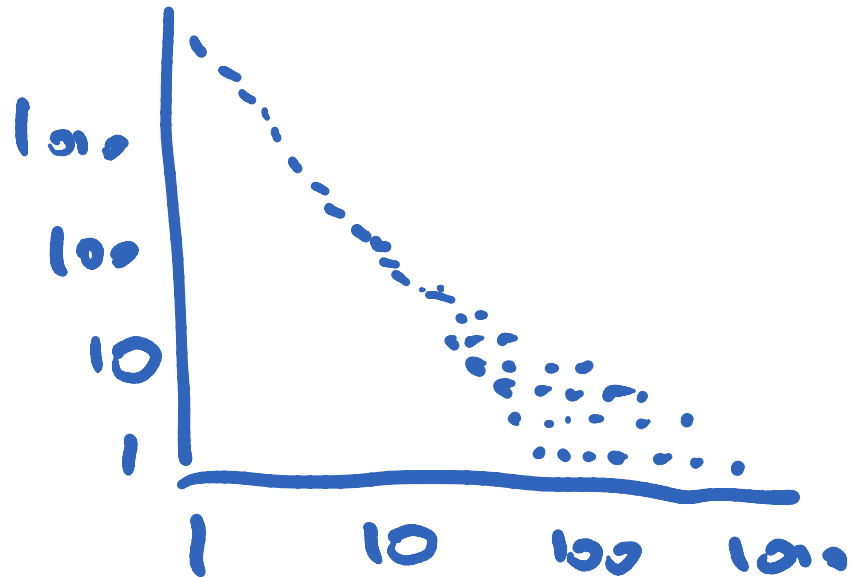
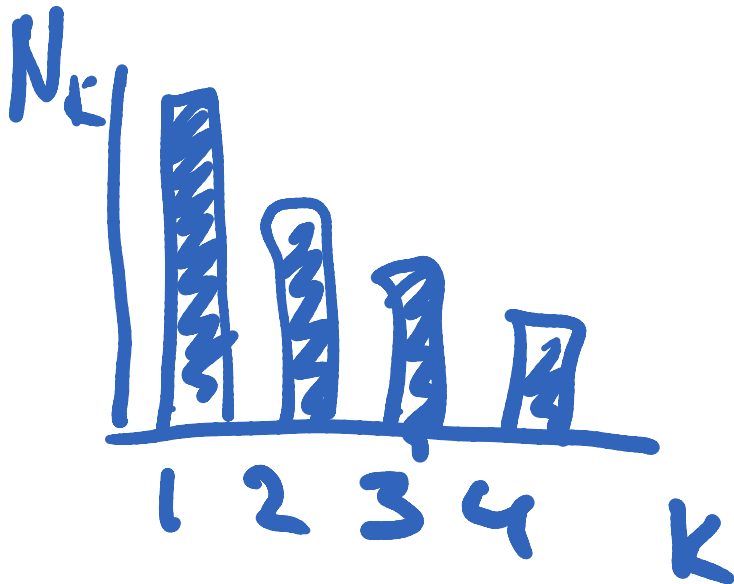


# Recap

- Type of graphs
- Connectivity/Giant component
- Diameter
- Clustering coefficient
- Betweenness Centrality
- **Degree distributions**

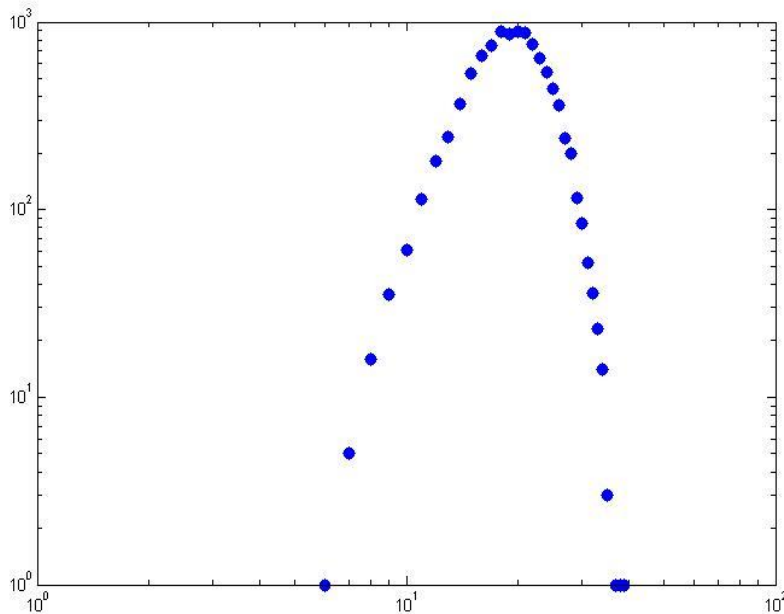
# Degree Distribution

- $N_k$  is the number of nodes with degree  $k$
- $P(k)$  is the probability that a randomly chosen node has degree  $k$ .
  - $P(k) = N_k / N$
  - Often power-law distributions (linear in loglog scale)

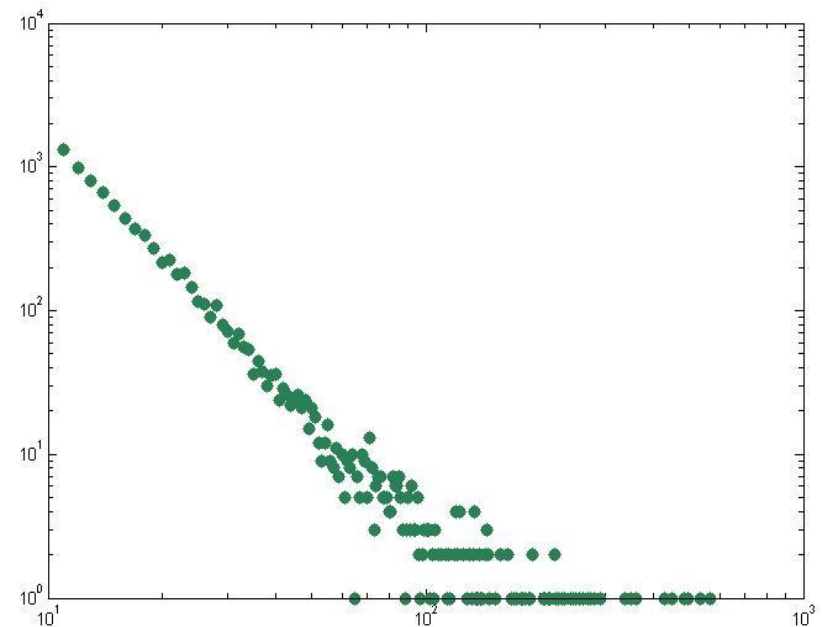


# More degree distributions

- Normal vs. power-law distributions
- $N=10k$  nodes, avg  $d=20$ ,



Random Graph

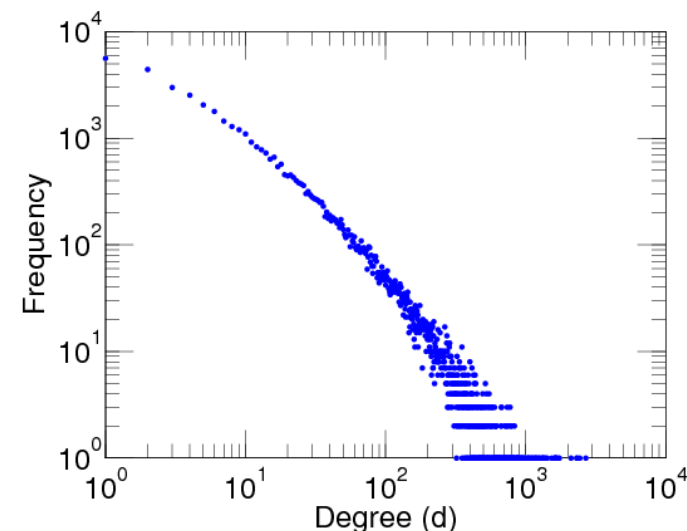
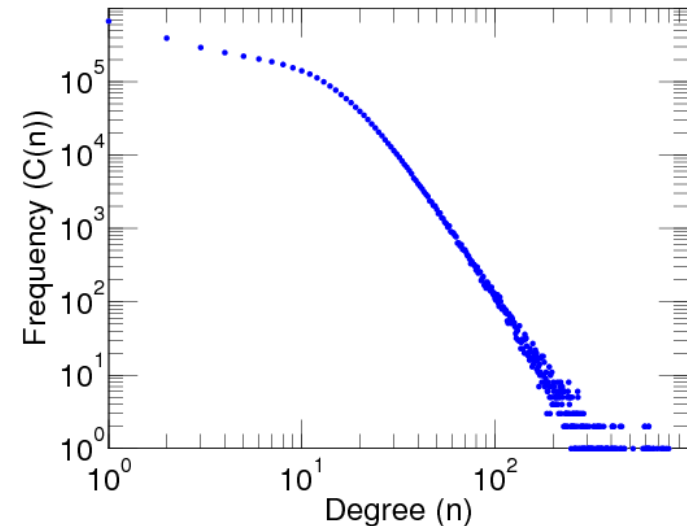


Preferential attachment (power law)

# Some real world examples

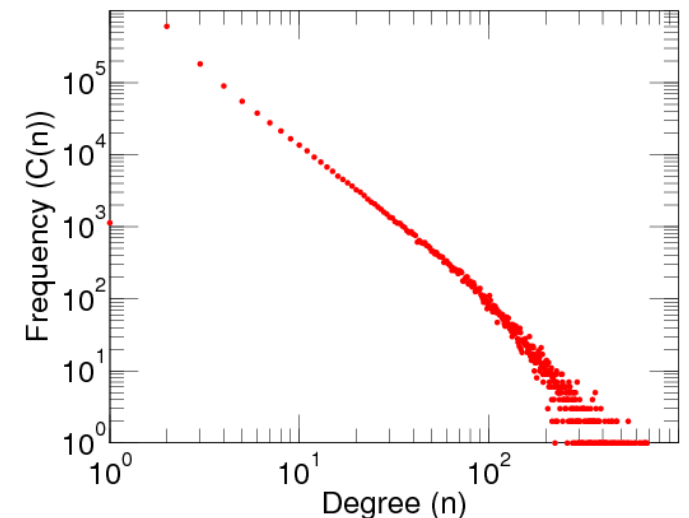
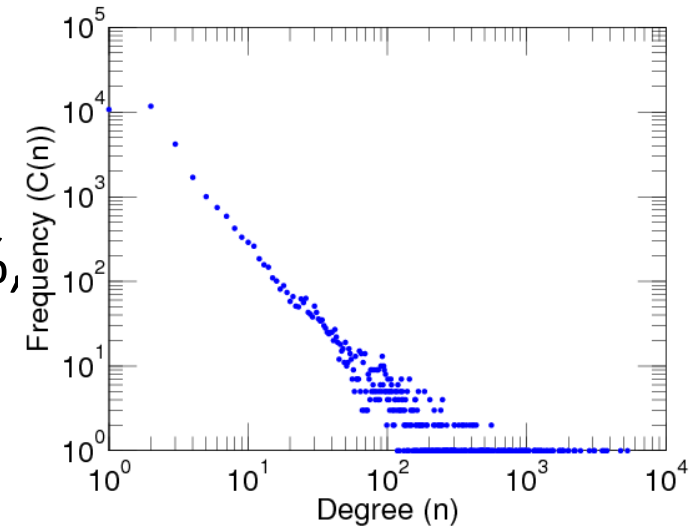
(from <http://konect.uni-koblenz.de>)

- US patents
  - Patent-patent citation
  - $N=3774768$
  - $E=16522438$
  - Effective diameter 9,79
- Facebook (user-user wall posts)
  - Directed
  - $N=63891$ ,  $E=876993$
  - $CC=19,1\%$  Effective diameter=7,25



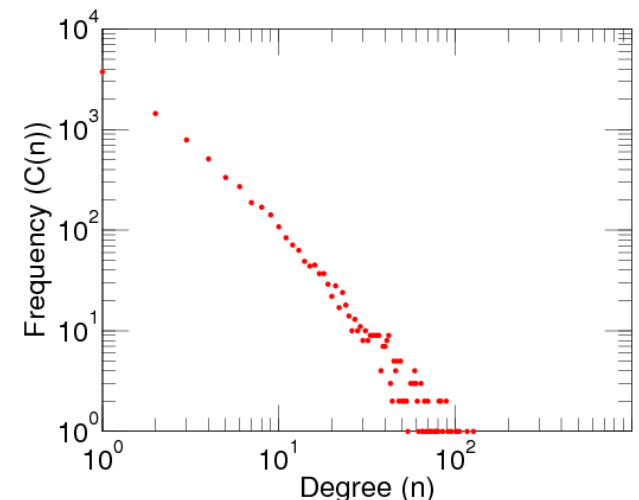
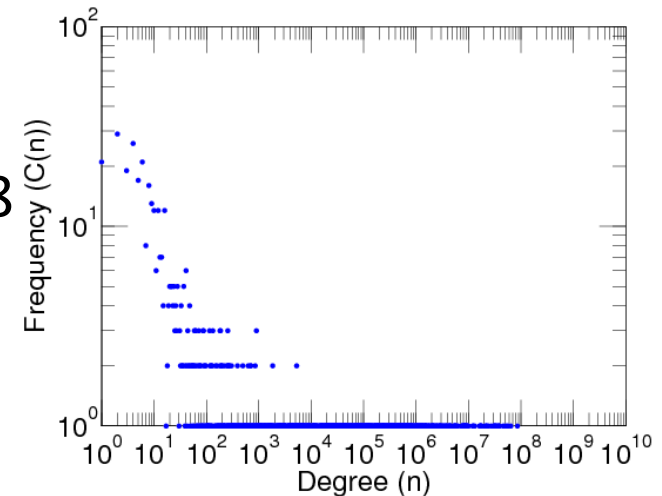
# Some real world examples (cont.)

- Internet Topology
  - AS-AS connection
  - $N=34761$ ,  $E=171403$ ,  $CC=4,851\%$ ,
  - Mean shortest path 3,77
- DBLP
  - Author-publication authorship
  - Bipartate
  - $N=4337293$
  - $E=6651968$
  - Effective diameter 15,3



# Some real world examples (cont.)

- USA airports (airport-airport flights)
  - $N=1574$ ,  $E=28236$ ,  $CC=38,4\%$ ,  $\dim=8$
- Sexual escorts (Buyer–escort contact)
  - bipartate graph
  - $N=16730$ ,  $E=50632$ ,  $\dim=6,05$
- **What do we see?**
  - *Sparse* networks, with *low diameter*, *hubs* and *non-random clustering*.



# How do we Approach Networks

- **Observations**

- Structure, properties, patterns, evolution of the networks
- Our empirical observations often find:
  - Sparse networks
  - Small diameter
  - Large clustering coefficients
  - Power law degree distributions
  - One giant component
- How to explain them?

- **Models**

- How do we model edge attachments, epidemics, communities etc?
- We will start from “simple” to more “sophisticated”.

- **Algorithms (applications)**

- How to route and search, how to find clusters, how to do efficient information collection/dissemination, gossip learning etc.

# Models of Graphs

- **$G(n,m)$  model**
  - Start with  $n$  isolated vertices
  - Place  $m$  edges among them at random.
  - $G(n,m)$  defines a family of graphs (not a particular graph)
- **$G(n, p)$  model (Erdos-Renyi random graph)**
  - Start with  $n$  isolated vertices
  - We place an edge between each distinct vertex pair with probability  $p$ .
  - $n$  and  $p$  do not uniquely determine the graph!
- **Q1: What's a degree of the above networks?**
- **Q2: Which family is bigger?**



# Erdos-Renyi random graph


- What can you say about the graph when we move  $p$  from 0 to 1?
  - Diameter?
  - How big is a giant component when  $p=0$  and when  $p=1$ ?
  - How does the giant component grow inbetween those  $p$  values?
    - Network undergoes “*phase transition*”

# Erdős-Renyi random graph (cont.)

as  $N$  becomes large:

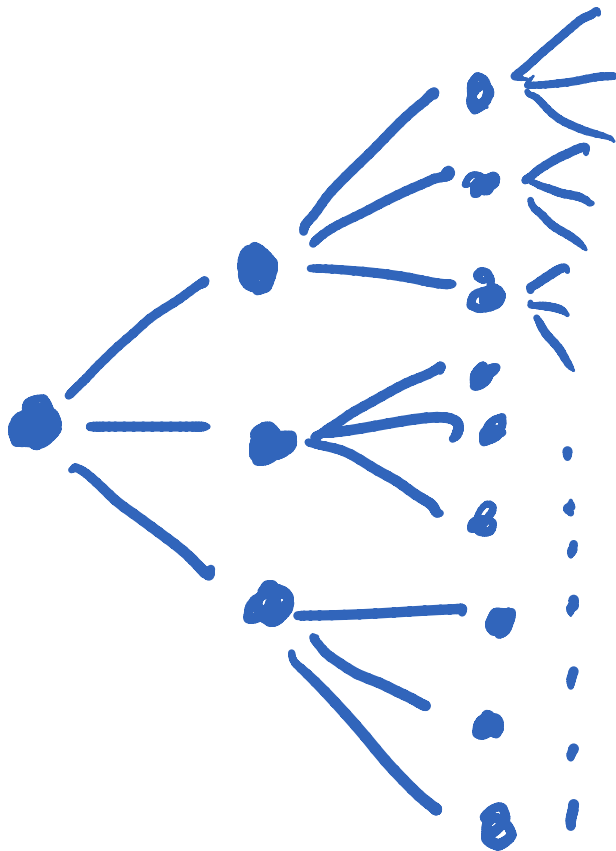
If  $p < 1/N$ , probability of a giant component goes to 0

If  $p > 1/N$ , probability of a giant component goes to 1, and all other components will have size at most  $\log(N)$

- at  $p \sim 1/N$ , average degree is  $\sim 1$  (sparse graph)
  - For regular Random graphs when  $d$  is  $\geq 3$  the graph is connected a.a.s.
- Any monotone property exhibits **Threshold phenomena** in Erdos-Renyi with respect to  $p$ .
  - E.g., network has a cycle of at least  $K$  vertices.
- Diameter: approx  $\log(N)/\log(d)$
- Clustering coefficient: 

<http://ccl.northwestern.edu/netlogo/models/GiantComponent>

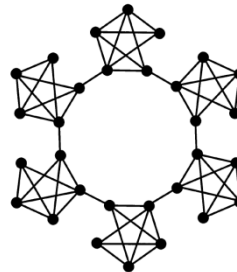
# Intuition on Diameter Calculation



- Take a tree with degree  $d$
- Nodes reached in
  - 1<sup>st</sup> step  $1+d$
  - 2<sup>nd</sup> step  $1+d+d^2$
  - 3<sup>rd</sup> step  $1+d+d^2+d^3$
  - $k^{\text{th}}$  step  $1+d+d^2+d^3+\dots+d^k \sim d^k$
- When do we reach  $N$  nodes?
  - $N=d^k$
  - $k=\log_d N$
- Intuition for Random graphs
  - Most of the degrees in random graph are between  $3d$  and  $d/3$  (chernoff bounds), so  $\log(3d)$  almost equal to  $\log(d/3)$  and equal to  $\log(d)$
  - Very few neighbors are neighbors themselves, so we do not hit many “covered” nodes.
    - Most of the nodes are hit on the last step.

# Random graphs and real world

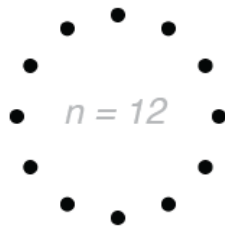
- **Erdos-Renyi can explain small diameter, connectivity under low density**
- Does not explain high clustering
  - No bias picking friends of a friend
- Does not explain power-law degree distributions
- We need a model that can explain biases towards selecting friend-of-a-friend
  - Any ideas?
    - E.g., caveman world



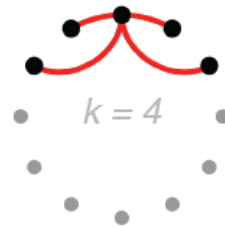
# Watts-Strogatz Model

- Regular graph with degree  $k$  connected to nearest neighbors

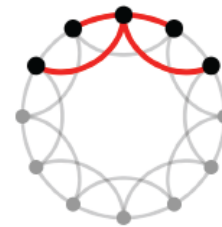
We start with a  
ring of  $n$  vertices



where each vertex  
is connected to its  
 $k$  nearest neighbors



like so.



- Can be also a grid, torus, or any other “geographical” structure which has high clusterisation and high path length
- With probability  $p$  rewire each edge in the network to a random node.***
  - Q: What happens when  $p=1$ ?***