# Homework 2

## due January 31 2017, 23:59

#### Task 1: Machine Epsilon

In general a computer stores a real number in the following way

$$x = (-1)^{s} \cdot (0.a_{1}a_{2}...a_{t}) \cdot \beta^{e} = (-1)^{s} \cdot m \cdot \beta^{e-t}, \quad a_{1} \neq 0$$

where s is either 0 or 1,  $\beta$  (a positive integer larger than or equal to 2) is the basis adopted by the specific computer at hand, m is an integer called mantissa whose length t is the maximum number of digits  $a_i$  (with  $0 \le a_i \le \beta - 1$ ) that are stored, and e is an integral number called the exponent. The numbers given in this form are called floating-point numbers, since the position of the decimal point is not fixed. The digits  $a_1a_2...a_p$  (with  $p \le t$ ) are called the p first significant digits of x. The accuracy with which floating-point numbers are stored depends then on  $\beta$  and t, and so does the amount of memory required to store them. For example, double precision real numbers are stored in registers of 8 Bytes: the sign s is stored in 1 bit, the exponent e in 11 bits, and the mantissa m in 52 bits. Note that, although there are 52 bits for m, we can count t = 53 digits when  $\beta = 2$ . As a matter of fact, since the first digit  $a_1$  of every floating point number must be different from 0, when  $\beta = 2$  it is worthless to store it as it must necessarily be 1. A round-off error is inevitably generated whenever a real number  $x \neq 0$  is replaced by its floating-point representative  $x_{num}$ , this error is always limited by

$$\frac{|x - x_{\text{num}}|}{|x|} \le \frac{1}{2}\varepsilon,$$

where  $\varepsilon = \beta^{1-t}$ , called machine epsilon.

The following code can be used in MATLAB to compute  $\varepsilon$ .

```
numprec=double(1.0); % Define 1.0 with double precision
numprec=single(1.0); % Define 1.0 with single precision
while(1 < 1 + numprec)
    numprec=numprec*0.5;
end
numprec=numprec*2</pre>
```

- a) Determine  $\varepsilon$  using the above program, both for single and double precision.
- b) Explain in detail what the code does. Why do we consider addition to 1?
- c) Explain the difference between single and double precision. How many Bytes are used to store a single precision number? How many for the mantissa?

### Task 2: Round-off Error

In this exercise, the errors involved in the numerical approximation of derivatives are examined. Using central finite differences the derivative of a function f(x) can be approximated as:

$$f'(x) \approx f'_{\text{num}}(x) = \frac{f(x + \Delta x) - f(x - \Delta x)}{2\Delta x} \tag{1}$$

a) Compute, numerically, the relative discretization error of the derivative of the function  $f(x) = \frac{1}{1+x} + x$  using equation (1). The relative discretization error is given by:

$$\xi_{\rm d} = \frac{|f'(x) - f'_{\rm num}(x)|}{|f'(x)|},\tag{2}$$

where f'(x) is the analytical derivative of f(x). Compute  $\xi_d$  at x = 2 for different stepsizes  $\Delta x \in [10^{-20}, 1]$ . Use both single and double precision for the calculation and present the results in a double logarithmic plot<sup>1</sup> ( $\xi_d$  vs.  $\Delta x$ ). Remember that *all* variables used here should be defined as double or single precision as in Task 1.

b) The absolute propagation error of an arithmetic operation  $\circ (+, -, \times \text{ or } /)$  between two numbers  $a_1$  and  $a_2$  can be evaluated as:  $a_1 \circ a_2 - a_{1,\text{num}} \circ a_{2,\text{num}}$ , where  $(\cdot)_{\text{num}}$  is the machine representation of the respective number.

Show that the propagation error of the addition of two positive numbers  $a_1$  and  $a_2$  is given by

$$\xi_{\rm p,add} = \frac{a_1}{a_1 + a_2} \varepsilon_{a_1} + \frac{a_2}{a_1 + a_2} \varepsilon_{a_2},\tag{3}$$

where  $\varepsilon_{a_j} := (a_j - a_{j,\text{num}})/a_j$  is the machine accuracy on the quantity  $a_{j,\text{num}}$ .

A general formula for the propagation error for a function  $g(a_1, a_2, \ldots, a_n)$  representing multiple arithmetic operations is given by:

$$\xi_{\rm p} = \sum_{j=1}^{n} \left| \frac{a_j}{g} \frac{\partial g}{\partial a_j} \right| \varepsilon_{a_j},\tag{4}$$

Show that when  $g = a_1 + a_2$  this formula results in equation (3).

c) Show that, when using the proposed central differences approximation, the relative discretization error (equation (2)) is given by:

$$\xi_{\rm d} \approx \frac{\Delta x^2 |f'''(x)|}{6|f'(x)|}$$

(Hint: Taylor expansion)

and that the propagation error (equation (4)) is given by:

$$\xi_{\rm p} \approx \frac{|f(x)|\varepsilon}{|f'(x)|\Delta x}$$

where  $\varepsilon$  is the machine accuracy. Find, analytically, the value of  $\Delta x$  that minimizes the total error

$$\xi_{\rm tot} = \xi_{\rm d} + \xi_{\rm p}$$

Plot  $\xi_d, \xi_p$  and  $\xi_{tot}$  together with the results from part a).

<sup>&</sup>lt;sup>1</sup>In MATLAB double logarithmic plots are obtained by the function loglog().

## Task 3 : Discretization in time

In this problem the stability and convergence order of three numerical time discretization methods is examined. Consider the first order, linear, test equation (the Dahlquist equation)

$$\begin{cases} u'(t) = f(u) = \lambda u(t), & 0 < t \le T, \\ u(0) = 1 \end{cases}$$
(5)

where  $\lambda = \lambda_{\Re} + i\lambda_{\Im} \in \mathbb{C}$ . The time interval [0, T] is discretized into N equally spaced parts:  $t_n = n\Delta t, n = 0, 1, \dots, N$ , where  $\Delta t$  is the step-size. The following numerical methods should be used:

• explicit Euler

$$u^{n+1} - u^n = \Delta t f(u^n)$$

• implicit Euler

$$u^{n+1} - u^n = \Delta t f(u^{n+1})$$

• Crank-Nicolson

$$u^{n+1} - u^n = \frac{1}{2}\Delta t \left[ f(u^{n+1}) + f(u^n) \right]$$

where  $u^n := u(t_n)$ .

- a) Solve the system (5) analytically (by hand) to obtain the exact solution  $u = u_{ex}$ .
- b) For  $\lambda = -\sqrt{3}/2 + i\pi$  and for the five cases N = 20, 40, 50, 100, and 200, compute the numerical solution iteratively until T = 10 for all the three methods. Plot the real part of the solutions together with the exact solution for each value of N. What do you observe?
- c) Now, consider  $\lambda \in \mathbb{R}$ . For each of the three considered schemes: (i) derive the expression of the amplification factor G(z), where  $z := \lambda \Delta t$ ; (ii) calculate  $\lim_{z \to -\infty} G(z)$ ; (iii) plot G(z) as a function of z together with the result for the exact amplification over the interval  $z \in [-10, 0.5]$ . Discuss the performance of the schemes in the limits  $z \to -\infty$  and  $z \to 0$ . Also, answer: why is the imaginary part of  $\lambda$  irrelevant for this analysis?
- d) For  $\lambda = -\sqrt{3}/2 + i$ , first do as in b) and explain the differences. Then, for each method, at a fixed time (chose t = 3) compute and plot the error  $|u_{\text{ex}} u_{\text{num}}|$  as a function of N in a double logarithmic plot and estimate the order of accuracy by considering the slope of the curve. (Hint:  $\log(x^p) = p \log(x)$ .)
- e) Based this task, discuss the usefulness, stability and accuracy of the methods.