

On the quantum threat to cryptography, its mitigation, and our quantum cryptanalysis research

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SWEDISH ARMED FORCES

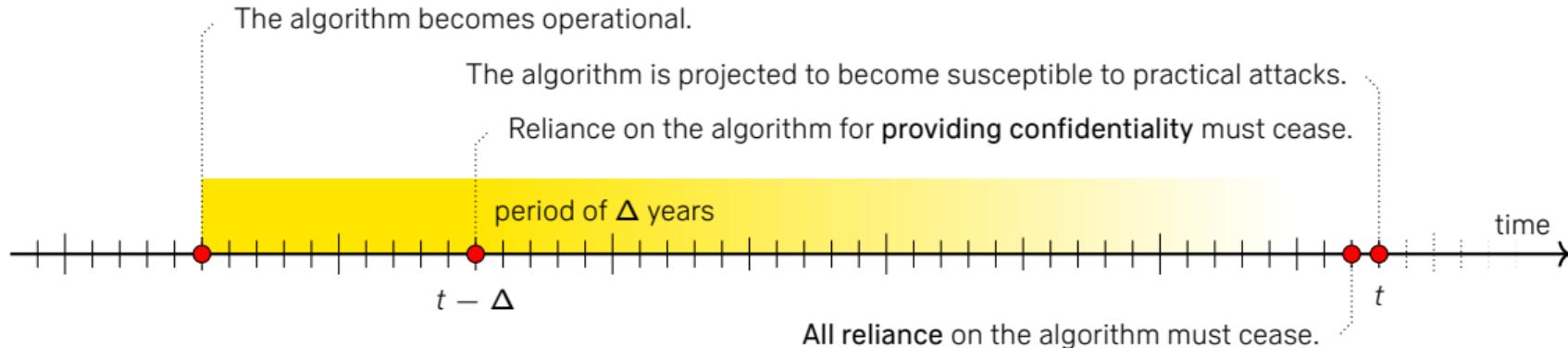


Introduction

Introduction

- ▶ Virtually all historically widely deployed commercial *asymmetric* cryptography will be broken if sufficiently capable quantum computers are built in the future.
- ▶ It is conceivable that such computers may be built sometime after the year 2030.
- ▶ Needless to say, it is very hard to make predictions about the future, but we need to make a prediction to set the time plan for mitigation efforts.

When are mitigating actions required at the latest?



Intermediary periods and confidentiality

- ▶ For plaintexts that we encrypt today to remain confidential for a period of Δ years, the algorithm we rely upon must remain secure for a period of Δ years.
- ▶ Prioritize taking mitigating actions for algorithms that provide confidentiality.

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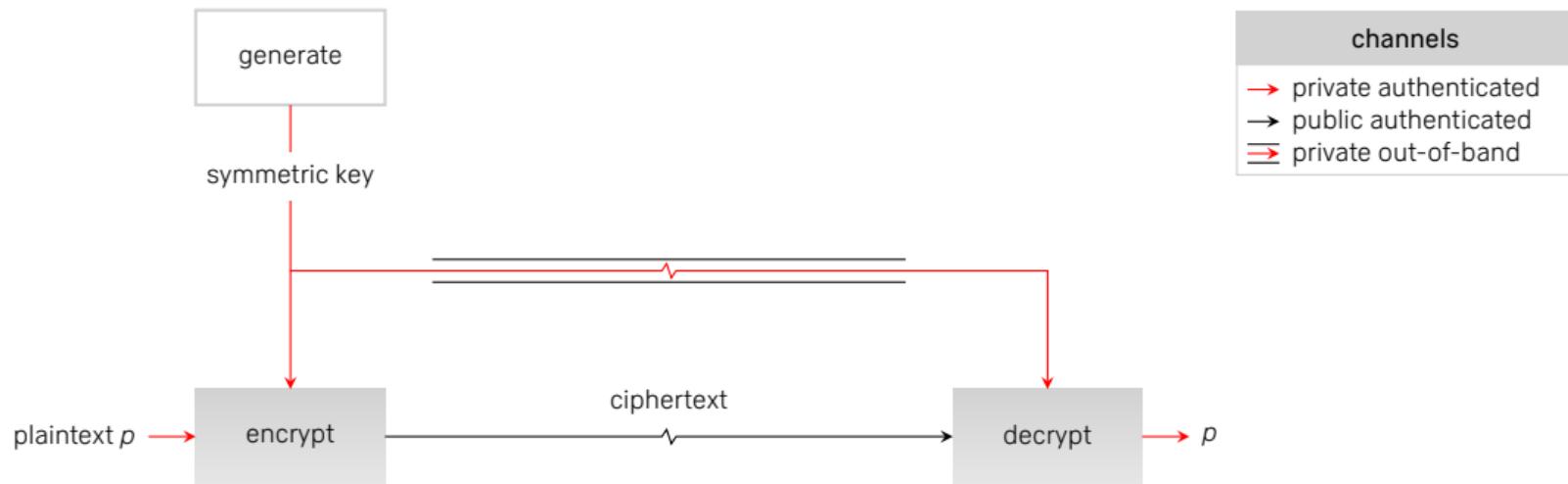
- Symmetric keying
- Asymmetric keying
- Hybrid keying

3. Quantum cryptanalysis

- Our research contributions
- Timeline projections
- Industry roadmaps
- Recent algorithmic developments

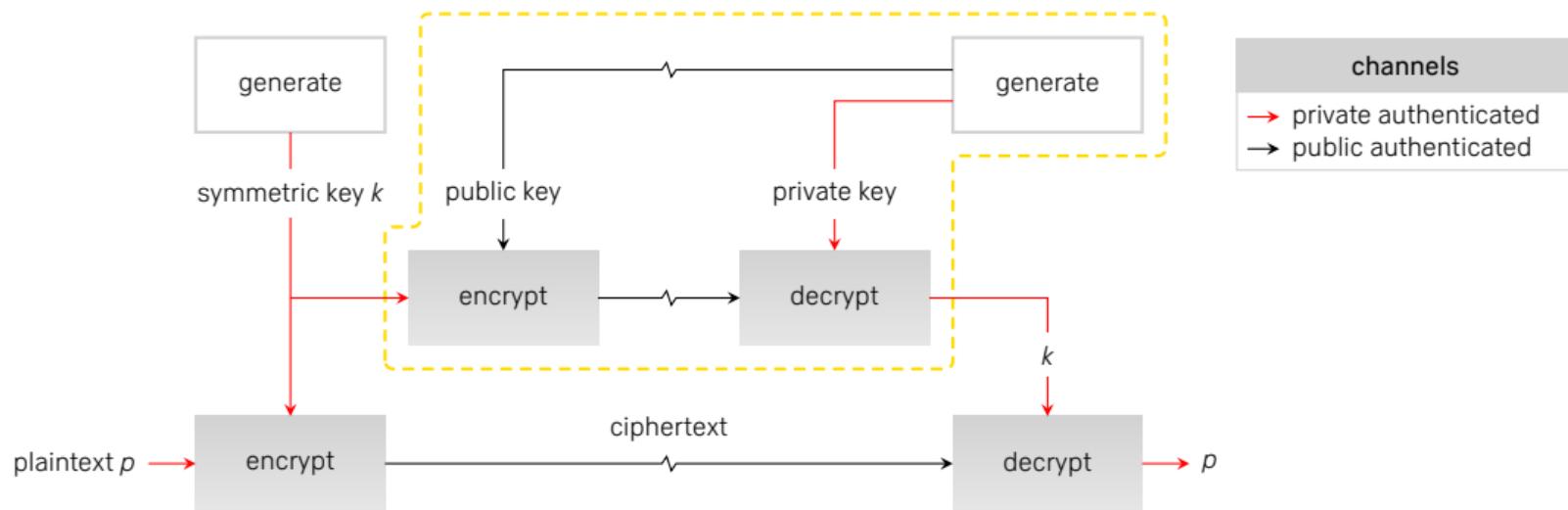
4. Summary and conclusion

Symmetric keying



1. Use symmetric keying, whenever feasible, with secure out-of-band key distribution.
 - ▶ Limit the use contexts and validity periods of keys. Provides robust security, but no forward secrecy (FS). Suitable baseline for closed high-security networks.

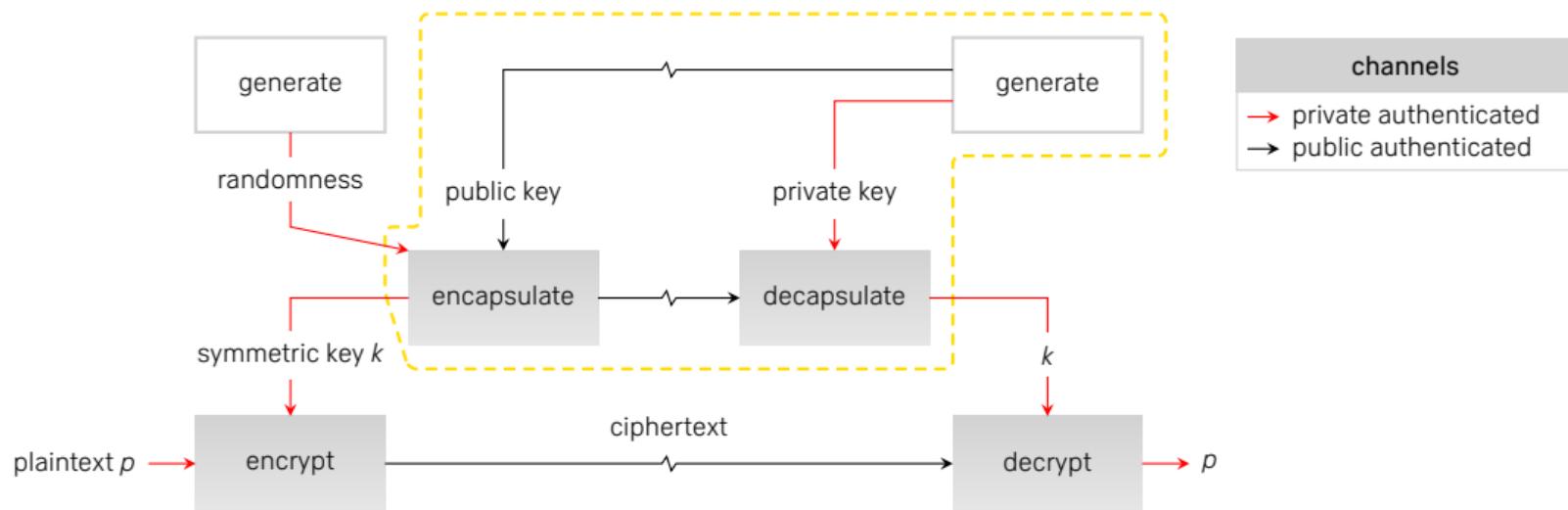
Asymmetric keying via public-key encryption



2. Use post-quantum secure asymmetric keying, e.g. via public-key encryption.

- Less robust than symmetric keying but can provide forward secrecy (FS). Suitable baseline for open networks when symmetric keying is not feasible.

Asymmetric keying via public-key encapsulation

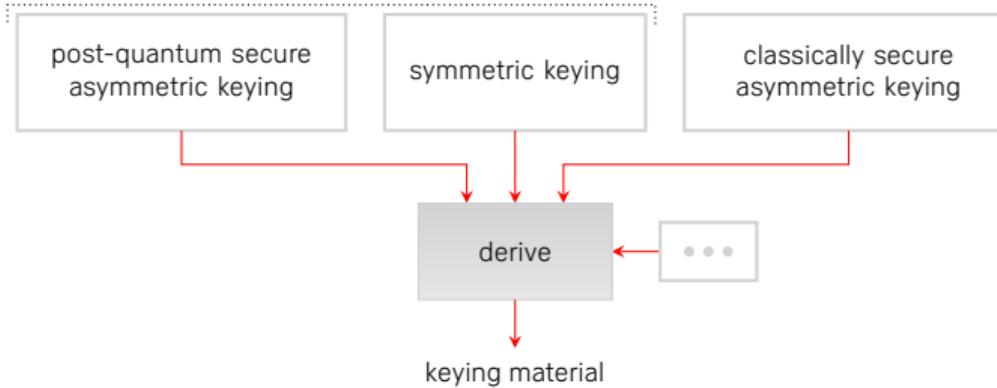


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Hybrid keying

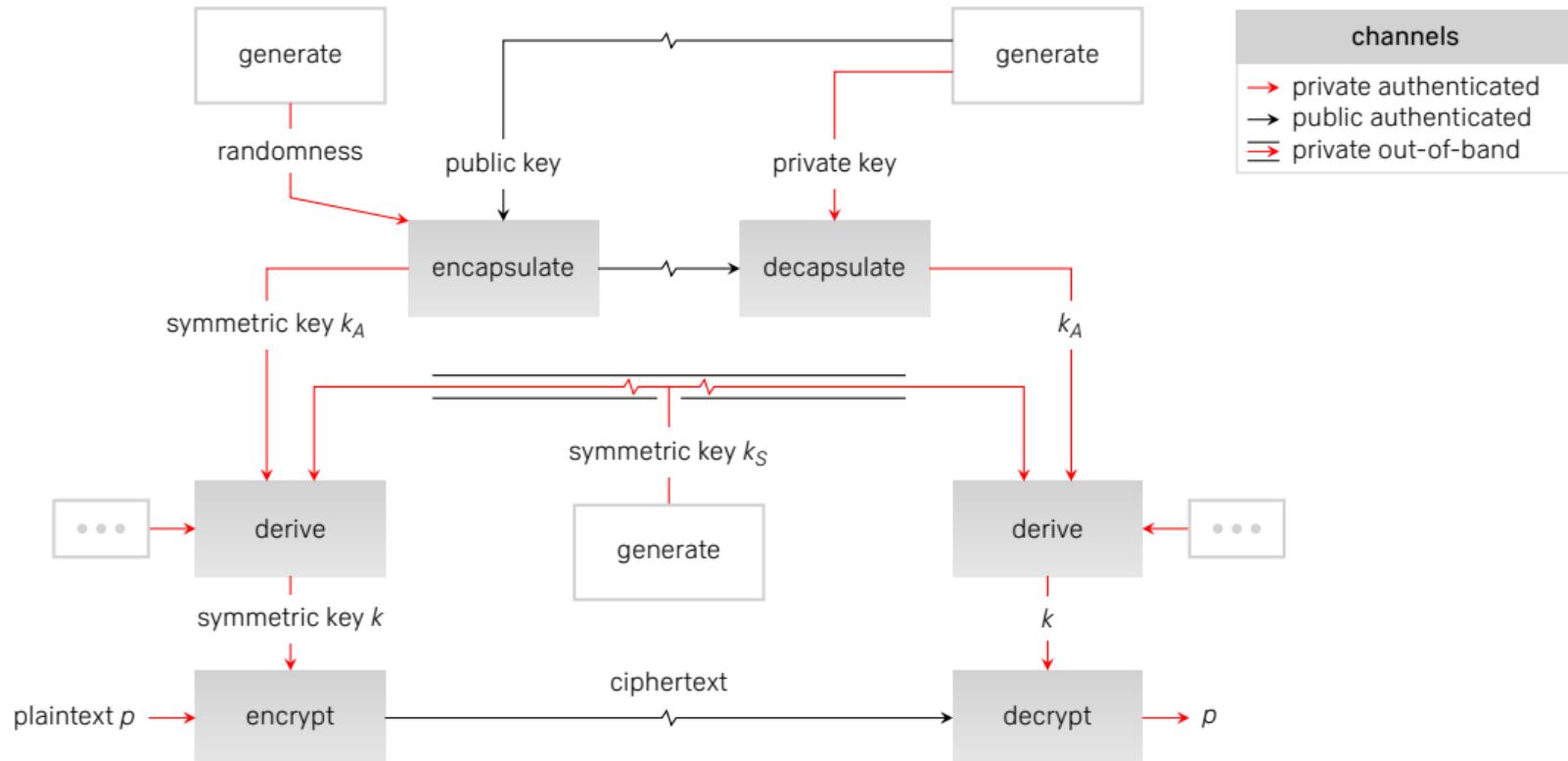
post-quantum secure — include at least one of these



3. Hybridize keying methods, e.g. via key derivation or layered encryption, with the aim of all methods having to be broken for the resulting hybrid method to be broken.

- ▶ At least one method must be post-quantum secure. Use symmetric keying as a baseline whenever feasible. Hybridize with asymmetric keying for FS.
- ▶ Keep current classically secure methods to ensure security cannot be degraded.

Hybrid symmetric and asymmetric keying



Further reading

- ▶ Be conservative. Prioritize. Use symmetric keying if feasible.
- ▶ Key encapsulation options:
 - ▶ FrodoKEM
 - ▶ ML-KEM
 - ▶ Classic McEliece
 - ▶ HQC
 - ▶ ...
- ▶ Signature options:
 - ▶ SLH-DSA
 - ▶ XMSS/LMS
 - ▶ ML-DSA
 - ▶ ...
- ▶ Hybridize all non-hash-based schemes. Avoid Level I-II for non-hash-based schemes.



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Primary quantum algorithms for cryptanalysis

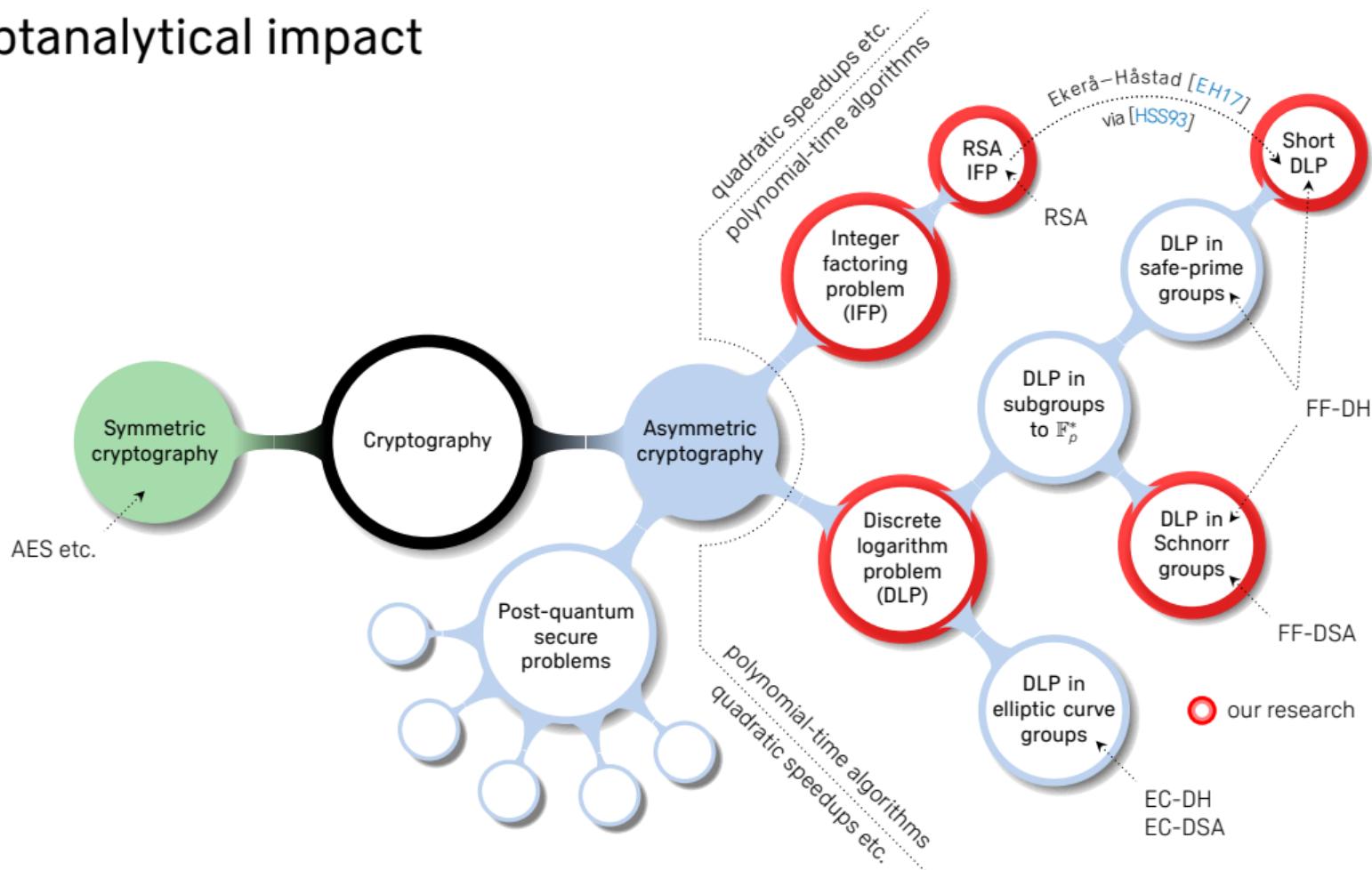
Shor's algorithms

- ▶ [Shor94] solve both the integer factoring problem (IFP), and the discrete logarithm problem (DLP) in finite cyclic groups, in polynomial time and space.
- ▶ Asymmetric cryptography based on either of these problems is vulnerable.

Grover's algorithm

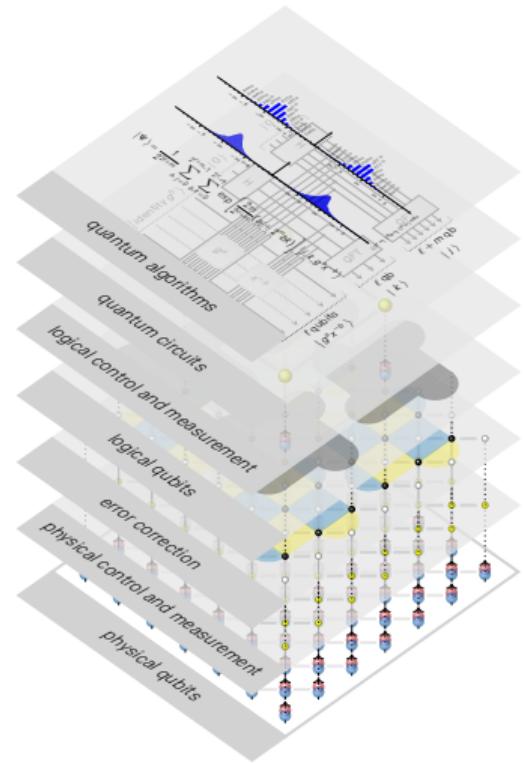
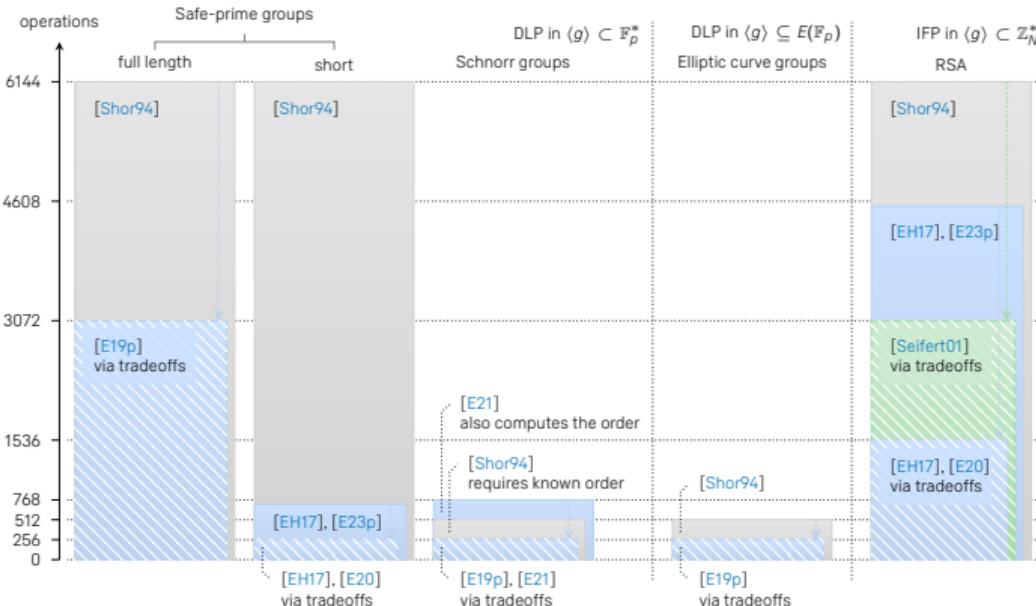
- ▶ [Grover96] provides a quadratic speedup for exhaustive search — in theory.
- ▶ In practice, due to overheads, the slow speed of quantum computers, and poor parallelization, it is not clear if [Grover96] provides a speedup. Easily mitigated.

Cryptanalytical impact



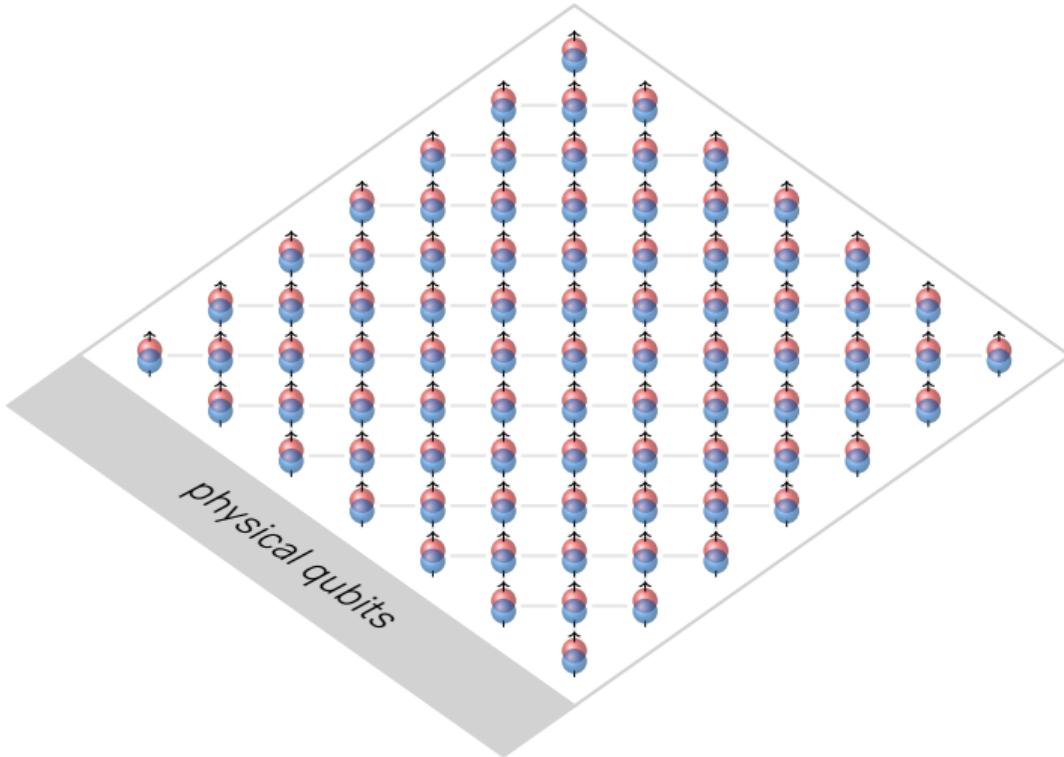
Our quantum cryptanalysis research

Fig.: Group operations per run for a 128-bit classical strength level

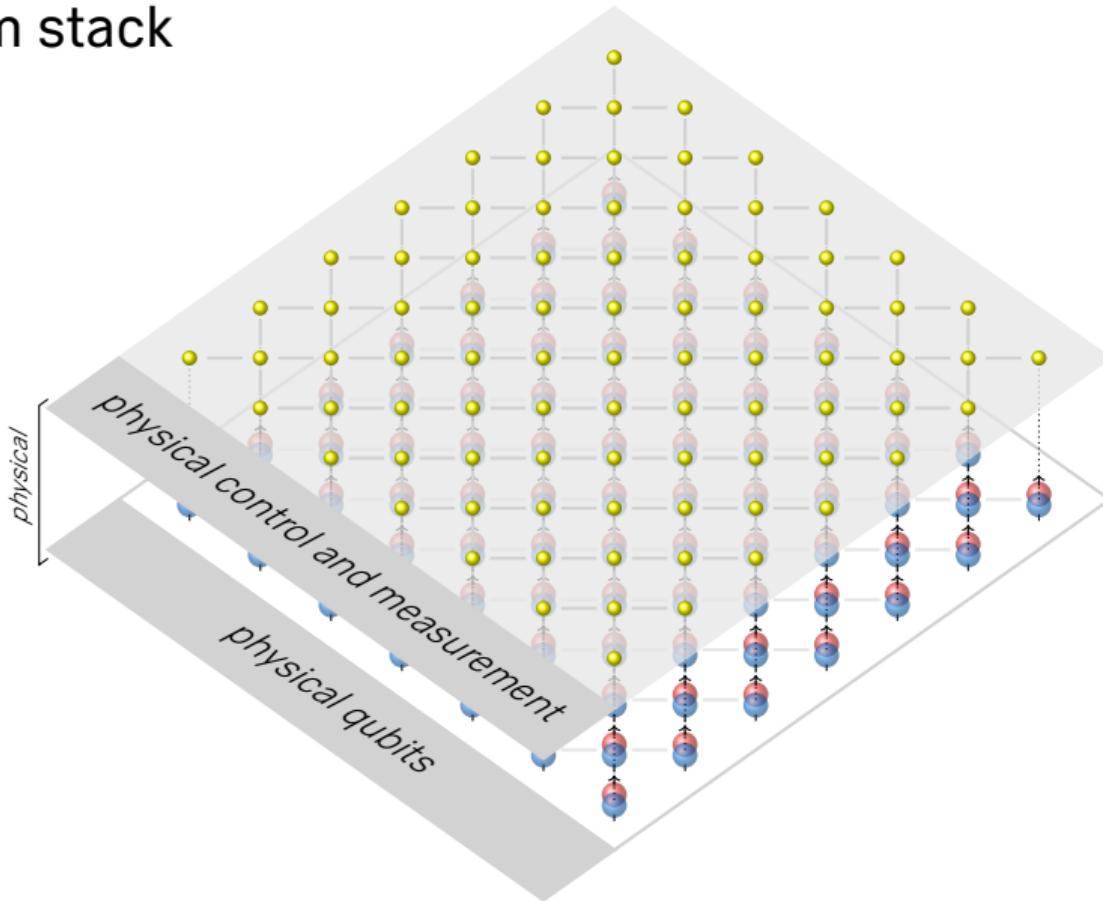


- I have developed state-of-the-art quantum algorithms for breaking widely deployed asymmetric cryptography and costed these to inform mitigation timelines.

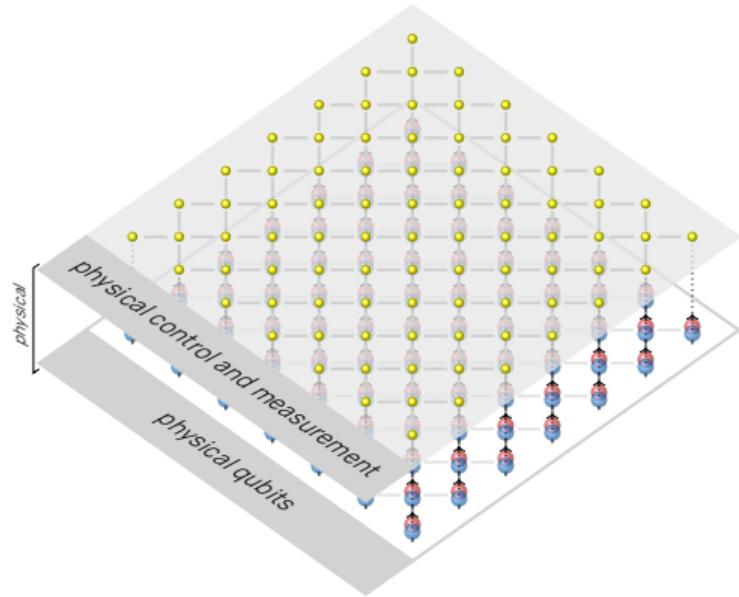
The quantum stack



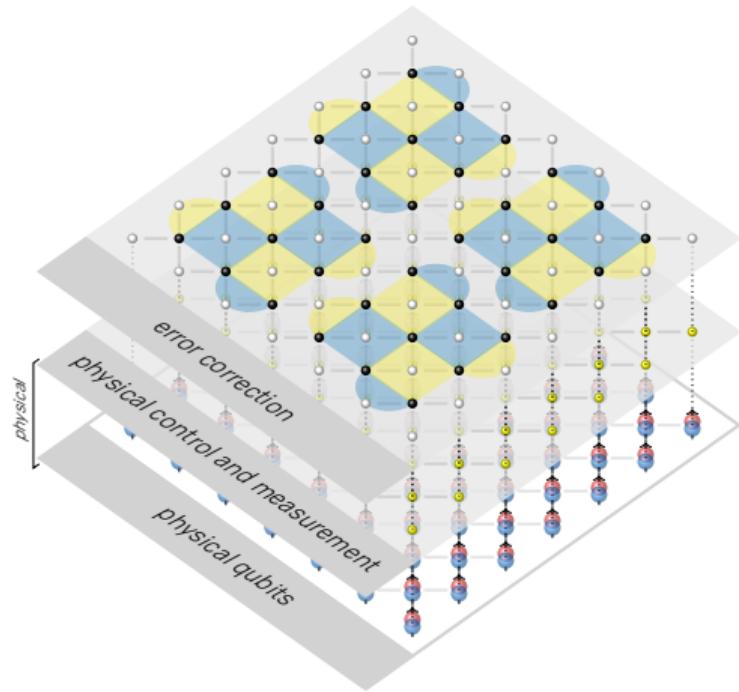
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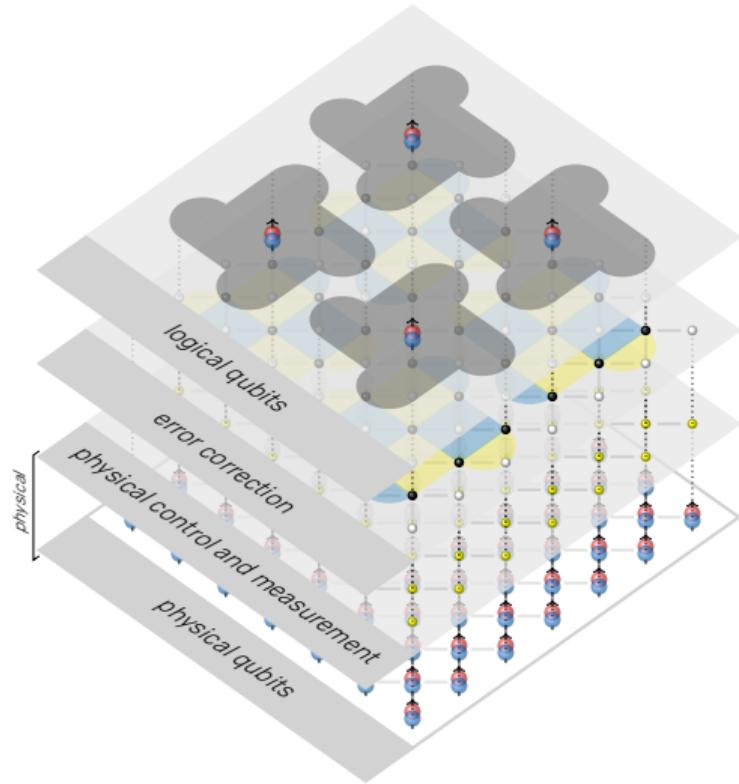
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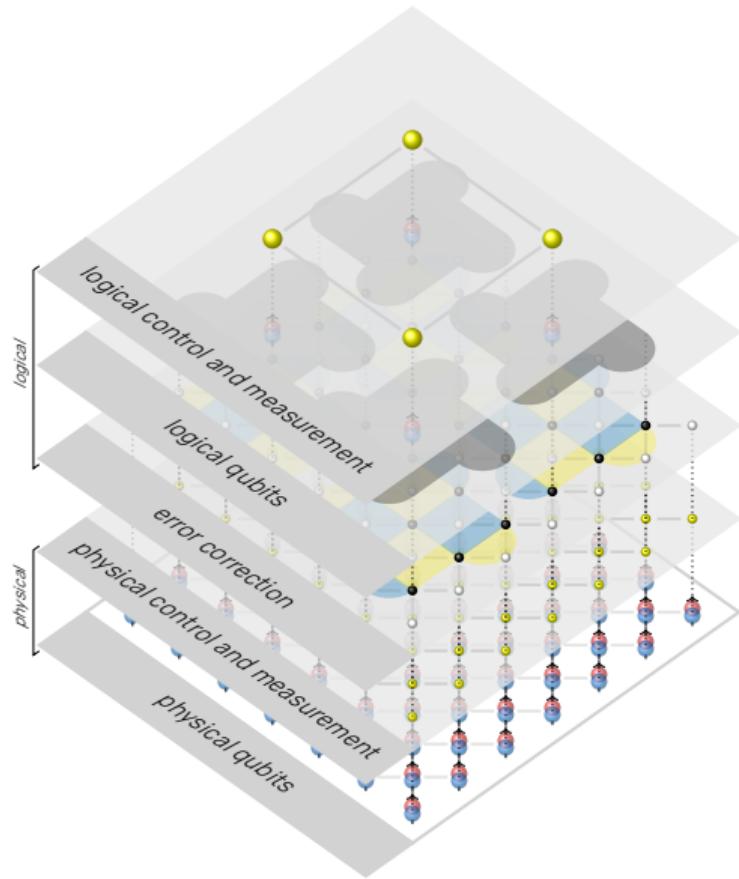
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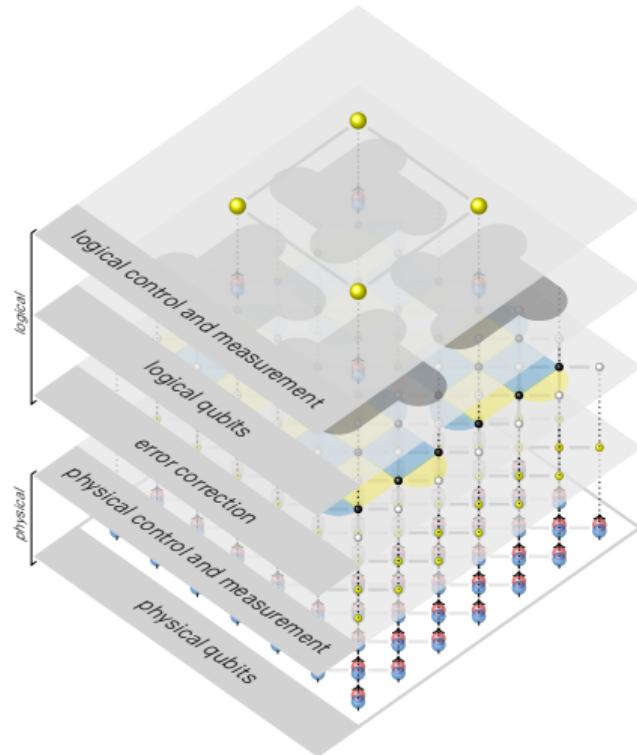
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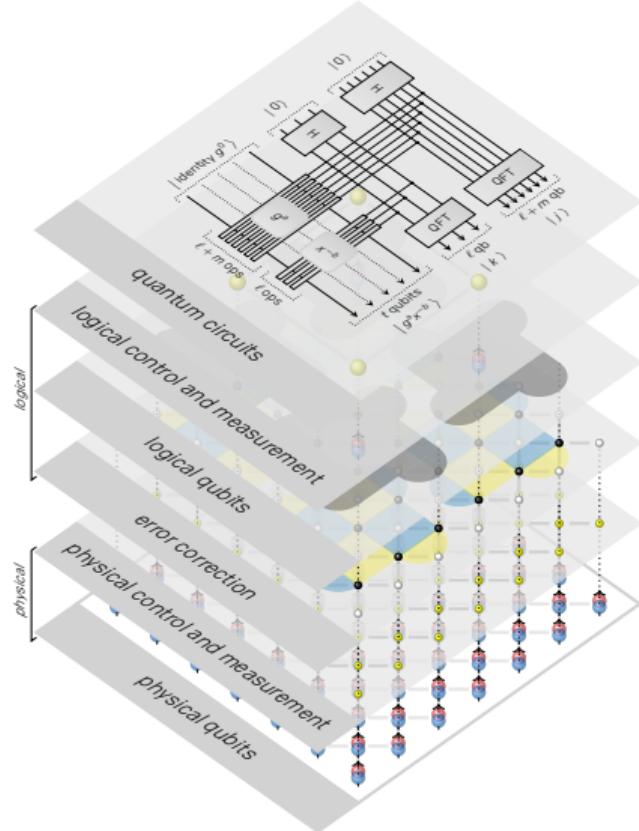
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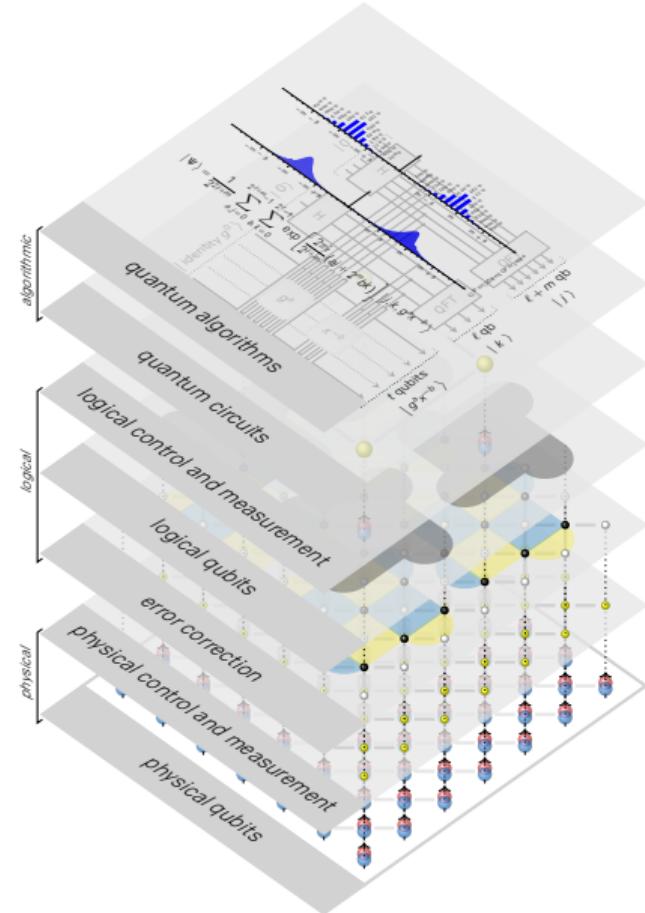
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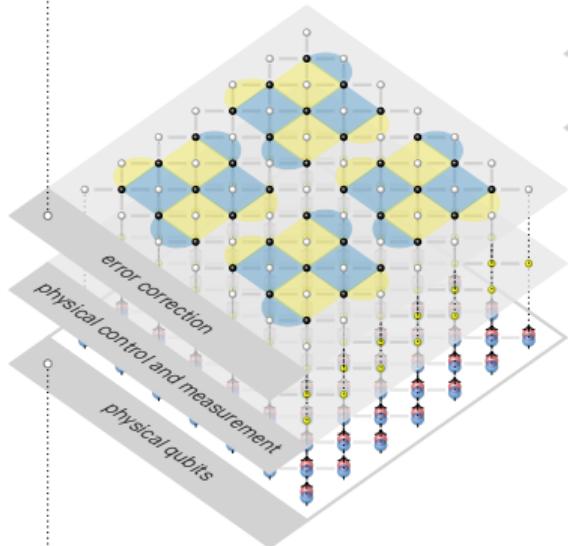


Full-stack cost estimates [GE21]

Efficient error correction



Austin Fowler
Craig Gidney

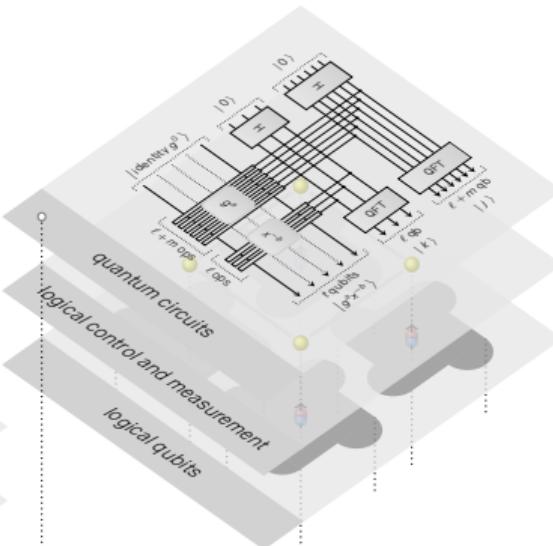


Plausible physical assumptions

Efficient approximate modular integer arithmetic



Craig Gidney



quantum algorithms



Martin Ekerå
Johan Håstad

Efficient quantum algorithms



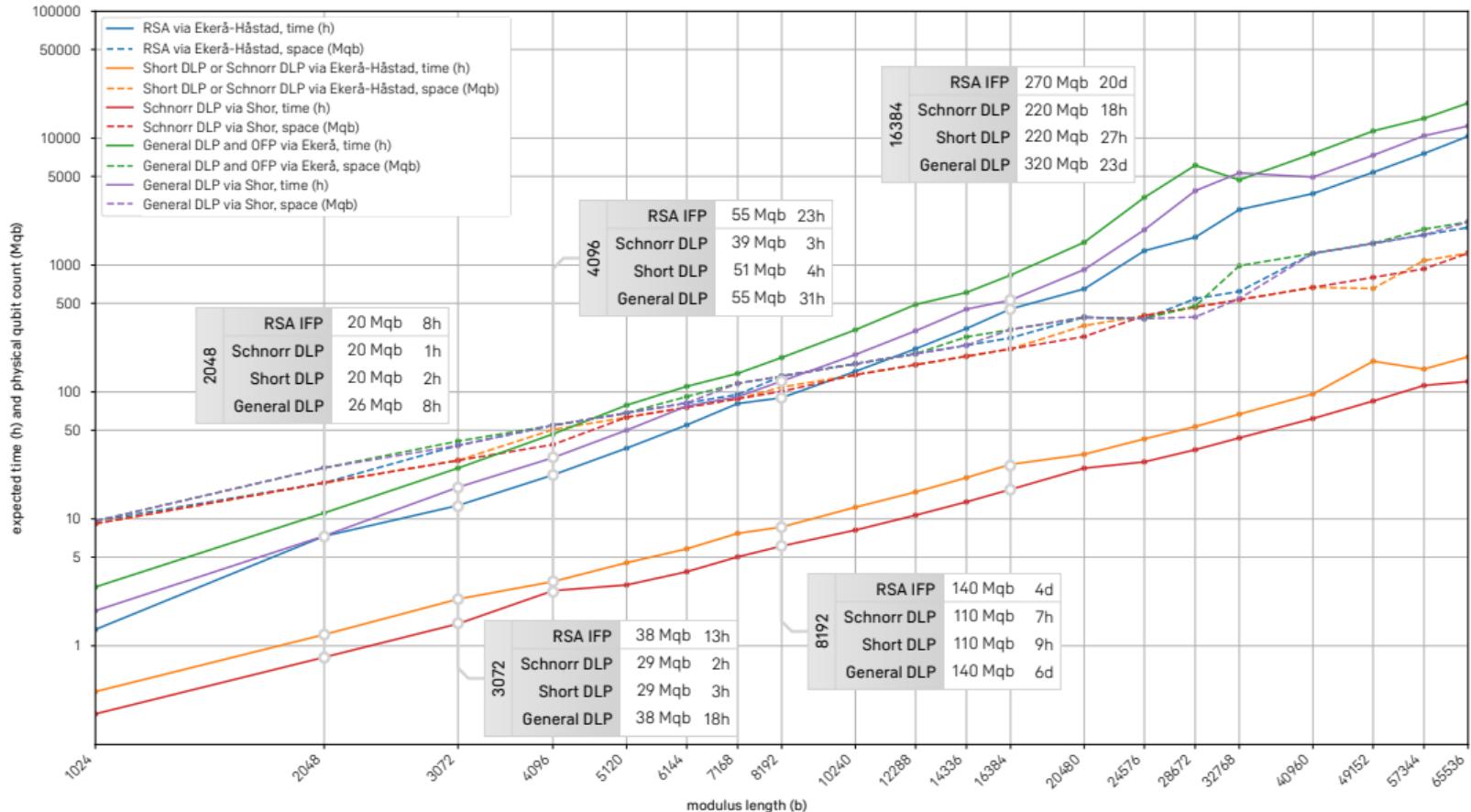
Martin Ekerå
Johan Håstad

Efficient classical post-processing and tight probability estimates



Martin Ekerå

Full-stack cost estimates [GE21]



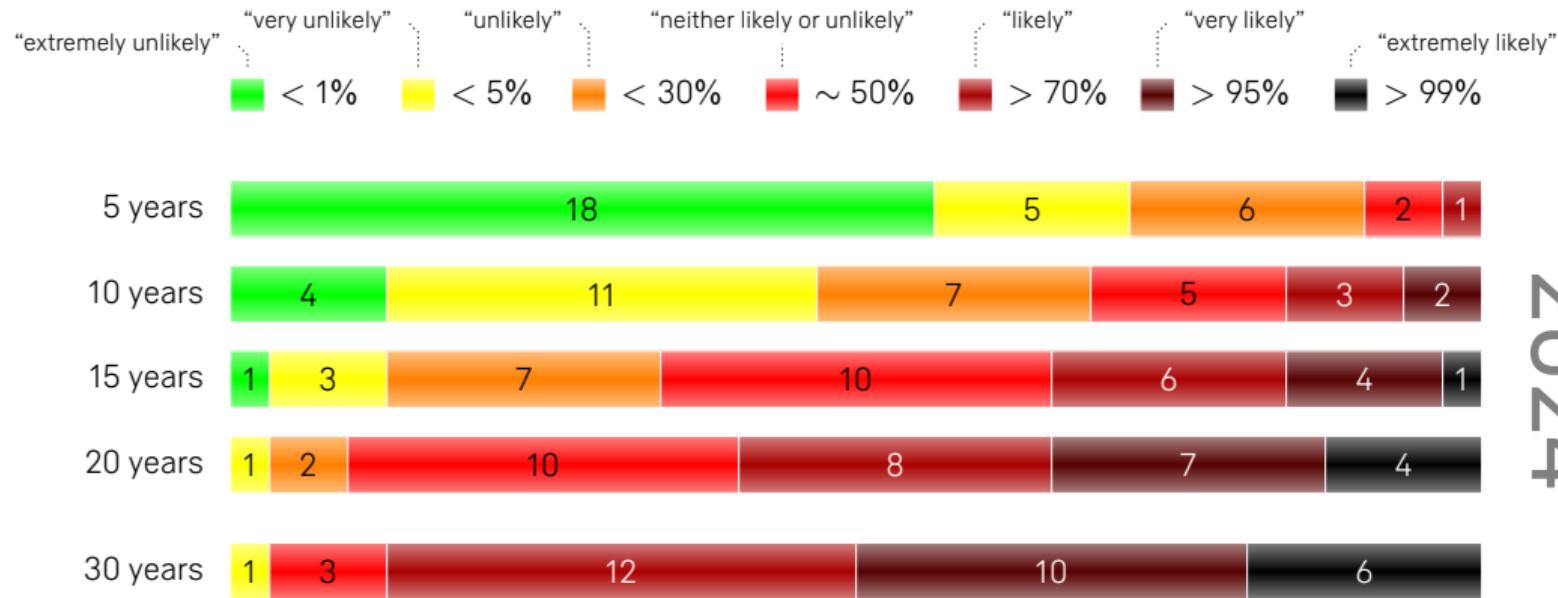
These estimates are from [GE21], see the paper and abstract for details on assumptions. Specifically, they are for factoring RSA integers, for solving the DLP in Schnorr groups, and for solving the general and short DLP in safe-prime groups, without making tradeoffs with respect to the number of runs required. The costs reported were obtained by optimizing the skewed volume, again see the paper for details. The classical strength level z is estimated using the model in FIPS 140-2 IG. For Schnorr groups, the order r is of length 22 bits. For safe-prime groups, the short exponent d is of length 22 bits.

What does this mean for the timeline?



Respondents 2019–2024: Dorit Aharonov • Alexandre Blais • Ignacio Cirac • Bill Coish • David DiVincenzo • Martin Ekerå • Artur Ekert • Daniel Gottesman • Andrea Morello • Tracy Northup • Stephanie Simmons • Peter Shor • Frank Wilhelm-Mauch • Shengyu Zhang — **Additional respondents 2024:** Sergio Boixo • Earl Campbell • Andrew Childs • Joe Fitzsimons • Jay Gambetta • Yvonne Gao • Aram Harrow • Winfried Hensinger • Elham Kashefi • Yi-Kai Liu • Klaus Mølmer • William John Munro • Nicolas Menicucci • Kae Nemoto • Francesco Petruccione • Simone Severini • Gregor Weihs • David J. Wineland

What is the likelihood of quantumly breaking RSA-2048 in 24 hours?



2024

A key question in this survey [M. Mosca and M. Piani, Quantum Threat Timeline Report] was: "Please indicate how likely you estimate it is that a quantum computer able to factorize a 2048-bit number in less than 24 hours will be built within the indicated number of years. (For reference, you might want to take into account recent estimates for resources that might be required for such a task, like the ones provided in [C. Gidney and M. Ekerå, Quantum 5, 433 (2021)].)"

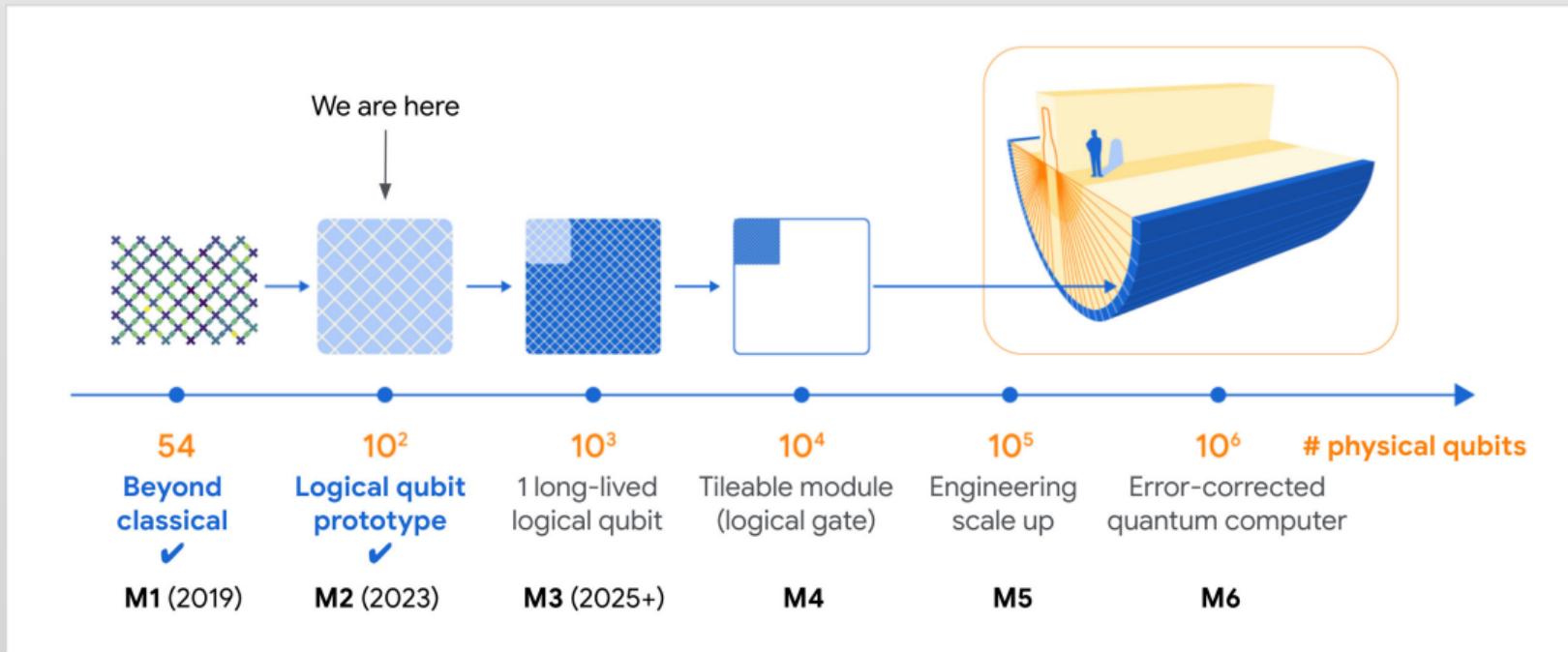
Roadmap from IBM Quantum (2024)

Development Roadmap

	2024	2025	2026	2027	2028	2029	2033+	
	Improve quantum circuit quality and speed to allow 5K gates with parametric circuits	Enhance quantum execution speed and parallelization with partitioning and quantum modularity	Improve quantum circuit quality to allow 7.5K gates	Improve quantum circuit quality to allow 10K gates	Improve quantum circuit quality to allow 15K gates	Improve quantum circuit quality to allow 100M gates	Beyond 2033, quantum-centric supercomputers will include 1000's of logical qubits unlocking the full power of quantum computing	
Data scientists	Platform	Qiskit Code Assistant	Qiskit Functions Service	Mapping collections	Specific libraries	General purpose QC libraries		
Quantum physicists	Qiskit Runtime Service	Heron (5K) Error mitigation 5k gates 133 qubits Classical modular $133 \times 3 = 399$ qubits	Flamingo (5K) Error mitigation 5k gates 156 qubits	Flamingo (7.5K) Error mitigation 7.5k gates 156 qubits	Flamingo (10K) Error mitigation 10k gates 156 qubits	Flamingo (15K) Error mitigation 15k gates 156 qubits	Starling (100M) Error correction 100M gates 200 qubits	Blue Jay (1B) Error correction 1B gates 2000 qubits

Cropped roadmap adapted from the roadmap in the "IBM Quantum 2024 State of the Union" by J. Gambetta et al.

Roadmap from Google Quantum AI (2022)



Roadmap presented by H. Neven in his talk “Google Quantum AI update” at Quantum Summer Symposium 2022. The high-resolution image was retrieved from the “Our quantum error correction milestone” article on the Google Quantum AI website. In a later revision, the 2025+ target for M3 was removed, and logical qubit error rates specified: 10^{-2} for M2, 10^{-6} for M3–M5, and 10^{-13} for M6. The original roadmap specified a 2029 target for M6.

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- ▶ It is conceivable that such computers may be built sometime after the year 2030.

Mitigation advice for vulnerable asymmetric cryptography

- ▶ Prioritize taking mitigating actions with respect to providing confidentiality.
- ▶ If feasible, use symmetric keying as a baseline, in combination with asymmetric keying. Otherwise, use post-quantum secure asymmetric keying as a baseline.
- ▶ Be mindful of the timeframes. Early mitigation is an affordable insurance.

