

# Huffman Coding

...

Prefix-free code & lossless data compression

# Compression

## Lossless

- Less effective
- Universal & domain Specific
- 100 % Quality remains

### Examples:

- ZIP
- Huffman Coding
- Lempel-Ziv-Welch (LZW)
- PNG

## Lossy

- More effective compression
- Domain specific only
- Quality lost

### Examples:

- MP3
- JPEG
- H.264
- Dolby Digital

# Text encoding: ASCII (*American Standard Code for Information Interchange*)

- 1 Character = 1 Byte
- Contains 127 different characters
- Does not contain Å Ä Ö

## Example:

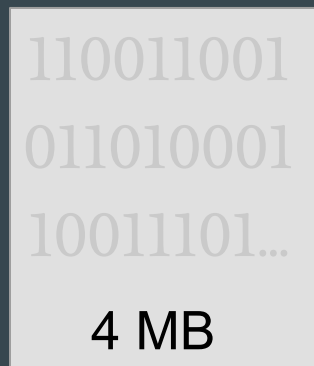
A = 0x41 = 0b0100 0001

Fox = 0x46 6F 78 = 0b0100 0110 01100 01111 0111 1000

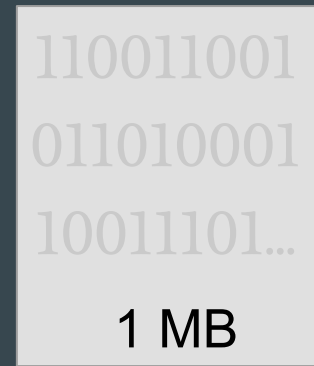
Dec	Hex	Char
..	..	..
74	4A	J
75	4B	K
76	4C	L
..	...	..

# Compression & Decompression

Original Data



Compressed Data



# Huffman Compression

Original Data

“Voodoo”



Protocol

<i>Char</i>	<i>Representation</i>
V	01
o	1
d	00



Compressed Data

0b01110011

# Prefix-free code

OK!

Char	Representation
V	01
o	1
d	00

Voodoo == 0b01110011

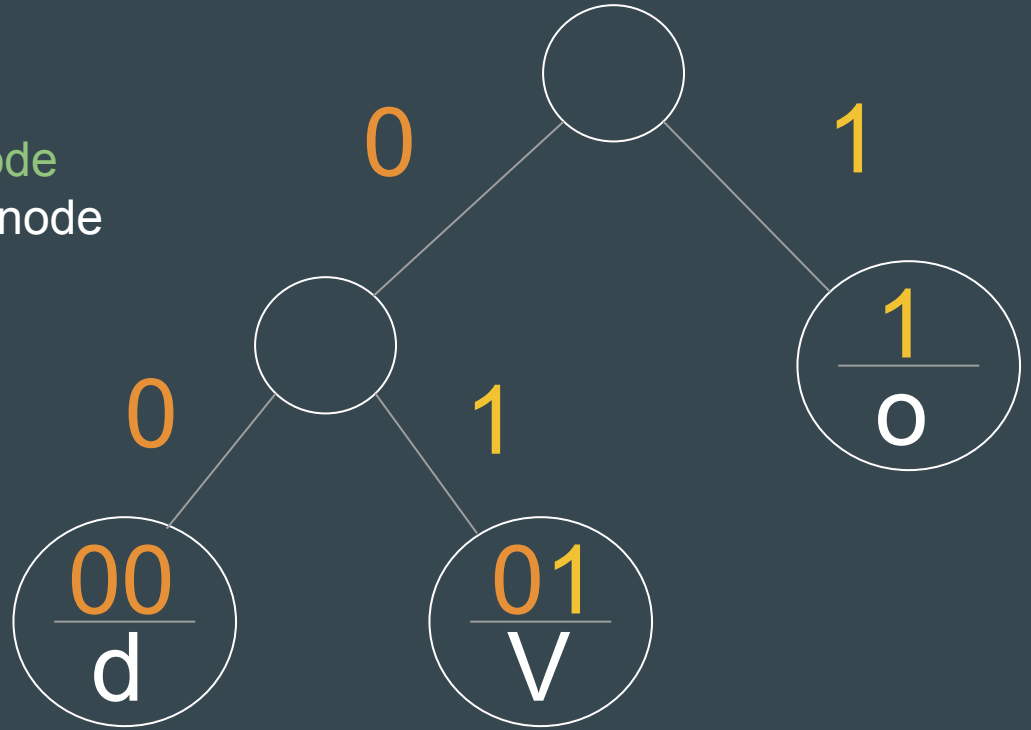
WRONG!

Char	Representation
V	11
o	1
d	01

Voodoo == 0b11110111  
oVoodoo == 0b11110111  
oVodoV == 0b11110111

# Huffman Tree

- Every leaf node is a character node
- A character node must be a leaf node
- Left == 0
- Right == 1



# Huffman Tree

“Voodoo”

<i>Char</i>	<i>Occurrences</i>
<b>v</b>	1
<b>o</b>	4
<b>d</b>	1

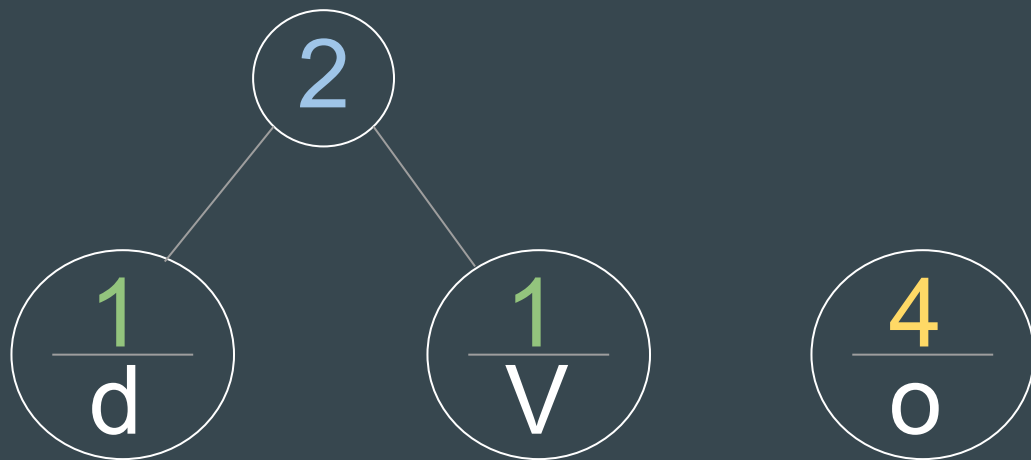




# Huffman Tree

“Voodoo”

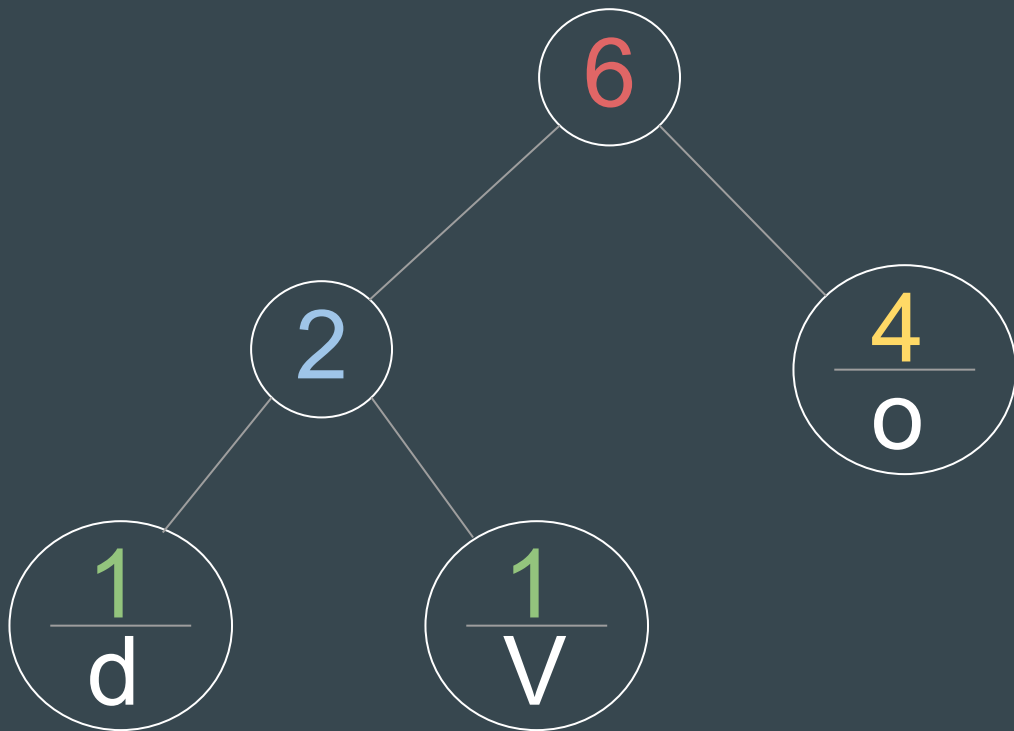
<i>Char</i>	<i>Occurrences</i>
<b>v</b>	1
<b>o</b>	4
<b>d</b>	1



# Huffman Tree

“Voodoo”

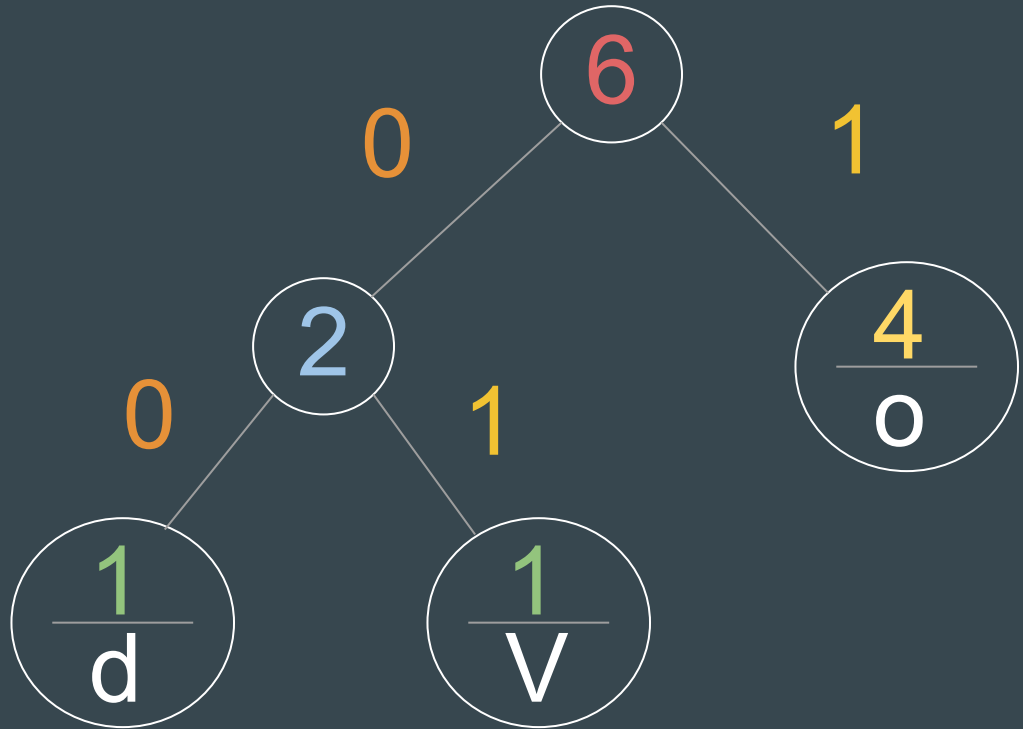
<i>Char</i>	<i>Occurrences</i>
<b>v</b>	1
<b>o</b>	4
<b>d</b>	1



# Huffman Tree

“Voodoo”

<i>Char</i>	<i>Representation</i>
<b>v</b>	01
<b>o</b>	1
<b>d</b>	00



# Huffman Compression

1. Calculate Occurrences
2. Huffman Tree
3. Protocol
4. Compress