

7.3 CODING

One of the most important ideas for studying complicated dynamics sounds strange at first. It involves throwing away some information by tracking orbits only approximately. The idea is to divide the phase space into finitely many pieces and to follow an orbit only to the extent of specifying which piece it is in at a given time. This is a bit like the itinerary of the harried tourist in Europe, who decides that it is Tuesday, so the place must be Belgium. A more technological analogy would be to look at the records of a cell phone addict and track which local transmitters were used at various times.

In these analogies one genuinely loses information, because the sequence of European countries or of local cellular stations does not pinpoint the traveller at any given moment. However, orbits in a dynamical system do not move around at whim, and the deterministic nature of the dynamics has the effect that a complete *itinerary* of this sort may (and often does) give all the information about a point. This is the process of *coding* of a dynamical system.

7.3.1 Linear Expanding Maps

The linear expanding maps

$$E_m: S^1 \rightarrow S^1, E_m(x) = mx \pmod{1}$$

from Section 7.1.1 are chaotic (Corollary 7.2.8), that is, they exhibit coexistence of dense orbits (Proposition 7.2.7) with a countable dense set of periodic orbits (Proposition 7.1.3). Thus the orbit structure is both complicated and highly nonuniform. Now we look at these maps from a different point of view, which in turn gives a deeper appreciation of just how complicated their orbit structure really is. To simplify notations, assume as before that $m = 2$.

Consider the binary intervals

$$\Delta_n^k := \left[\frac{k}{2^n}, \frac{k+1}{2^n} \right] \quad \text{for } n = 1, \dots \quad \text{and } k = 0, 1, \dots, 2^n - 1.$$

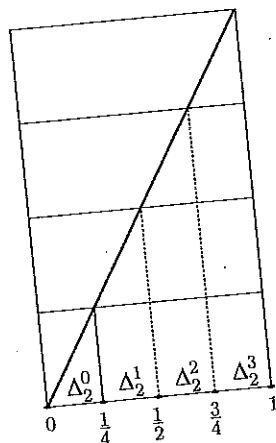


Figure 7.3.1. Linear coding.

Figure 7.3.1 illustrates this for $n = 2$. Let $x = 0.x_1x_2 \dots$ be the binary representation of $x \in [0, 1]$. Then $2x = x_1.x_2x_3 \dots \pmod{1}$. Thus

$$(7.3.1) \quad E_2(x) = 0.x_2x_3 \dots \pmod{1}.$$

This is the first and easiest example of coding, which we will discuss in greater detail shortly.

7.3.2 Implications of Coding

We briefly derive a few new facts about linear expanding maps that are best seen via this coding.

1. Proof of Transitivity via Coding. First, we use this representation to give another proof of topological transitivity by describing explicitly the binary representation of a number whose orbit under the iterates of E_2 is dense. Consider an integer k , $0 \leq k \leq 2^n - 1$. Let $k_0 \dots k_{n-1}$ be the binary representation of k , maybe with several zeroes at the beginning. Then $x \in \Delta_n^k$ if and only if $x_i = k_i$ for $i = 0, \dots, n-1$. Therefore we write $\Delta_{k_0 \dots k_{n-1}} := \Delta_n^k$ from now on. Now put the binary representations of all numbers from 0 to $2^n - 1$ (with zeroes in front if necessary) one after another and form a finite sequence, which we denote by ω_n , that is, ω_n is obtained by concatenating all 2^n binary sequences of length n . Having done this for every $n \in \mathbb{N}$, put the sequences ω_n , $n = 1, 2, \dots$ in that order, call the resulting infinite sequence ω , and consider the number x with the binary representation $0.\omega$. Since by construction moving ω to the left and cutting off the first digits produces at various moments binary representations of any n -digit number, this means that the orbit of the point x under the iterates of the map E_2 intersects every interval $\Delta_{k_0 \dots k_{n-1}}$ and hence is dense.

This construction extends to any $m \geq 2$. To construct a dense orbit for E_m with $m \leq -2$, we notice that $E_m^2 = E_{m^2}$. Obviously the orbit of any point under the iterates of a square of a map is a subset of the orbit under the iterates of the map itself; thus if the former is dense, so is latter. So we apply our construction to the map E_{m^2} and obtain a point with dense orbit under E_m .

2. Exotic Asymptotics. Next we use this approach to show that besides periodic and dense orbits there are other types of asymptotic behavior for orbits of expanding maps. One can construct such orbits for E_2 , but the simplest and most elegant example appears for the map E_3 .

Proposition 7.3.1 *There exists a point $x \in S^1$ such that the closure of its orbit with respect to the map E_3 in additive notation coincides with the standard middle-third Cantor set K . In particular, K is E_3 -invariant and contains a dense orbit.*

Proof The middle-third Cantor set K can be described as the set of all points on the unit interval that have a representation in base 3 with only 0's and 2's as digits (see Section 2.7.1). Similarly to (7.3.1), the map E_3 acts as the shift of digits to the left in the base 3 representation. This implies that K is E_3 -invariant. It remains to show that E_3 has a dense orbit in K .

Every point in K has a unique representation in base 3 without 1's. Let $x \in K$ and

$$(7.3.2) \quad 0.x_1x_2x_3 \dots$$

be such a representation. Let $h(x)$ be the number whose representation in base 2 is

$$0.\frac{x_1}{2} \frac{x_2}{2} \frac{x_3}{2} \dots,$$

that is, it is obtained from (7.3.2) by replacing 2's by 1's. Thus we have constructed a map $h: K \rightarrow [0, 1]$ that is continuous, nondecreasing [that is, $x > y$ implies $h(x) \geq h(y)$], and one-to-one, except for the fact that binary rationals have two preimages each (compare Section 2.7.1 and Section 4.4.1). Furthermore, $h \circ E_3 = E_2 \circ h$. Let $D \subset [0, 1]$ be a dense set of points that does not contain binary rationals. Then $h^{-1}(D)$ is dense in K because, if Δ is an open interval such that $\Delta \cap K \neq \emptyset$, then $h(\Delta)$ is a nonempty interval open, closed, or semiclosed and hence contains points of D . Now take any $x \in [0, 1]$ whose E_2 -orbit is dense; the E_3 -orbit of $h^{-1}(x) \in K$ is dense in K . \square

3. Nonrecurrent Points. Another interesting example is the construction of a nonrecurrent point, that is, such a point x that for some neighborhood U of x all iterates of x avoid U (see Definition 6.1.8). In fact, there is a dense set of nonrecurrent points for the map E_2 .

Pick any fixed sequence $(\omega_0, \dots, \omega_{n-1})$ of 0's and 1's and add a tail of 0's if $\omega_{n-1} = 1$, or of 1's if $\omega_{n-1} = 0$. Call the resulting infinite sequence ω . As before, let x be the number with binary representation $0.\omega$. Thus, x lies in a prescribed interval $\Delta_{\omega_0 \dots \omega_{n-1}}$ and by construction $x \neq 0$. On the other hand, $E_2^n x = 0$ and hence $E_2^m x = 0$ for all $m \geq n$, so x is a nonrecurrent point.

Thus, we have found that E_m is chaotic and topologically mixing, that its periodic and nonrecurrent orbits are dense, and that E_3 has orbits whose closure is a Cantor set.

7.3.3 A Two-Dimensional Cantor Set

We now describe a map in the plane that naturally gives rise to a two-dimensional Cantor set (previously encountered in Problem 2.7.5) on which ternary expansion of the coordinates provides all information about the dynamics. This *horseshoe map* plays a central role in our further development.

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Consider a map defined on the unit square $[0, 1] \times [0, 1]$ by the following construction: First apply the linear transformation $(x, y) \mapsto (3x, y/3)$ to get a horizontal strip whose left and right thirds will be rigid in the next transformation. Holding the left third fixed, bend and stretch the middle third such that the right third falls rigidly on the top third of the original unit square. This results in a "G"-shape. For points that are in and return to the unit square, this map is given analytically by

$$(x, y) \mapsto \begin{cases} (3x, y/3) & \text{if } x \leq 1/3 \\ (3x - 2, (y + 2)/3) & \text{if } x \geq 2/3. \end{cases}$$

The inverse can be written as

$$(x, y) \mapsto \begin{cases} (x/3, 3y) & \text{if } y \leq 1/3 \\ ((x + 2)/3, 3y - 2) & \text{if } y \geq 2/3. \end{cases}$$

Geometrically, the inverse looks like an "e"-shape rotated counterclockwise by 90° .

To iterate this map one triples the x -coordinate repeatedly and always assumes that the resulting value is either at most $1/3$ or else at least $2/3$, that is, that the first ternary digit is 0 or 2, but not 1. (If the expansion is not unique, one requires such a choice to be possible.) Comparing with the construction of the ternary Cantor set in Section 2.7.1, one sees that the x -coordinate lies in the ternary Cantor set C . Looking at the inverse one sees likewise that, in order for all preimages to be defined, the y -coordinate lies in the Cantor set as well. Therefore this map is defined for all positive and negative iterates on the two-dimensional Cantor set $C \times C$. There is a straightforward way of using ternary expansion to code the dynamics. For a point (x, y) the map shifts the ternary expansion of x one step to the left, dropping the first term, and shifts the ternary expansion of y to the right. It is natural to fill in the now-ambiguous first digit of the shifted y -coordinate with the entry from the x -coordinate that was just dropped. This retains all information, and the best way of visualizing the result is to write the expansion of the y -coordinate in reverse and in front of that of the x -coordinate. This gives a bi-infinite string of 0's and 2's (remember, no 1's allowed), which is shifted by the map. Of course, one should verify that the inverse acts by shifting in the opposite direction.

7.3.4 Sequence Spaces

Now we are ready to discuss the concept of coding in general. We mean by coding a representation of points in the phase space of a discrete-time dynamical system or an invariant subset by sequences (not necessarily unique) of symbols from a certain "alphabet," in this case the symbols $0, \dots, N - 1$. So we should acquaint ourselves with these spaces.

Denote by $\Omega_N^{\mathbb{R}}$ the space of sequences $\omega = (\omega_i)_{i=0}^{\infty}$ whose entries are integers between 0 and $N - 1$. Define a metric by

$$(7.3.3) \quad d_\lambda(\omega, \omega') := \sum_{i=0}^{\infty} \frac{\delta(\omega_i, \omega'_i)}{\lambda^i},$$

where $\delta(k, l) = 1$ if $k \neq l$, $\delta(k, k) = 0$, and $\lambda > 2$. The same definition can be made for two-sided sequences by summing over $i \in \mathbb{Z}$:

$$(7.3.4) \quad d_\lambda(\omega, \omega') := \sum_{i \in \mathbb{Z}} \frac{\delta(\omega_i, \omega'_i)}{\lambda^{|i|}},$$

for some $\lambda > 3$. This means that two sequences are close if they agree on a long stretch of entries around the origin.

Consider the symmetric cylinder defined by

$$C_{\alpha_1 \dots \alpha_{n-1}} := \{\omega \in \Omega_N \mid \omega_i = \alpha_i \text{ for } |i| < n\}.$$

Fix a sequence $\alpha \in C_{\alpha_1 \dots \alpha_{n-1}}$. If $\omega \in C_{\alpha_1 \dots \alpha_{n-1}}$, then

$$d_\lambda(\alpha, \omega) = \sum_{i \in \mathbb{Z}} \frac{\delta(\alpha_i, \omega_i)}{\lambda^{|i|}} = \sum_{|i| \geq n} \frac{\delta(\alpha_i, \omega_i)}{\lambda^{|i|}} \leq \sum_{|i| \geq n} \frac{1}{\lambda^{|i|}} = \frac{1}{\lambda^{n-1}} \frac{2}{\lambda - 1} < \frac{1}{\lambda^{n-1}}.$$

Thus $C_{\alpha_1 \dots \alpha_{n-1}} \subset B_{d_\lambda}(\alpha, \lambda^{1-n})$, the λ^{1-n} -ball around α . If $\omega \notin C_{\alpha_1 \dots \alpha_{n-1}}$, then

$$d_\lambda(\alpha, \omega) = \sum_{i \in \mathbb{Z}} \frac{\delta(\alpha_i, \omega_i)}{\lambda^{|i|}} \geq \lambda^{1-n}$$

because $\omega_i \neq \alpha_i$ for some $|i| < n$. Thus $\omega \notin B_{d_\lambda}(\alpha, \lambda^{1-n})$, and the symmetric cylinder is the ball of radius λ^{1-n} around any of its points:

$$(7.3.5) \quad C_{\alpha_1 \dots \alpha_{n-1}} = B_{d_\lambda}(\alpha, \lambda^{1-n}).$$

Therefore, balls in Ω_N are described by specifying a symmetric stretch of entries around the initial one.

For one-sided sequences this discussion works along the same lines [one only needs $\lambda > 2$ in (7.3.4)] and λ^{1-n} -balls are described by specifying a string of n initial entries.

Our examples [see (7.3.1)] suggest to represent points in the phase space by sequences in such a way that the sequences representing the image of a point are obtained from those representing the point itself by the shift (translation) of the symbols. In this way the given transformation corresponds to the *shift transformation*

$$(7.3.6) \quad \begin{aligned} \sigma : \Omega_N &\rightarrow \Omega_N, & (\sigma \omega)_i &= \omega_{i+1} \\ \sigma^R : \Omega_N^R &\rightarrow \Omega_N^R, & (\sigma^R \omega)_i &= \omega_{i+1}. \end{aligned}$$

We often write σ_N for the shift σ on Ω_N and likewise σ_N^R for σ^R on Ω_N^R . For invertible discrete-time systems, any coding involves sequences of symbols extending in both directions; while for noninvertible systems, one-sided sequences do the job. Section 7.3.7 studies these shifts as dynamical systems.

Among the shift transformations that arise from coding there is also a new kind of combinatorial model for a dynamical system that is described by the possibility or impossibility of certain successions of events.

Definition 7.3.2 Let $A = (a_{ij})_{i,j=0}^{N-1}$ be an $N \times N$ matrix whose entries a_{ij} are either 0's or 1's. (We call such a matrix a 0-1 matrix.) Let

$$(7.3.7) \quad \Omega_A := \{\omega \in \Omega_N \mid a_{\omega_n \omega_{n+1}} = 1 \text{ for } n \in \mathbb{Z}\}.$$

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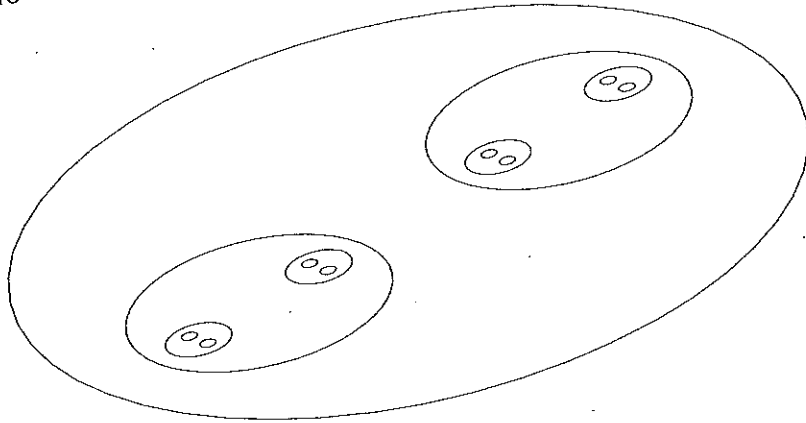


Figure 7.3.2. Obtaining a Cantor set.

The space Ω_A is closed and shift-invariant, and the restriction

$$\sigma_N|_{\Omega_A} =: \sigma_A$$

is called the *topological Markov chain* determined by A .

This is a particular case of a *subshift of finite type*.

7.3.5 Coding

Sequences representing a given point of the phase space are called the *codes* of that point. We have several examples of coding: for the map E_m on the whole circle by sequences from the *alphabet* $\{0, \dots, |m| - 1\}$; for the restriction of the map E_3 to the middle-third Cantor set K by one-sided sequences of 0's and 1's; and for the ternary horseshoe in Section 7.3.3 by bi-infinite sequences of 0's and 2's. In both cases we used one-sided sequences, all sequences appeared as codes of some points, and each code represented only one point. There was, however, an important difference: In the first case, which involved for positive m a representation in base m , a point could have either one or two codes; in the latter there was only one code.

This shows that the space of binary sequences is a Cantor set (Definition 2.7.4). In fact, this also holds for the other sequence spaces.

7.3.6 Conjugacy and Factors

This situation can be roughly described by saying that the shift (Ω_2^R, σ^R) "contains" the map f up to a continuous coordinate change. (We already encountered such a situation in Theorem 4.3.20.)

Definition 7.3.3 Suppose that $g: X \rightarrow X$ and $f: Y \rightarrow Y$ are maps of metric spaces X and Y and that there is a continuous surjective map $h: X \rightarrow Y$ such that $h \circ g = f \circ h$. Then f is said to be a *factor* of g via the *semiconjugacy* or *factor map* h . If this h is a homeomorphism, then f and g are said to be *conjugate* and h is said to be a *conjugacy*.

These notions made a brief appearance in Section 4.3.5 in connection with modeling an arbitrary homeomorphism of the circle by a rotation. The notion of conjugacy is natural and central; two conjugate maps are obtained from one another by a continuous change of coordinates. Hence all properties that are independent of such changes of coordinates are unchanged, such as the numbers of periodic orbits for each period, sensitive dependence (Exercise 7.2.5), topological transitivity, topological mixing, and hence also being chaotic. Such properties are said to be topological invariants. Later in this book we will encounter further important topological invariants such as topological entropy (Definition 8.2.1).

7.3.7 Dynamics of Shifts and Topological Markov Chains

We now study the properties of shifts and topological Markov chains introduced in (7.3.6) and Definition 7.3.2 in more detail. These are important because many interesting dynamical systems are coded by shifts or topological Markov chains. To such dynamical systems the results of this section have immediate applications.

Proposition 7.3.4 *Periodic points for the shifts σ_N and σ_N^R are dense in Ω_N and Ω_N^R , correspondingly, $P_n(\sigma_N) = P_n(\sigma_N^R) = N^n$, and both σ_N and σ_N^R are topologically mixing.*

Proof Periodic orbits for a shift are periodic sequences, that is, $(\sigma_N)^m \omega = \omega$ if and only if $\omega_{n+m} = \omega_n$ for all $n \in \mathbb{Z}$. In order to prove density of periodic points, it is enough to find a periodic point in every ball (symmetric cylinder), because every open set contains a ball. To find a periodic point in $C_{\alpha_{-m}, \dots, \alpha_m}$, take the sequence ω defined by $\omega_n = \alpha_{n'}$ for $|n'| \leq m$, $n' = n \pmod{2m+1}$. It lies in this cylinder and has period $2m+1$.

Every periodic sequence ω of period n is uniquely determined by its coordinates $\omega_0, \dots, \omega_{n-1}$. There are N^n different finite sequences $(\omega_0, \dots, \omega_{n-1})$.

To prove topological mixing, we show that $\sigma_N^n(C_{\alpha_{-m}, \dots, \alpha_m}) \cap C_{\beta_{-m}, \dots, \beta_m} \neq \emptyset$ for $n > 2m+1$, say, $n = 2m+k+1$ with $k > 0$. Consider any sequence ω such that

$$\omega_i = \alpha_i \text{ for } |i| \leq m, \quad \omega_i = \beta_{i-n} \text{ for } i = m+k+1, \dots, 3m+k+1.$$

Then $\omega \in C_{\alpha_{-m}, \dots, \alpha_m}$ and $\sigma_N^n(\omega) \in C_{\beta_{-m}, \dots, \beta_m}$.

The arguments for the one-sided shift are analogous. \square

There is a useful geometric representation of topological Markov chains. Connect i with j by an arrow if $a_{ij} = 1$ to obtain a *Markov graph* G_A with N vertices and several oriented edges. We say that a finite or infinite sequence of vertices of G_A is an *admissible path* or *admissible sequence* if any two consecutive vertices in the sequence are connected by an oriented arrow. A point of Ω_A corresponds to a doubly infinite path in G_A with marked origin; the topological Markov chain σ_A corresponds to moving the origin to the next vertex. The following simple combinatorial lemma is a key to the study of topological Markov chains:

Lemma 7.3.5 *For every $i, j \in \{0, 1, \dots, N-1\}$, the number N_{ij}^m of admissible paths of length $m+1$ that begin at x_i and end at x_j is equal to the entry a_{ij}^m of the matrix A^m .*

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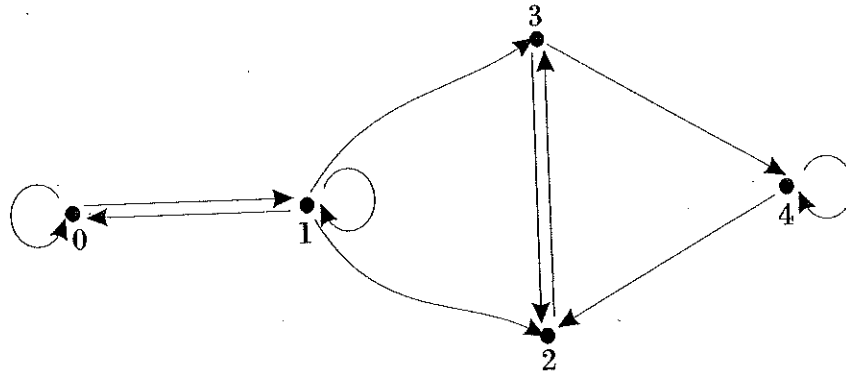


Figure 7.3.3. A Markov graph.

Proof We use induction on m . First, it follows from the definition of the graph G_A that $N_{ij}^1 = a_{ij}$. To show that

$$(7.3.8) \quad N_{ij}^{m+1} = \sum_{k=0}^{N-1} N_{ik}^m a_{kj},$$

take $k \in \{0, \dots, N-1\}$ and an admissible path of length $m+1$ connecting i and k . It can be extended to an admissible path of length $m+2$ connecting i to j (by adding j) if and only if $a_{kj} = 1$. This proves (7.3.8). Now, assuming by induction that $N_{ij}^m = a_{ij}^m$ for all ij , we obtain $N_{ij}^{m+1} = a_{ij}^{m+1}$ from (7.3.8). \square

Corollary 7.3.6 $P_n(\sigma_A) = \text{tr } A^n$.

Proof Every admissible closed path of length $m+1$ with marked origin, that is, a path that begins and ends at the same vertex of G_A , produces exactly one periodic point of σ_A of period m . \square

Because the eigenvalue of largest absolute value dominates the trace, it determines the exponential growth rate:

Proposition 7.3.7 $p(\sigma_A) = r(A)$, where $r(A)$ is the spectral radius.

Proof " \leq " is clear. To show " \geq " we need to avoid cancellations: If $\lambda_j = re^{2\pi i \varphi_j}$ ($1 \leq j \leq k$) are the eigenvalues of maximal absolute value then there is a sequence $m_n \rightarrow \infty$ such that $m_n \varphi_j \rightarrow 0 \pmod{1}$ for all j (recurrence for toral translations, Section 5.1), so $\sum \lambda_i^{m_n} \sim r^{m_n}$. \square

Example 7.3.8 The Markov graph in Figure 7.3.3 produces three fixed points, $\overline{0}$, $\overline{1}$, and $\overline{4}$. $\overline{01}$ and $\overline{23}$ give four periodic points with period 2. The period-3 orbits are generated by $\overline{011}$, $\overline{001}$, $\overline{234}$.

Topological Markov chains can be classified according to the recurrence properties of various orbits they contain. Now we concentrate on those topological Markov chains that possess the strongest recurrence properties.

Definition 7.3.9 A matrix A is said to be *positive* if all its entries are positive. A 0-1 matrix A is said to be *transitive* if A^m is positive for some $m \in \mathbb{N}$. A topological Markov chain σ_A is said to be *transitive* if A is a transitive matrix.

Lemma 7.3.10 *If A^m is positive, then so is A^n for any $n \geq m$.*

Proof If $a_{ij}^m > 0$ for all i, j , then for each j there is a k such that $a_{kj} = 1$. Otherwise, $a_{ij}^m = 0$ for every n and i . Now use induction. If $a_{ij}^n > 0$ for all i, j , then $a_{ij}^{n+1} = \sum_{k=0}^{N-1} a_{ik}^n a_{kj} > 0$ because $a_{kj} = 1$ for at least one k . \square

Lemma 7.3.11 *If A is transitive and $\alpha_{-k}, \dots, \alpha_k$ is admissible, that is, $a_{\alpha_i \alpha_{i+1}} = 1$ for $i = -k, \dots, k-1$, then the intersection $\Omega_A \cap C_{\alpha_{-k}, \dots, \alpha_k} =: C_{\alpha_{-k}, \dots, \alpha_k, A}$ is nonempty and moreover contains a periodic point.*

Proof Take m such that $a_{\alpha_k, \alpha_{-k}}^m > 0$. Then one can extend the sequence α to an admissible sequence of length $2k + m + 1$ that begins and ends with α_{-k} . Repeating this sequence periodically, we obtain a periodic point in $C_{\alpha_{-k}, \dots, \alpha_k, A}$. \square

Proposition 7.3.12 *If A is a transitive matrix, then the topological Markov chain σ_A is topologically mixing and its periodic orbits are dense in Ω_A ; in particular, σ_A is chaotic and hence has sensitive dependence on initial conditions.*

Proof The density of periodic orbits follows from Lemma 7.3.11. To prove topological mixing, pick open sets $U, V \subset \Omega_A$ and nonempty symmetric cylinders $C_{\alpha_{-k}, \dots, \alpha_k, A} \subset U$ and $C_{\beta_{-k}, \dots, \beta_k, A} \subset V$. Then it suffices to show that $\sigma_A^n(C_{\alpha_{-k}, \dots, \alpha_k, A}) \cap C_{\beta_{-k}, \dots, \beta_k, A} \neq \emptyset$ for any sufficiently large n . Take $n = 2k + 1 + m + l$ with $l \geq 0$, where m is as in Definition 7.3.9. Then $a_{\alpha_k, \beta_{-k}}^{m+l} > 0$ by Lemma 7.3.10, so there is an admissible sequence of length $4k + 2 + m + l$ whose first $2k + 1$ symbols are identical to $\alpha_{-k}, \dots, \alpha_k$ and the last $2k + 1$ symbols to $\beta_{-k}, \dots, \beta_k$. By Lemma 7.3.11, this sequence can be extended to a periodic element of Ω_A which belongs to $\sigma_A^n(C_{\alpha_{-k}, \dots, \alpha_k, A}) \cap C_{\beta_{-k}, \dots, \beta_k, A}$. \square

Example 7.3.13 The matrix $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ is not transitive because all its powers are upper triangular and hence there is no path from 1 to 0. In fact, the space Ω_A is countable and consists of two fixed points $(\dots, 0, \dots, 0, \dots)$ and $(\dots, 1, \dots, 1, \dots)$, and a single heteroclinic orbit connecting them (consisting of the sequences that are 1 up to some place and 0 thereafter).

Example 7.3.14 For the matrix

$$\begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \end{pmatrix}$$

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every orbit alternates between entries from the first group $\{0, 1\}$ on the one hand and from the second group $\{2, 3\}$ on the other hand, that is, the parity (even-odd) must alternate. Therefore no power of the matrix has all entries positive.

EXERCISES

■ **Exercise 7.3.1** Prove that E_2 has a nonperiodic orbit all of whose even iterates lie in the left half of the unit interval.

■ **Exercise 7.3.2** Prove that E_2 has a uncountably many orbits for which no segment of length 10 has more than one point in the left half of the unit interval.

■ **Exercise 7.3.3** Prove that linear maps that are conjugate in the sense of linear algebra are topologically conjugate in the sense of Definition 7.3.3.

■ **Exercise 7.3.4** Write down the Markov matrix for Figure 7.3.3 and check Corollary 7.3.6 up to period 3.

■ **Exercise 7.3.5** Consider the metric

$$(7.3.9) \quad d_\lambda^i(\alpha, \omega) := \sum_{i \in \mathbb{Z}} \frac{|\alpha_i - \omega_i|}{\lambda^{|i|}}$$

on Ω_N . Show that for $\lambda > 2N - 1$ the cylinder $C_{\alpha_1, \dots, \alpha_{n-1}}$ is a λ^{1-n} -ball for d_λ^i .

■ **Exercise 7.3.6** Repeat the previous exercise for one-sided shifts (with $\lambda > N$).

■ **Exercise 7.3.7** Consider the metric

$$(7.3.10) \quad d_\lambda''(\alpha, \omega) := \lambda^{-\max\{n \in \mathbb{N} \mid \alpha_i = \omega_i \text{ for } |i| \leq n\}}$$

[and $d_\lambda''(\alpha, \alpha) = 0$] on Ω_N . Show that the cylinder $C_{\alpha_1, \dots, \alpha_{n-1}}$ is a ball for d_λ'' .

■ **Exercise 7.3.8** Find the supremum of sensitivity constants for a transitive topological Markov chain with respect to the metric d_λ'' .

■ **Exercise 7.3.9** Find the supremum of sensitivity constants for a transitive topological Markov chain with respect to the metric d_λ'' .

■ **Exercise 7.3.10** Show that for $m < n$ the shift on Ω_m is a factor of the shift on Ω_n .

■ **Exercise 7.3.11** Prove that the quadratic map f_4 on $[0, 1]$ is not conjugate to any of the maps f_λ for $\lambda \in [0, 4)$.

■ **Exercise 7.3.12** Show that the topological Markov chains determined by the matrices

$$\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \quad \text{and} \quad \begin{pmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix}$$

are conjugate.

■ **Exercise 7.3.13** Find the smallest positive value of $p(\sigma_A)$ for a transitive topological Markov chain with two states (that is, with a 2×2 matrix A).

■ PROBLEMS FOR FURTHER STUDY

■ **Problem 7.3.14** Find all factors of an irrational rotation R_α of the circle.

■ **Problem 7.3.15** Find the smallest value of $p(\sigma_A)$ for a transitive topological Markov chain with three states (that is, with a 3×3 matrix A).

7.4 MORE EXAMPLES OF CODING

We now carry out a coding construction for several familiar dynamical systems.

7.4.1 Nonlinear Expanding Maps

There is a correspondence between general (not necessarily linear) expanding maps of the circle (Section 7.1.3) and a shift on a sequence space. The construction is similar to the one from Section 7.3.1. There is some effort involved, but there is a beautiful prize at the end: We obtain a complete classification of a large class of maps in terms of a simple invariant.

To keep notations simple, we consider an expanding map $f: S^1 \rightarrow S^1$ of degree 2. By Proposition 7.1.9, f has exactly one fixed point p . (For maps of higher degree, we could pick any one of the fixed points.) Since $\deg(f) = 2$, there is exactly one point $q \neq p$ such that $f(q) = p$. The points p and q divide the circle into two arcs. Starting from p in the positive direction, denote the first arc by Δ_0 and the second arc by Δ_1 . Define the coding for $x \in S^1$ as follows: x is represented by the sequence $\omega \in \Omega_2^R$ for which

$$(7.4.1) \quad f^n(x) \in \Delta_{\omega_n}.$$

This representation is unique unless $f^n(x) \in \{p, q\} = \Delta_0 \cap \Delta_1$. This lack of uniqueness is similar to the case of binary rationals for the map E_2 . Suppose a point x has an iterate in $\{p, q\}$. Then either $x = p$ and $f^n(x) = p$ for all $n \in \mathbb{N}$, or else the point q must appear before p in the sequence of iterates, that is, $f^n(x) \notin \{p, q\}$ for all n less than some k and then $f^k(x) = q$ and $f^{k+1}(x) = p$. In this case we make the following convention. p has two codes, all 0's and all 1's, and q has two codes, 01111111... and 1000000..., and any x such that $F^k(x) = q$ has two codes given by the first $k - 1$ digits uniquely defined by (7.4.1), followed by either of the codes for q .

Actually, going the other way around is better:

Proposition 7.4.1 *If $f: S^1 \rightarrow S^1$ is an expanding map of degree 2, then f is a factor of σ^R on Ω_2^R (Definition 7.3.3), that is, there is a surjective continuous map $h: \Omega_2^R \rightarrow S^1$ such that $f^n(h(\omega)) \in \Delta_{\omega_n}$ for all $n \in \mathbb{N}_0$, that is, $h \circ \sigma^R = f \circ h$.*

Proof That the domain of h is Ω_2^R requires that every sequence of 0's and 1's appears as the code of some point. First, f maps each of the two intervals Δ_0 and Δ_1 onto S^1 almost injectively, the only identification being at the ends. Let

$$\Delta_{00} \text{ be the core of } \Delta_0 \cap f^{-1}(\Delta_0),$$

$$\Delta_{01} \text{ be the core of } \Delta_0 \cap f^{-1}(\Delta_1),$$

$$\Delta_{10} \text{ be the core of } \Delta_1 \cap f^{-1}(\Delta_0),$$

$$\Delta_{11} \text{ be the core of } \Delta_1 \cap f^{-1}(\Delta_1).$$

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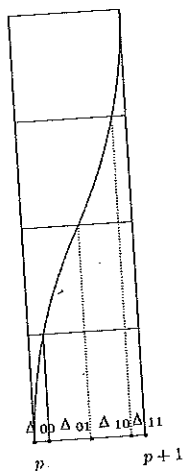


Figure 7.4.1. Nonlinear coding.

What we mean by "core" is that each indicated intersection consists of an interval as well as an isolated point (p or q), and we discard this extraneous point. Each of these four intervals is mapped onto S^1 by f^2 , again the only identification being at the ends. By definition, any point from Δ_{ij} has ij as the first two symbols of its code. Proceeding inductively we construct for any finite sequence $\omega_0, \dots, \omega_{n-1}$ the interval

$$(7.4.2) \quad \Delta_{\omega_0, \dots, \omega_{n-1}} := \text{the core of } \Delta_{\omega_0} \cap f^{-1}(\Delta_{\omega_1}) \dots \cap f^{1-n}(\Delta_{\omega_{n-1}}),$$

which is mapped by f^n onto S^1 with identification of the endpoints. Now take any infinite sequence $\omega = \omega_1, \dots \in \Omega_2^R$. The intersection $\bigcap_{n=1}^{\infty} \Delta_{\omega_0, \dots, \omega_{n-1}}$ of the nested closed intervals $\Delta_{\omega_0, \dots, \omega_{n-1}}$ is nonempty, and any point in this intersection has the sequence ω as its code.

So far we have only used the fact that f is a monotone map of degree 2. To show that h is well defined, we use the expanding property to check that $\bigcap_{n=1}^{\infty} \Delta_{\omega_0, \dots, \omega_{n-1}}$ consists of a single point, hence a point with a given code is unique.

If $g: I \rightarrow S^1$ is an injective map of an open interval I with a nonnegative derivative, then by the Mean-Value Theorem A.2.3 $l(g(I)) = \int_I g'(x) dx = g'(\xi)l(I)$ for some $\xi \in I$. Thus, in our case, there is a ξ_n such that

$$1 = l(S^1) = \int_{\Delta_{\omega_0, \dots, \omega_{n-1}}} (f^n)'(x) dx = (f^n)'(\xi_n) \cdot l(\Delta_{\omega_0, \dots, \omega_{n-1}}).$$

Since f is expanding $|(f^n)'| > \lambda^n$ for some $\lambda > 1$, hence $l(\Delta_{\omega_0, \dots, \omega_{n-1}}) < \lambda^{-n} \rightarrow 0$ as $n \rightarrow \infty$ and $\bigcap_{n=1}^{\infty} \Delta_{\omega_0, \dots, \omega_{n-1}}$ consists of a single point x_ω . This gives a well-defined surjective map $h: \Omega_2^R \rightarrow S^1, \omega \mapsto x_\omega$.

Give Ω_2^R the metric d_4 from (7.3.3). We showed in Section 7.3.4 that if $\epsilon = \lambda^{-n}$ and $\delta = 4^{-n}$, then $d(\omega, \omega') < \delta$ implies that $\omega_i = \omega'_i$ for $i < n$ and hence $|x_\omega - x_{\omega'}| \leq l(\Delta_{\omega_0, \dots, \omega_{n-1}}) < \lambda^{-n} = \epsilon$. Thus h is continuous. That $h(\sigma^R(\omega)) = f(h(\omega))$ is clear from the construction. \square

7.4.2 Classification via Coding

Proposition 7.4.1 and the discussion preceding it established a semiconjugacy between the one-sided 2-shift and the expanding map f on S^1 , that is,

Proposition 7.4.2 *Let $f: S^1 \rightarrow S^1$ be an expanding map of degree 2. Then f is a factor of the one-sided 2-shift (Ω_2^R, σ_R) via a semiconjugacy $h: \Omega_2^R \rightarrow S^1$. If $h(\omega) = h(\omega') =: x$, then there exists an $n \in \mathbb{N}_0$ such that $f^n(x) \in \{p, q\}$, where $p = f(p) = f(q)$, $q \neq p$.*

The last sentence of this proposition says that h is "very close" to being a conjugacy: There are only countably many image points where injectivity fails.

An important feature of this coding is that it is obtained in a uniform way for all expanding maps, and that the absence of injectivity occurs at points defined by their dynamics, namely, the fixed point and its preimages. This leads us to the prize promised at the beginning:

Theorem 7.4.3 *If $f, g: S^1 \rightarrow S^1$ are expanding maps of degree 2, then f and g are topologically conjugate; in particular, every expanding map of S^1 of degree 2 is conjugate to E_2 .*

Proof We have semiconjugacies $h_f, h_g: \Omega_2^R \rightarrow S^1$ for f and g . For $x \in S^1$, consider the set $H_x := h_g(h_f^{-1}(\{x\}))$. If x is a point of injectivity of h_f , that is, $h_f^{-1}(\{x\})$ is a single point, then so is H_x . Otherwise, x is a preimage of the fixed point under some iterate of f and $h_f^{-1}(\{x\})$ consists of a collection of sequences that are mapped under h_g to a single point. Therefore, H_x always consists of precisely one point $h(x)$. The bijective map $h: S^1 \rightarrow S^1$ thus defined is clearly a conjugacy: $h \circ f = g \circ h$. It is continuous because h_f sends open sets to open sets, that is, the image of a sequence and all sufficiently closeby sequences contains a small interval. Exchanging f and g shows that h^{-1} is also continuous. \square

This holds for any degree via an appropriate coding. It is the first major conjugacy result that establishes conjugacy with a specific model for all maps from a certain class. The Poincaré Classification Theorem 4.3.20 comes close, but requires extra assumptions (such as the existence of the second derivative; see Section 4.4.3) to produce a conjugacy with a rotation.

7.4.3 Quadratic Maps

For $\lambda > 4$ consider the quadratic map

$$f: \mathbb{R} \rightarrow \mathbb{R}, \quad x \rightarrow \lambda x(1 - x).$$

If $x < 0$, then $f(x) < x$ and $f'(x) > \lambda > 4$, so $f^n(x) \rightarrow -\infty$. When $x > 1$, $f(x) < 0$ and hence $f^n(x) \rightarrow -\infty$. Thus the set of points with bounded orbits is $\bigcap_{n \in \mathbb{N}_0} f^{-n}([0, 1])$.

Proposition 7.4.4 *If $\lambda > 2 + \sqrt{5}$ and $f: \mathbb{R} \rightarrow \mathbb{R}$, $x \rightarrow \lambda x(1 - x)$, then there is a homeomorphism $h: \Omega_2^R \rightarrow \Lambda := \bigcap_{n \in \mathbb{N}_0} f^{-n}([0, 1])$ such that $h \circ \sigma^R = f \circ h$, that is, $f|_\Lambda$ is conjugate to the 2-shift.*

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Proof Let

$$\Delta_0 = \left[0, \frac{1}{2} - \sqrt{\frac{1}{4} - \frac{1}{\lambda}} \right] \quad \text{and} \quad \Delta_1 = \left[\frac{1}{2} + \sqrt{\frac{1}{4} - \frac{1}{\lambda}}, 1 \right].$$

Then $f^{-1}([0, 1]) = \Delta_0 \cup \Delta_1$ by solving the quadratic equation $f(x) = 1$. Likewise, $f^{-2}([0, 1]) = \Delta_{00} \cup \Delta_{01} \cup \Delta_{11} \cup \Delta_{10}$ consists of four intervals, and so forth. Consider the partition of Λ by Δ_0 and Δ_1 . These pieces do not overlap and

$$\begin{aligned} |f'(x)| &= |\lambda(1 - 2x)| = 2\lambda \left| x - \frac{1}{2} \right| \geq 2\lambda \sqrt{\frac{1}{4} - \frac{1}{\lambda}} \\ &= \sqrt{\lambda^2 - 4\lambda} > \sqrt{(2 + \sqrt{5})^2 - 4(2 + \sqrt{5})} = 1 \end{aligned}$$

on $\Delta_0 \cup \Delta_1$. Thus, for any sequence $\omega = (\omega_0, \omega_1, \dots)$, the diameter of the intersections

$$\bigcap_{n=0}^N f^{-n}(\Delta_{\omega_n})$$

decreases (exponentially) as $N \rightarrow \infty$. This shows that for a sequence $\omega = (\omega_0, \omega_1, \dots)$ the intersection

$$(7.4.3) \quad h(\{\omega\}) = \bigcap_{n \in \mathbb{N}_0} f^{-n}(\Delta_{\omega_n})$$

consists of exactly one point and this map $h: \Omega_2^R \rightarrow \Lambda$ is a homeomorphism. \square

Remark 7.4.5 It turns out that Proposition 7.4.4 holds whenever $\lambda > 4$ (Proposition 11.4.1), but this is significantly less straightforward to prove than the present result. The situation present in either case, where a map folds an interval entirely over itself, is referred to as a one-dimensional *horseshoe*, in analogy to the geometry seen in the next subsection.

7.4.4 Linear Horseshoe

We now describe Smale's original "horseshoe," which provides one of the best examples of perfect coding. (In Section 7.3.3 a special case was constructed, in which ternary expansion provides the coding.)

Let Δ be a rectangle in \mathbb{R}^2 and $f: \Delta \rightarrow \mathbb{R}^2$ a diffeomorphism of Δ onto its image such that the intersection $\Delta \cap f(\Delta)$ consists of two "horizontal" rectangles Δ_0 and Δ_1 and the restriction of f to the components $\Delta^i := f^{-1}(\Delta_i)$, $i = 0, 1$, of $f^{-1}(\Delta)$ is a hyperbolic linear map, contracting in the vertical direction and expanding in the horizontal direction. This implies that the sets Δ^0 and Δ^1 are "vertical" rectangles. One of the simplest ways to achieve this effect is to bend Δ into a "horseshoe," or rather into the shape of a permanent magnet (Figure 7.4.2), although this method produces some inconveniences with orientation. Another way, which is better from the point of view of orientation, is to bend Δ roughly into a paper clip shape (Figure 7.4.3). This is an exaggerated version of the ternary horseshoe in Section 7.3.3, which also leaves some extra margin. If the horizontal and vertical rectangles lie strictly inside Δ , then the maximal invariant subset $\Lambda = \bigcap_{n=-\infty}^{\infty} f^{-n}(\Delta)$ of Δ is contained in the interior of Δ .

Proposition 7.4.6 $f|_{\Lambda}$ is topologically conjugate to σ_2 .

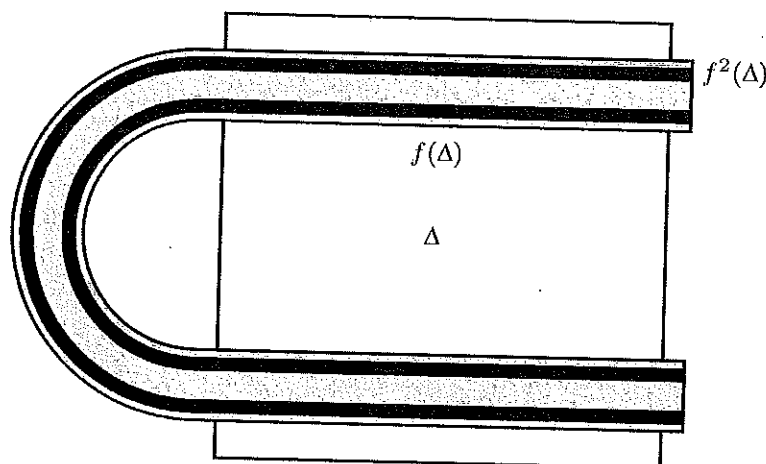


Figure 7.4.2. The horseshoe.

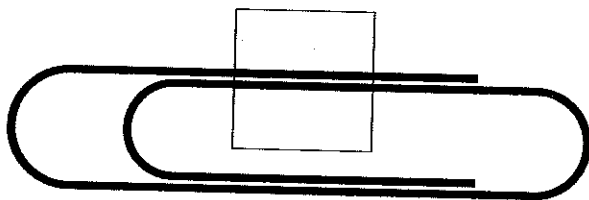


Figure 7.4.3. The paper clip.

Proof We use Δ^0 and Δ^1 as the “pieces” in the coding construction and start with positive iterates. The intersection $\Delta \cap f(\Delta) \cap f^2(\Delta)$ consists of four thin horizontal rectangles: $\Delta_{ij} = \Delta_i \cap f(\Delta_j) = f(\Delta^i) \cap f^2(\Delta^j)$, $i, j \in \{0, 1\}$ (see Figure 7.4.2). Continuing inductively, one sees that $\bigcap_{i=0}^n f^i(\Delta)$ consists of 2^n thin disjoint horizontal rectangles whose heights are exponentially decreasing with n . Each such rectangle has the form $\Delta_{\omega_1, \dots, \omega_n} = \bigcap_{i=1}^n f^i(\Delta^{\omega_i})$, where $\omega_i \in \{0, 1\}$ for $i = 1, \dots, n$. Each infinite intersection $\bigcap_{n=1}^{\infty} f^n(\Delta^{\omega_n})$, $\omega_n \in \{0, 1\}$, is a horizontal segment, and the intersection $\bigcap_{n=1}^{\infty} f^n(\Delta)$ is the product of the horizontal segment with a Cantor set in the vertical direction. Similarly, one defines and studies vertical rectangles $\Delta^{\omega_0, \dots, \omega_{-n}} = \bigcap_{i=0}^n f^{-i}(\Delta^{\omega_{-i}})$, the vertical segments $\bigcap_{n=0}^{\infty} f^{-n}(\Delta^{\omega_{-n}})$, and the set $\bigcap_{n=0}^{\infty} f^{-n}(\Delta)$, which is the product of a segment in the vertical direction with a Cantor set in the horizontal direction.

The desired invariant set $\Lambda = \bigcap_{n=-\infty}^{\infty} f^{-n}(\Delta)$ is the product of two Cantor sets and hence is a Cantor set itself (Problem 2.7.5), and the map

$$h: \Omega_2 \rightarrow \Lambda, \quad h(\{\omega\}) = \bigcap_{n=-\infty}^{\infty} f^{-n}(\Delta^{\omega_n})$$

is a homeomorphism that conjugates the shift σ_2 and the restriction of the diffeomorphism f to the set Λ . \square

Since periodic points and topological mixing are invariants of topological conjugacy, Proposition 7.4.6 and Proposition 7.3.4 immediately give substantial information about the behavior of f on Λ .

Corollary 7.4.7 *Periodic points of f are dense in Λ , $P_n(f|_\Lambda) = 2^n$, and the restriction of f to Λ is topologically mixing.*

Remark 7.4.8 Any map for which there is a perfect coding is defined on a Cantor set, because the perfect coding establishes a homeomorphism between the phase space and a sequence space, which is a Cantor set.

7.4.5 Coding of the Toral Automorphism

The idea of coding can be applied to hyperbolic toral automorphisms. To simplify notations and keep the construction more visual, we consider the standard example. Among our examples, this is the first where the coding is ingenious, even though it is geometrically simple. Section 10.3 describes a construction whose dynamical implications are quite similar to those obtained here, but where the geometry is complicated and almost always fractal.

Theorem 7.4.9 *For the map*

$$F(x, y) = (2x + y, x + y) \pmod{1}$$

of the 2-torus from Section 7.1.4 there is a semiconjugacy $h: \Omega_A \rightarrow \mathbb{T}^2$ with

$$F \circ h = h \circ \sigma_5|_{\Omega_A}, \quad \text{where}$$

$$(7.4.4) \quad A = \begin{pmatrix} 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 \end{pmatrix}.$$

Proof Draw segments of the two eigenlines at the origin until they cross sufficiently many times and separate the torus into disjoint rectangles. Specifically, extend a segment of the contracting line in the fourth quadrant until it intersects the segment of the expanding line twice in the first quadrant and once in the third quadrant (see Figure 7.4.4). The resulting configuration is a decomposition of the torus into two rectangles $R^{(1)}$ and $R^{(2)}$. Three pairs among the seven vertices of the plane configuration are identified, so there are only four different points on the torus that serve as vertices of the rectangles; the origin and three intersection points. Although $R^{(1)}$ and $R^{(2)}$ are not disjoint, one can apply the method used for the horseshoe, using $R^{(1)}$ and $R^{(2)}$ as basic rectangles. The expanding and contracting eigendirections play the role of the "horizontal" and "vertical" directions, correspondingly. Figure 7.4.5 shows that the image $F(R^{(i)})$ ($i = 1, 2$) consists of several "horizontal" rectangles of full length. The union of the boundaries $\partial R^{(1)} \cup \partial R^{(2)}$ consists of the segments of the two eigenlines at the origin just described. The image of the contracting segment is a part of that segment. Thus, the images of $R^{(1)}$ and $R^{(2)}$ have to be "anchored" at parts of their "vertical" sides; that is, once one of the images "enters" either $R^{(1)}$ or $R^{(2)}$, it has to stretch all the way through it. By matching things up along the contracting direction one sees that $F(R^{(1)})$ consists of three components, two in $R^{(1)}$ and one in $R^{(2)}$. The image of $R^{(2)}$ has two components, one in each

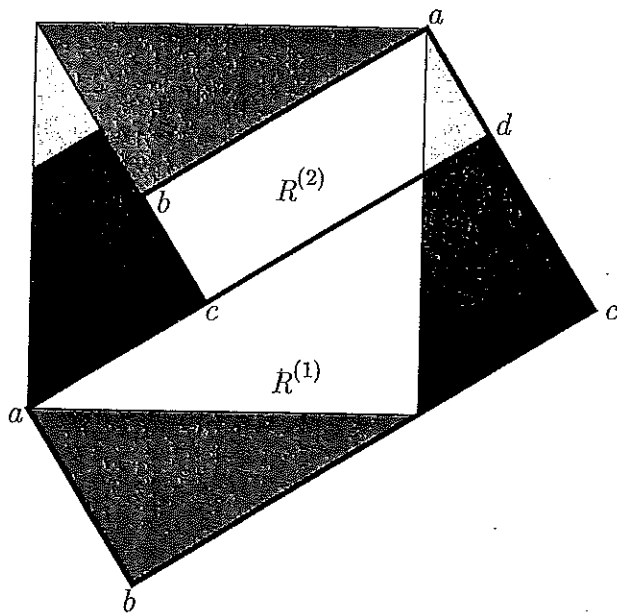


Figure 7.4.4. Partitioning the torus.

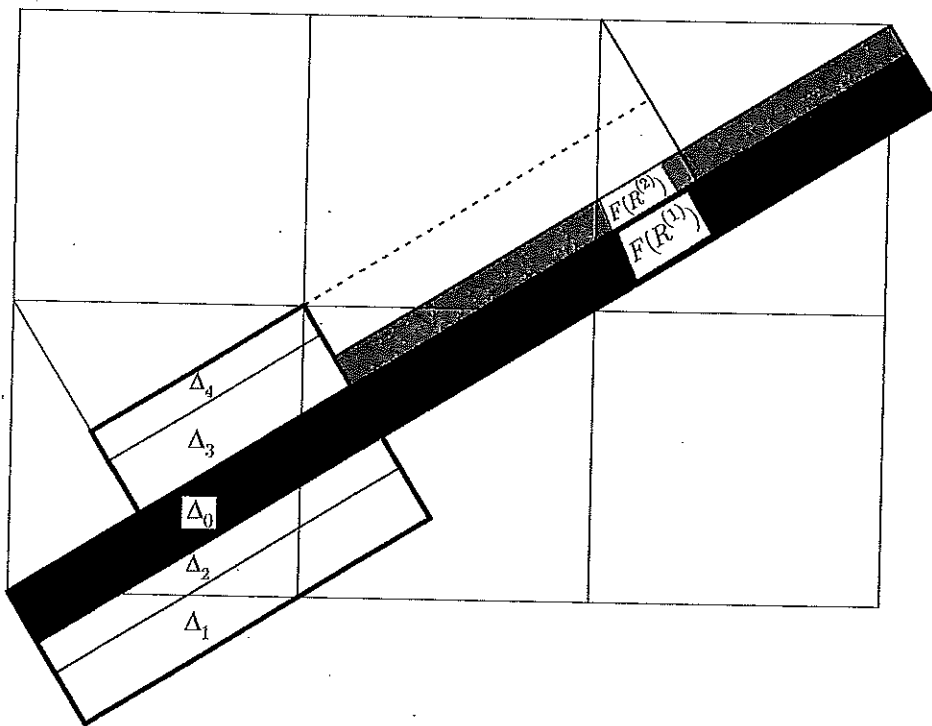


Figure 7.4.5. The image of the partition.

rectangle (see Figure 7.4.5). The fact that $F(R^{(1)})$ has two components in $R^{(1)}$ would cause problems if we were to use $R^{(1)}$ and $R^{(2)}$ for coding construction (more than one point for some sequences), but we use these five components $\Delta_0, \Delta_1, \Delta_2, \Delta_3, \Delta_4$ (or their preimages) as the pieces in our coding construction. There is exactly one rectangle $\Delta_{\omega_{-\ell} \dots \omega_0 \omega_1 \dots \omega_k}$ defined by $\bigcap_{n=-\ell}^k F^{-n}(\Delta_{\omega_n})$, not several. (As in the case of expanding maps in Section 7.3.1, we have to discard extraneous pieces, in this case line segments.) Due to the contraction of F in the "vertical" direction, $\Delta_{\omega_{-\ell} \dots \omega_0 \omega_1 \dots \omega_k}$ has "height" less than $((3 - \sqrt{5})/2)^\ell$, and due to the contraction of F^{-1} in the "horizontal" direction $\Delta_{\omega_{-\ell} \dots \omega_0 \omega_1 \dots \omega_k}$ has "width" less than $((3 - \sqrt{5})/2)^k$. These go to zero as $\ell \rightarrow \infty$ and $k \rightarrow \infty$, so the intersection $\bigcap_{n \in \mathbb{Z}} F^{-n}(\Delta_{\omega_n})$ defines at most one point $h(\omega)$. On the other hand, because of the "Markov" property described previously, that is, the images going full length through rectangles, the following is true: If $\omega \in \Omega_5$ and $F^{-1}(\Delta_{\omega_n})$ overlaps $\Delta_{\omega_{n+1}}$ for all $n \in \mathbb{Z}$, then there is such a point $h(\omega)$ in $\bigcap_{n \in \mathbb{Z}} F^{-n}(\Delta_{\omega_n})$. Thus, we have a coding, which, however, is not defined for all sequences of Ω_5 .

Instead, we have to restrict attention to the subspace Ω_A of Ω_5 that contains only those sequences where any two successive entries constitute an "allowed transition", that is, 0, 1, 2 can be followed by 0, 1, or 3, and 3 and 4 can be followed by 2 or 4. This is exactly the topological Markov chain (Definition 7.3.2) for (7.4.4). \square

Theorem 7.4.10 *The semiconjugacy between σ_A and F is one-to-one on all periodic points except for the fixed points. The number of preimages of any point not negatively asymptotic to the fixed point is bounded.*

Proof We describe carefully the identifications arising from our semiconjugacy, that is, what points on the torus have more than one preimage. First, obviously, the topological Markov chain σ_A has three fixed points, namely, the constant sequences of 0's, 1's, and 4's, whereas the toral automorphism F has only one, the origin. It is easy to see that all three fixed points are indeed mapped to the origin. As we have seen in Proposition 7.1.10, $P_n(F) = \lambda_1^n + \lambda_1^{-n} - 2$, and accordingly $P_n(\sigma_A) = \text{tr } A^n = \lambda_1^n + \lambda_1^{-n} = P_n(F) + 2$ (Corollary 7.3.6), where $\lambda_1 = (3 + \sqrt{5})/2$ is the maximal eigenvalue for both the 2×2 matrix $\begin{pmatrix} 2 & 1 \\ 1 & 1 \end{pmatrix}$ and for the 5×5 matrix (7.4.4). To see that the eigenvalues are the same, consider $A - \lambda \text{Id}$, subtract column 4 from the first two columns and column 5 from the third, and then add rows 1 and 2 to row 4 and row 3 to row 5:

$$\begin{pmatrix} 1-\lambda & 1 & 0 & 1 & 0 \\ 1 & 1-\lambda & 0 & 1 & 0 \\ 1 & 1 & -\lambda & 1 & 0 \\ 0 & 0 & 1 & -\lambda & 1 \\ 0 & 0 & 1 & 0 & 1-\lambda \end{pmatrix} \rightarrow \begin{pmatrix} -\lambda & 0 & 0 & 1 & 0 \\ 0 & -\lambda & 0 & 1 & 0 \\ 0 & 0 & -\lambda & 1 & 0 \\ \lambda & \lambda & 0 & -\lambda & 1 \\ 0 & 0 & \lambda & 0 & 1-\lambda \end{pmatrix} \\ \rightarrow \begin{pmatrix} -\lambda & 0 & 0 & 1 & 0 \\ 0 & -\lambda & 0 & 1 & 0 \\ 0 & 0 & -\lambda & 1 & 0 \\ 0 & 0 & 0 & 2-\lambda & 1 \\ 0 & 0 & 0 & 1 & 1-\lambda \end{pmatrix}$$

Furthermore, one can see that every point $q \in \mathbb{T}^2$ whose positive and negative iterates avoid the boundaries $\partial R^{(1)}$ and $\partial R^{(2)}$ has a unique preimage, and vice versa. In particular, periodic points other than the origin (which have rational coordinates) fall into this category. The points of Ω_A whose images are on those boundaries or their iterates under F fall into three categories corresponding to the three segments of stable and unstable manifolds through 0 that define parts of the boundary. Thus sequences are identified in the following cases: They have a constant infinite right (future) tail consisting of 0's or 4's, and agree otherwise – this corresponds to a stable boundary piece – or else an infinite left (past) tail (of 0's and 1's, or of 4's), and agree otherwise – this corresponds to an unstable boundary piece. \square

■ EXERCISES

■ **Exercise 7.4.1** Prove that for $\lambda \geq 1$ every bounded orbit of the quadratic map f_λ is in $[0, 1]$.

■ **Exercise 7.4.2** Give a detailed argument that (7.4.3) defines a homeomorphism.

■ **Exercise 7.4.3** Construct a Markov partition for $\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$ that consists of two squares.

■ **Exercise 7.4.4** Construct a Markov partition and describe the corresponding topological Markov chain for the automorphism F_L , where $L = \begin{pmatrix} 1 & 1 \\ 2 & 1 \end{pmatrix}$.

■ **Exercise 7.4.5** Given a 0-1 $n \times n$ -matrix A , describe a system of n rectangles $\Delta_1, \dots, \Delta_n$ in \mathbb{R}^2 and map $f: \Delta := \bigcup_{i=1}^n \Delta_i \rightarrow \mathbb{R}^2$ such that the restriction of f to the set of points that stay inside Δ for all iterates of f is topologically equivalent to the topological Markov chain σ_A .

■ **Exercise 7.4.6** Check that the process (7.4.2) of discarding extraneous points in the coding construction amounts to taking $\Delta_{\omega_0, \dots, \omega_{n-1}} = \bigcap_{i=0}^{n-1} \text{Int}(f^{-i}(\Delta_{\omega_i}))$, and $\{h(\omega)\} := \bigcap_{n \in \mathbb{N}} \Delta_{\omega_0, \dots, \omega_{n-1}}$.

■ PROBLEMS FOR FURTHER STUDY

■ **Problem 7.4.7** Show that the assertion of Theorem 7.4.3 remains true for any map f of degree 2 such that $f' \geq 1$ and $f' = 1$ only at finitely many points.

■ **Problem 7.4.8** Prove the assertion of Theorem 7.4.9 for some 0-1 matrix A for any automorphism

$$F_L: \mathbb{T}^2 \rightarrow \mathbb{T}^2, x \mapsto Lx \pmod{1},$$

where L is an integer 2×2 matrix with determinant $+1$ or -1 and with real eigenvalues different from ± 1 .

7.5 UNIFORM DISTRIBUTION

We now investigate whether the notion of the uniform distribution of orbits that appeared in previous chapters for rotations of the circle and translations of the torus has any meaning for the group of examples discussed in the present chapter, such as linear or nonlinear expanding maps of the circle, shifts, and automorphisms of the torus.