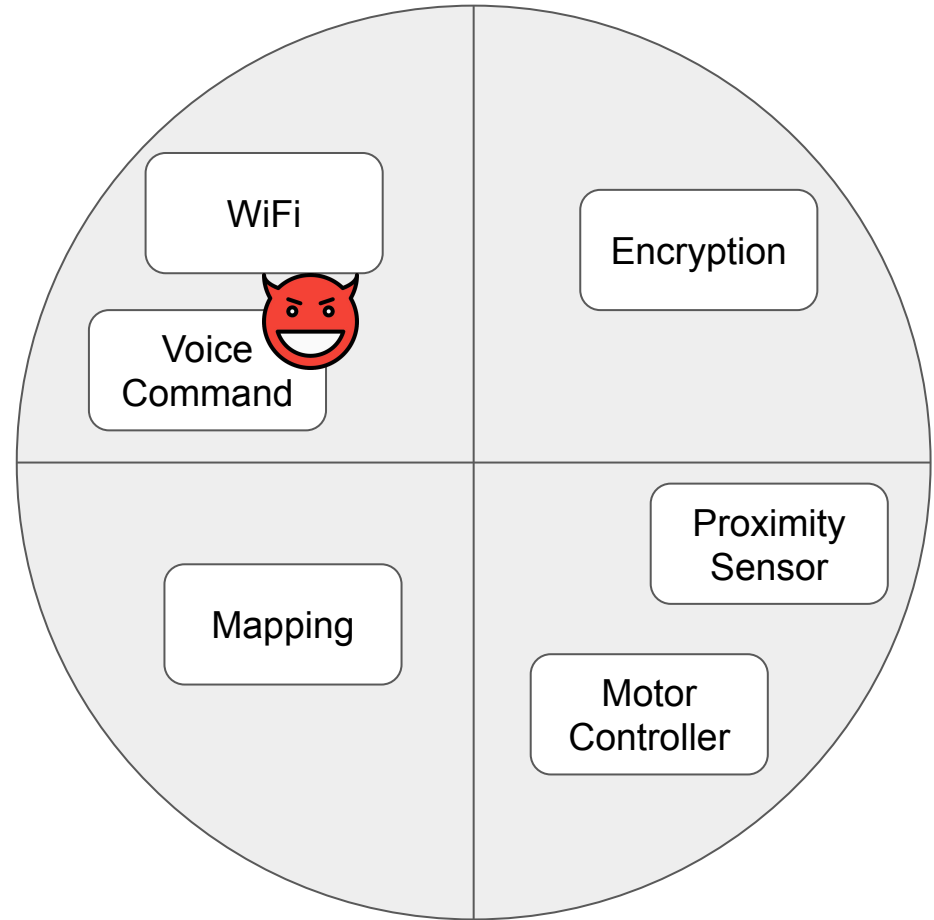
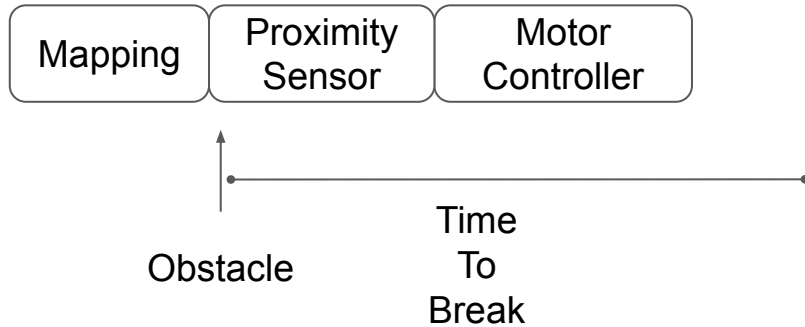
Notoriously buggy

- 27.8 million lines of code
- 75 thousands code commits/year
- Linus Torvalds 3.19% of commits
- 4,189 different contributors



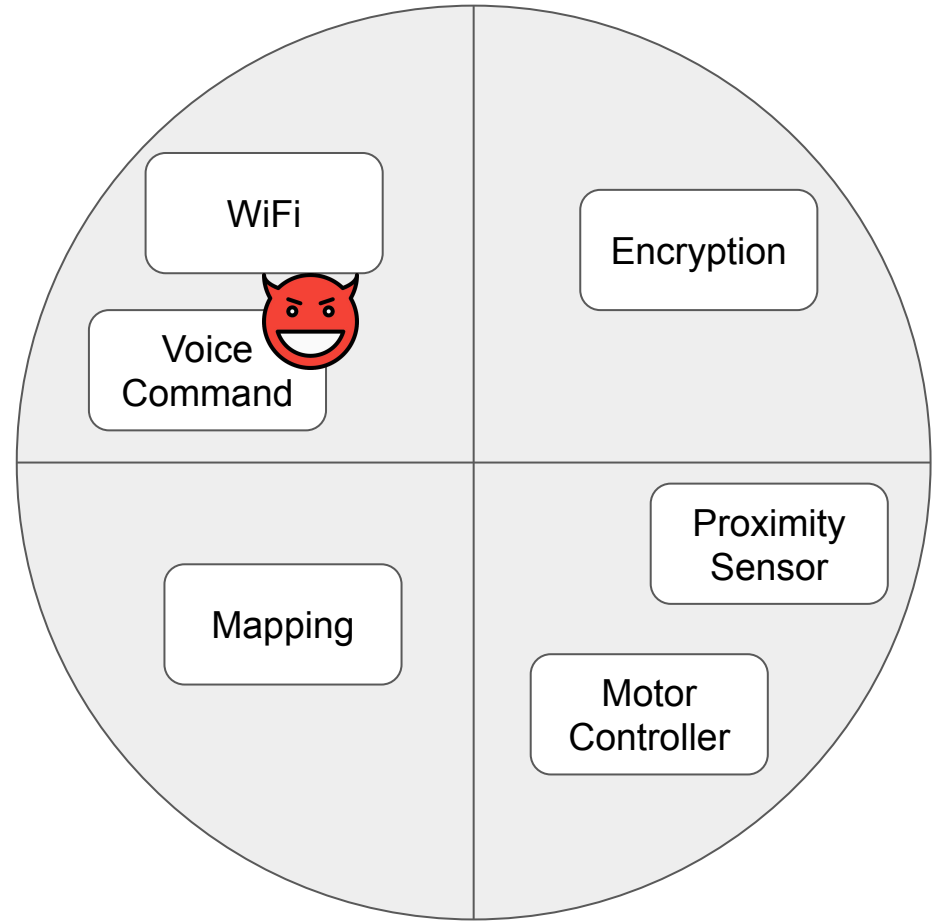
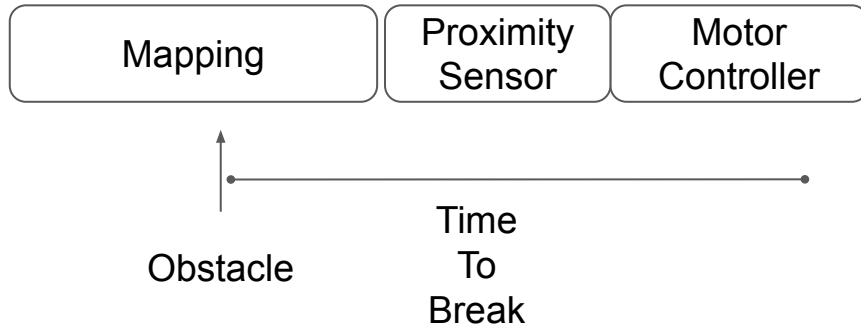
Microkernels

- Functionally correct
- Strong spatial isolation
- NO non-functional guarantees



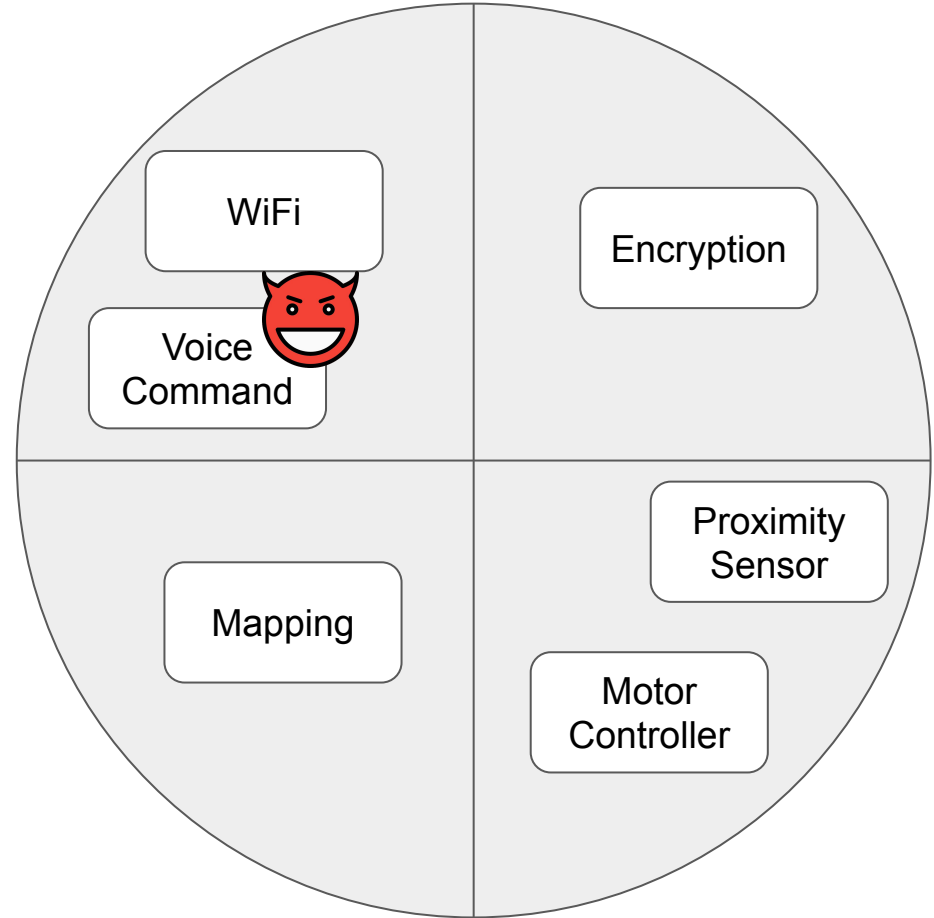
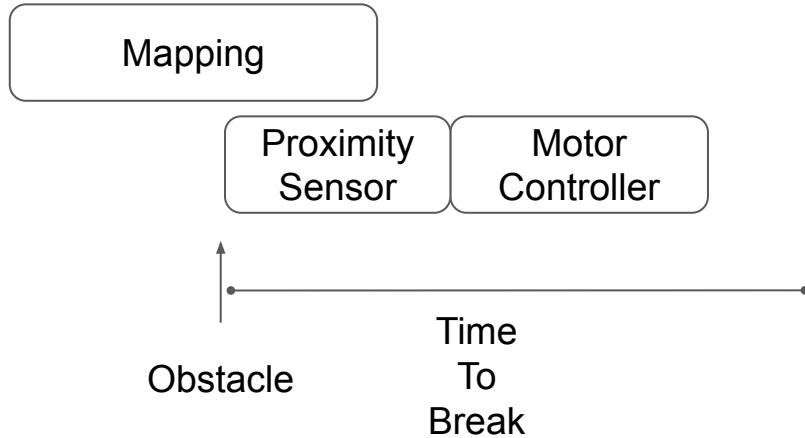
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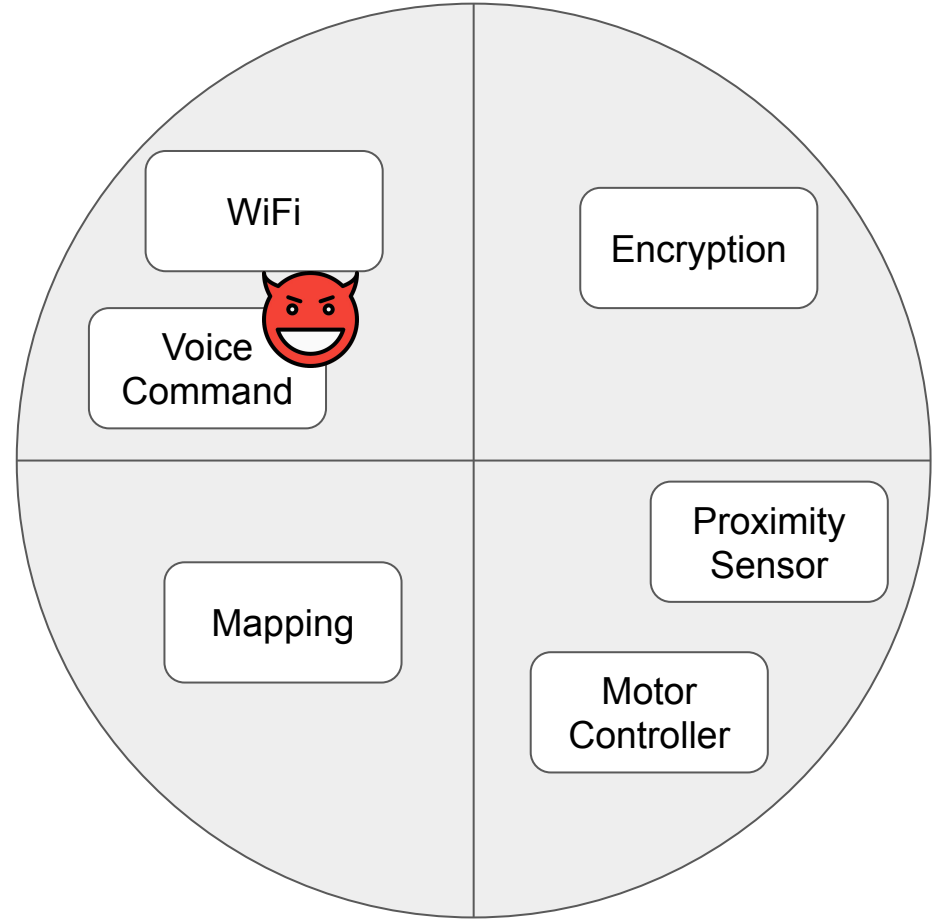
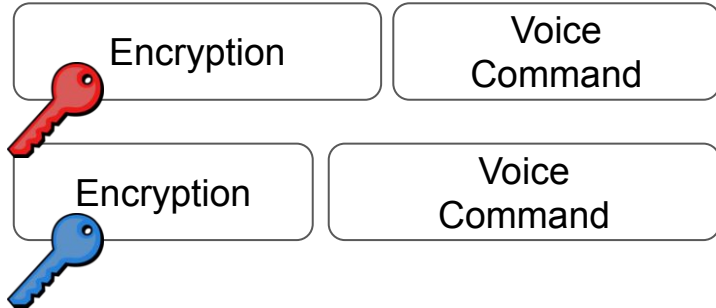
Real Time Systems

- Priorities to meet deadlines
- No protection against side channel



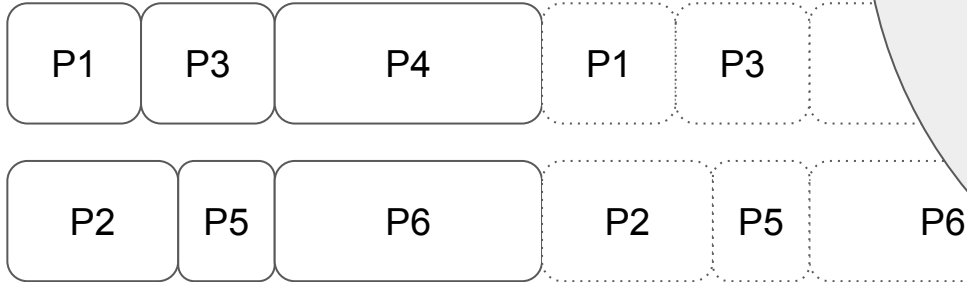
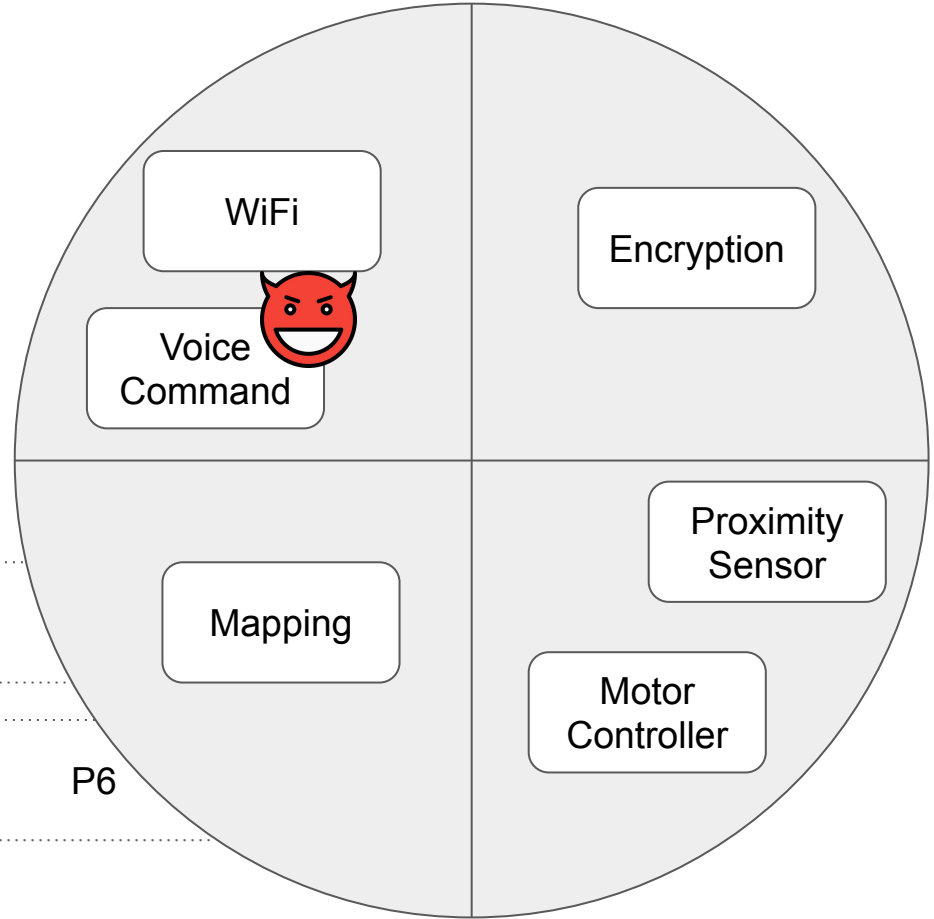
Real Time Systems

- Priorities to meet deadlines
- No protection against side channel



Separation Kernel

- Both strong spatial and temporal isolation
- Used in Avionics (ARINC 653)
- Everything is static



Minor Frame

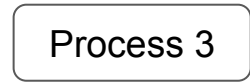
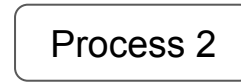
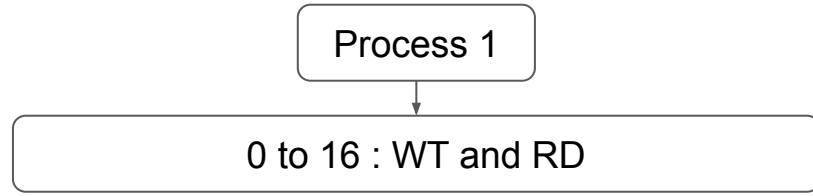
Mayor Frame

A new model for separation kernels

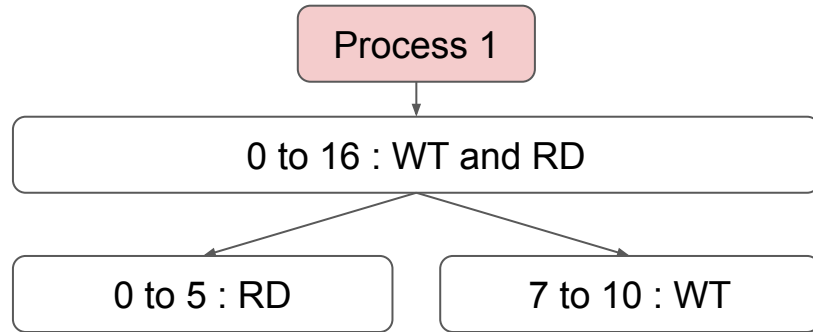
Use the capability model for all types of resources:

- Process management
- Process communication
- Memory
- Time

Capability model



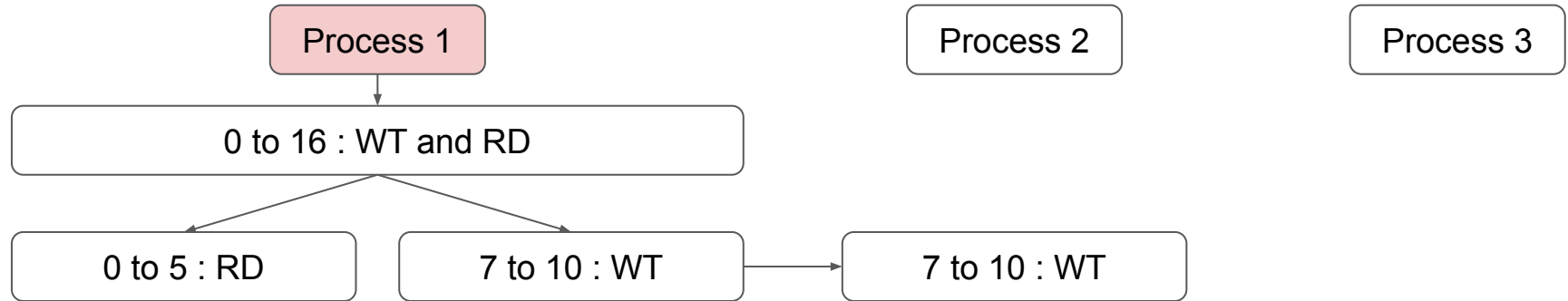
Capability model: derivation



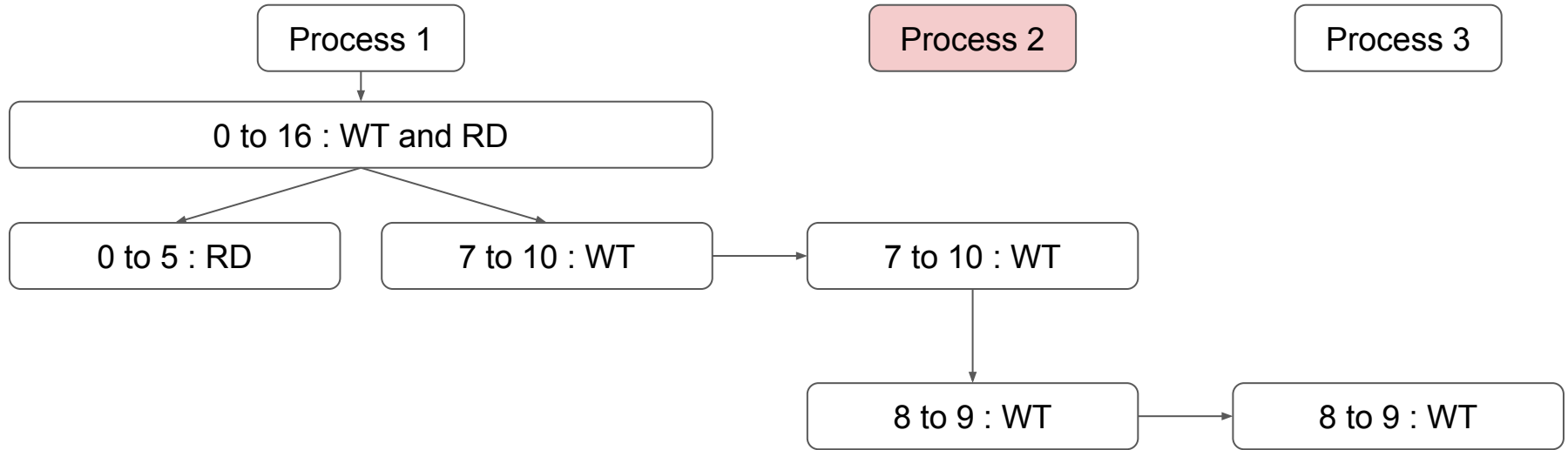
Process 2

Process 3

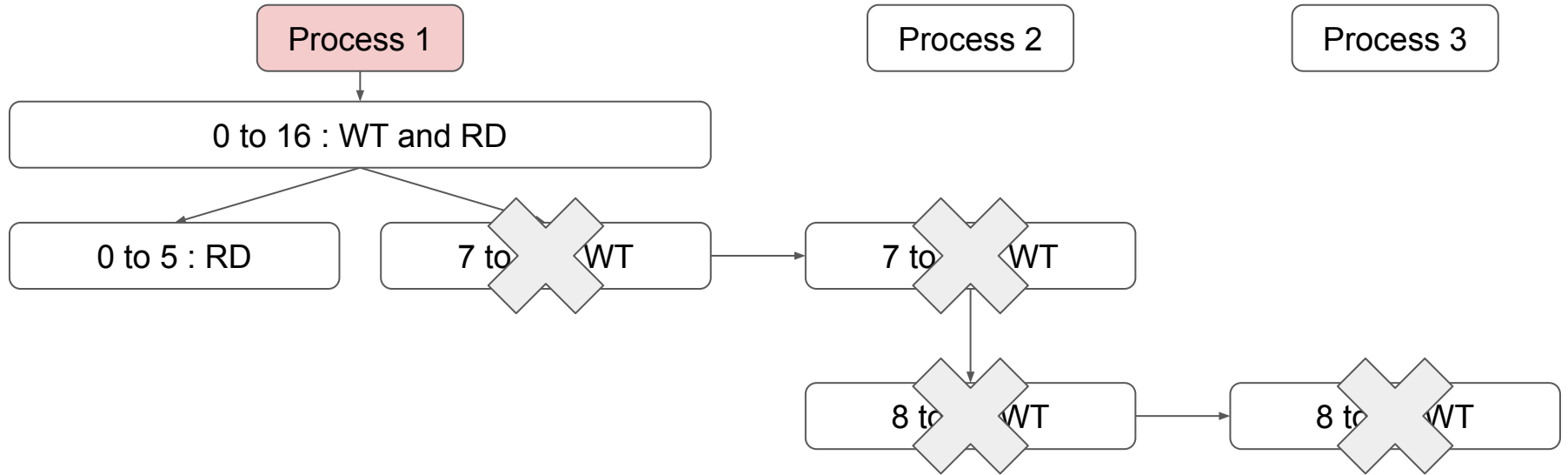
Capability model: delegation



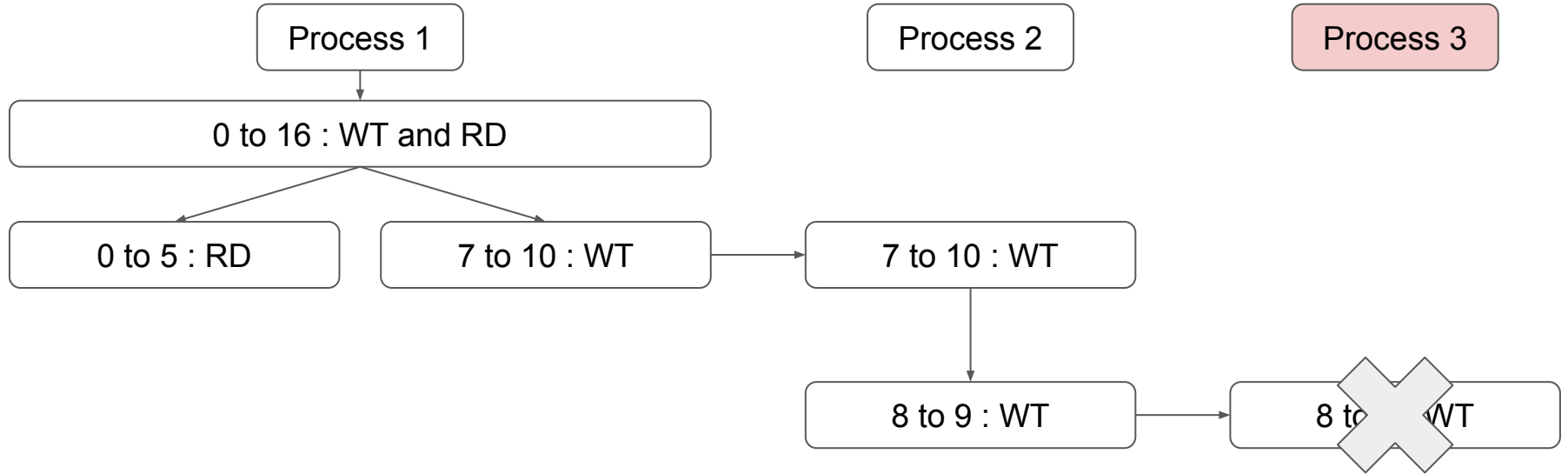
Capability model: delegation



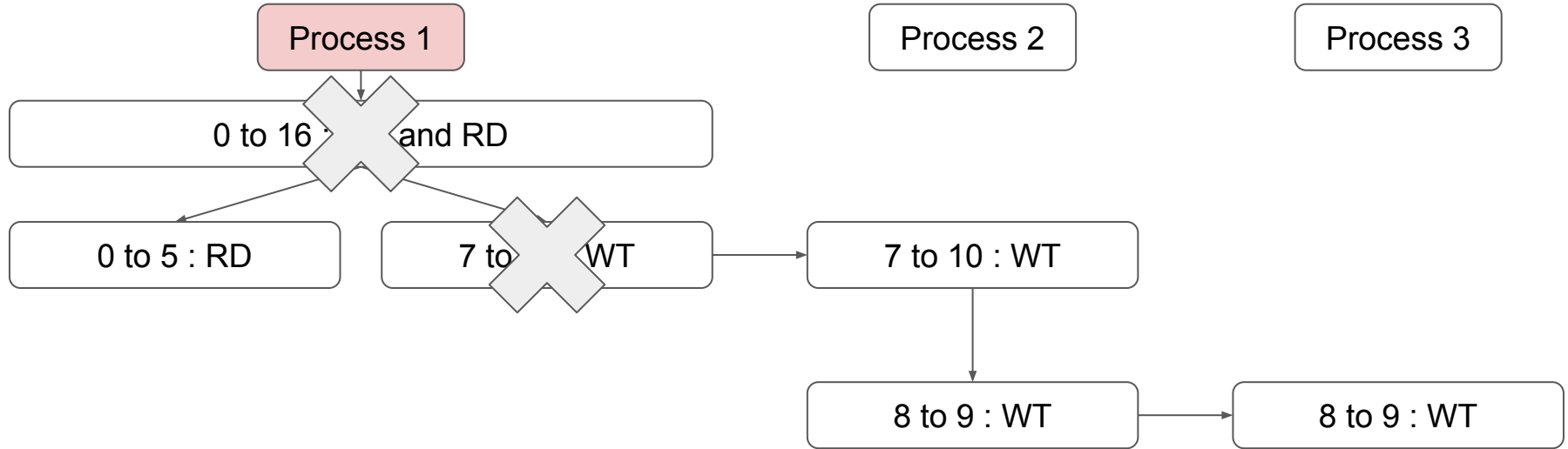
Capability model: revocation



Capability model: deletion



Capability model: drop



Timing Capability



Mayor Frame

P1



0 to 100

P2

Timing Capability

P1

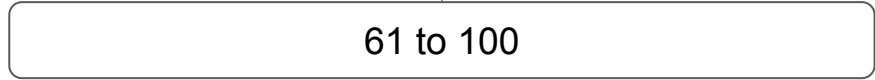
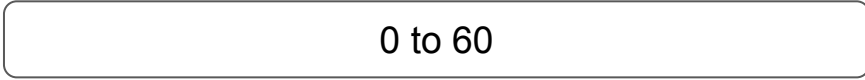
0 to 60

61 to 100

Mayor Frame

P2

61 to 100

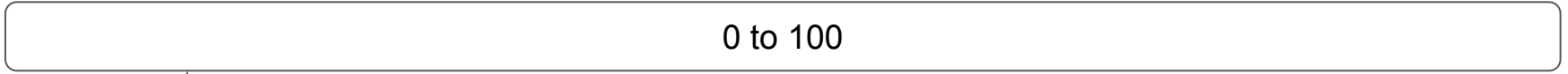


Few APIs but expressive

- Timing Capabilities are per core.
- A process can run only on one core, but it can migrate.
- Kernel is (mostly) preemptive. We use slack time that should cover the non-preemptive parts of the kernel
- 2kLoC of C/Assembly for 64-bit RISC-V

Remote Procedure Call

P1



0 to 100

P2



Remote Procedure Call: Derive

P1



0 to 100

P2



0 to 90

91 to 100



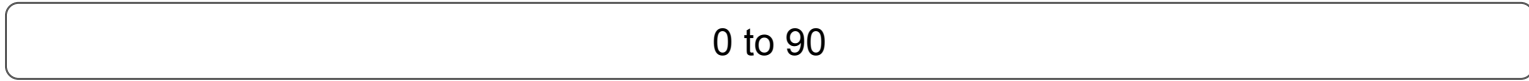
Remote Procedure Call: Delegate

P1



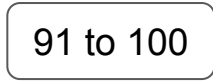
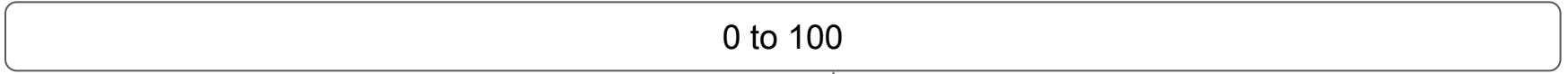
91 to 100

P2



Remote Procedure Call: Delete

P1



P2

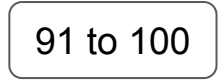
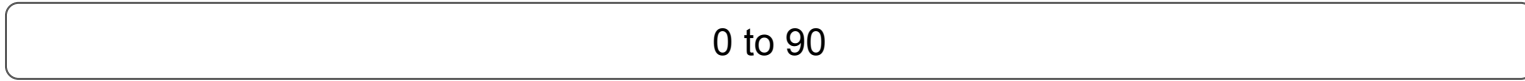


Remote Procedure Call: Revoke

P1

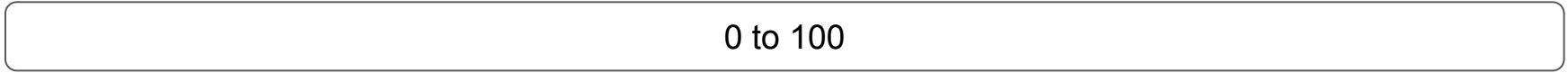


P2

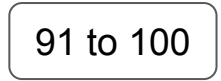


Remote Procedure Call: Revoke

P1



P2



Ongoing work

- WCET analysis of non-preemptive kernel
- Demonstrating application:
 - Cyber physical system with multiple modes of operations
 - Reconfigurations to handle faults
- Userland (Posix) and drivers
- Formal verification of kernel system calls
 - Model of RISC-V memory model
- Formal verification of non-functional properties
 - WCET
 - Constant time execution