Extending Dataflow Streaming for Stateful Serverless

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Philipp Haller: Background

- Associate professor at KTH (2014–2018 Assistant professor)
 - PhD 2010 EPFL, Switzerland
- 2005–2014 Scala language team **=**
 - 2012–2014 Typesafe, Inc. (now Lightbend, Inc.)
 - Co-author Scala language specification
- Focus on *concurrent and distributed programming*
 - Creator of Scala Actors, co-author of Scala's futures and Scala Async

2019: ACM SIGPLAN Programming Languages Software Award for Scala Core contributors: Martin Odersky, Adriaan Moors, Aleksandar Prokopec, Heather Miller, Iulian Dragos, Nada Amin, Philipp Haller, Sebastien Doeraene, Tiark Rompf







Scala Actors and Akka

LEASES

JUN21

2006

JUN5

2007

OCT 13 2007

DEC 19

2007

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Philipp Haller creates Scala Actors (the original standard library actors). His work becomes a major influence to Akka and the main reason for Jonas Bonér to choose Scala as the platform for Akka.

First commit: https://github.com/scala/scala/commit/0d8b14c6055e76c0bff3b65d0f428d711abe1f5a

CONFERENCE MPER Actors that Unify Threads and Events

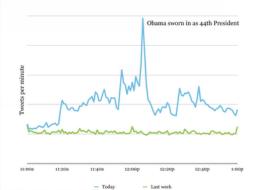
Haller, Philipp: Odersky, Martin Editors: Vitek, Jan: Murphy, Anty L

Presented at: International Conference on Coordination Models and Languages, Paptos, Cyprus, 5-8 June 2007 Published itt: Proceedings of the Shi International Conference on Coordination Models and Languages (COORDINATION), p. 171-190 Series: Letture Notes in Computer Science (LNCS) 4467 : Somarr, 2007

There is an impedance mismatch between message-passing concurrency and virtual machines, such as the JVM. Whis usually map their threads to heavyweight OS processes. Without a lightweight process abstraction, users are often forced to write parts of concurrent applications in an evert-driven

Scala Actors used, e.g., in core message queue system of Twitter:

Philipp Haller publishes his influential paper on Scala's Actors; 'Actors that Unify Threads and Events'. http://infoscience



https://www.lightbend.com/akka-five-year-anniversary^{19/hotswap-cod}

ERLANG http://

Scala



Blog post:

http://jonasboner.com/2007/10/30/interview-with-joe-armstrong

Jonas starts tinkering with Scala Actors.

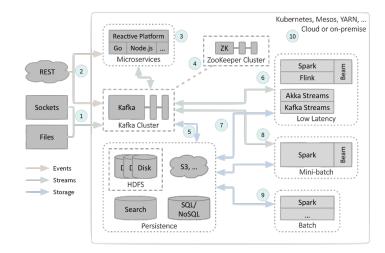
The use of actors is common in industry



All Modern Services are Distributed

Uber (1) (1) (2)

- Each of these systems is a distributed system itself
- User data and services scattered across multiple systems
- This is not suited for classic monolith architectures: microservices architecture to the rescue*



Source: Dean Wampler: Fast Data Architectures For Streaming Applications (2nd edition), O'Reilly

*Till Rohrmann: Keynote: Rethinking how distributed applications are built. DEBS 2022.

F.T.C. Scrutinizes Tech Companies' Data Use: Live Updates

Google's apps crash in a worldwide outage.

Dec 14 2020

Failures in cloud-based distributed systems can be catastrophic.

Internet users worldwide received a jarring reminder on Monday about just how reliant they were on Google, when the Silicon Valley giant suffered a major outage for about an hour, sending many of its most popular services offline.

At a time when more people than ever are working from home because of the pandemic, Google services including Calendar, Gmail, Hangouts, Maps, Meet and YouTube all crashed, halting productivity and sending angry users to Twitter to vent about the loss of services. Students struggled to sign into virtual classrooms.

Reality of Distributed Systems

- Reliability: computers crash, messages get lost
- Scalability: workloads increase or decrease
- **Cloud and edge:** execution in heterogeneous environments
- **Response time:** services require low latency
- **Privacy:** systems manage sensitive, regulated data (GDPR, CCPA)

We are asking too much of distributed software programmers!

Limitations of Distributed Programming Models

		Distributed Programming Patterns					Guarantees		Distributed Execution		
		Cyclic Dependencie s	Dynamic Communication Topology	Dataflow Composition	Typed Communication	Request/ Reply with Futures	Exactly-once Processing	Serializable Updates	Decentralize d Deployments	Data Parallelism	Task Parallelism
	Dataflows	-	-	Х	Х	-	Х	-	-	Х	Х*
	Actors	х	Х	X*	X*	Х	-	-	Х	-	Х
Stateful Functions	Stateful Serverless	X*	Х	-	Х	Х*	Х	-	Х	Х	Х

* Supported with restrictions

J Spenger, P Carbone, P Haller. *Portals: an Extension of Dataflow Streaming for Stateful Serverless.* 2022, preprint

Philipp Haller

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	Actors	х	Х	Х*	Χ*	Х	-	-	Х	-	Х
Stateful Functions	Stateful Serverless	X*	Х	-	Х	X*	Х	-	Х	Х	Х

* Supported with restrictions

No current programming system is well-equipped for the complete job!

The Stateful Serverless Dream

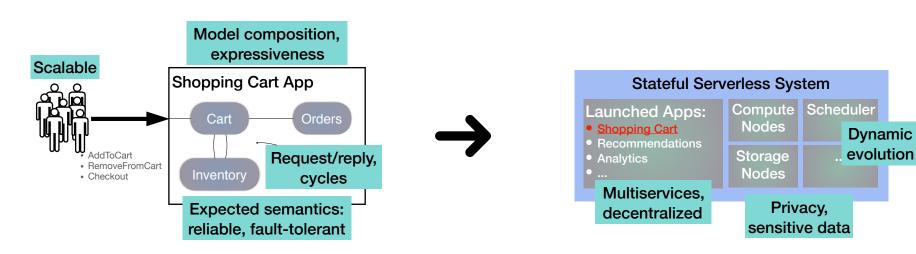
- The programmer should only need to write business logic
- The stateful serverless system should automate everything else:
 - **Reliability:** exactly-once-processing guarantees
 - Scalability: scale up and down with demand
 - **Execution:** cloud, edge, performance, latency
 - **Privacy:** primitives for handling sensitive data



Automatically:

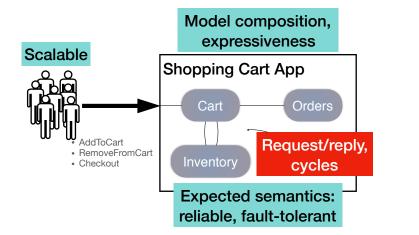
- End-to-end processing guarantees: checkpointing, recovery
- Manage running applications
- · Manage multiple, decentralized deployments
- Scale up/down, dynamic reconfiguration
- Handle requests for live application updates, privacy

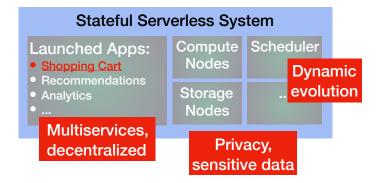
Requirements and challenges



Requirements and challenges

Requirements not supported by Dataflow Streaming





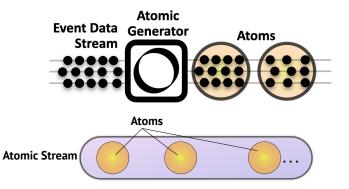
Extending Dataflow Streaming for Stateful Serverless

The **Portals programming model** introduces new abstractions:

- Atomic Streams
- "Portals"
- Workflows
- Live consistent updates (serializable)

Atomic Streams

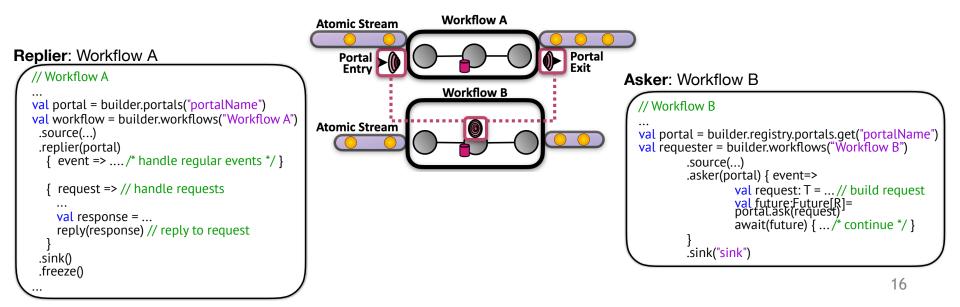
- Totally ordered, distributed stream of atoms.
 - Atom: Sequence of events, transactional unit of computation.
- Atomic Streams enforce end-to-end exactly-onceprocessing guarantees.
 - The Atomic Processing Contract: "The consumer/producer must always consume and process the whole atom, before consuming and processing the next atom."

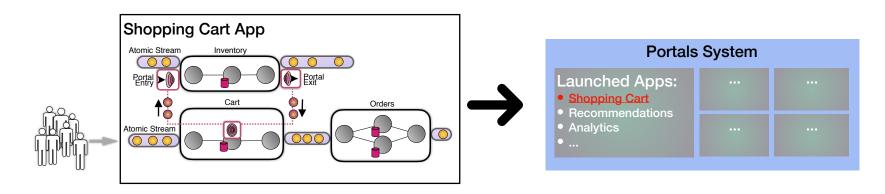


"Portals"

Request-reply style programming with workflows, includes futures API

• Portals enable request/reply, futures



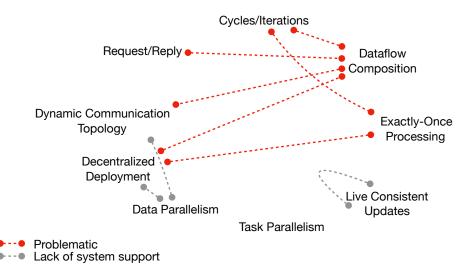


Semantically sound application logic

Fully automated deployment

When to use Portals

Applications that have/need certain combinations that are problematic. Common solution: resort to plumbing together different systems.



Use Cases

- Complex event processing applications
- ML model training and serving
- Dynamic workflow reconfiguration
- Sagas, distributed transactions
- Serializable updates (e.g., for consistent execution of GDPR requests)
- Secure workflows / privacy-preserving computing (future work)

Outlook

Portals programming model

- Express/simulate other distributed programming models in Portals
- Operational semantics and soundness of Portals
- Integration of Secure Multi-Party Computation (future)
- Portals system
 - Exploit use-cases
 - Performance optimization & evaluation
 - Release Portals 1.0: distributed, decentralized runtime
 - Sign up for launch at <u>www.portals-project.org</u>



Sign up for the launch at

Summary

ME www.portals-project.org



Key takeaways

- Dataflow streaming a great candidate for composing stateful serverless services
 - Not so great for cycles, request/reply-style communication, decentralized dynamic deployments
- The **Portals programming model** extends dataflow streaming:
 - Atomic streams ensure processing guarantees
 over decentralized dynamic deployments
 - **Portals** enable request/reply-style communication with futures

